

## Inside the Octagon

Analyzing System Guard Runtime Attestation

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## Shout outs

#### System Guard is the result of many many people's hard work!

The entire "Octagon" team

The DRTM team across Windows

Enclaves and VBS team

Microsoft Core UEFI team

Intel, AMD, QC teams

## Why attestation?

All code executes with integrity.

User identities cannot be compromised, spoofed, or stolen.

Attacker with casual physical access cannot modify data or code on the device.













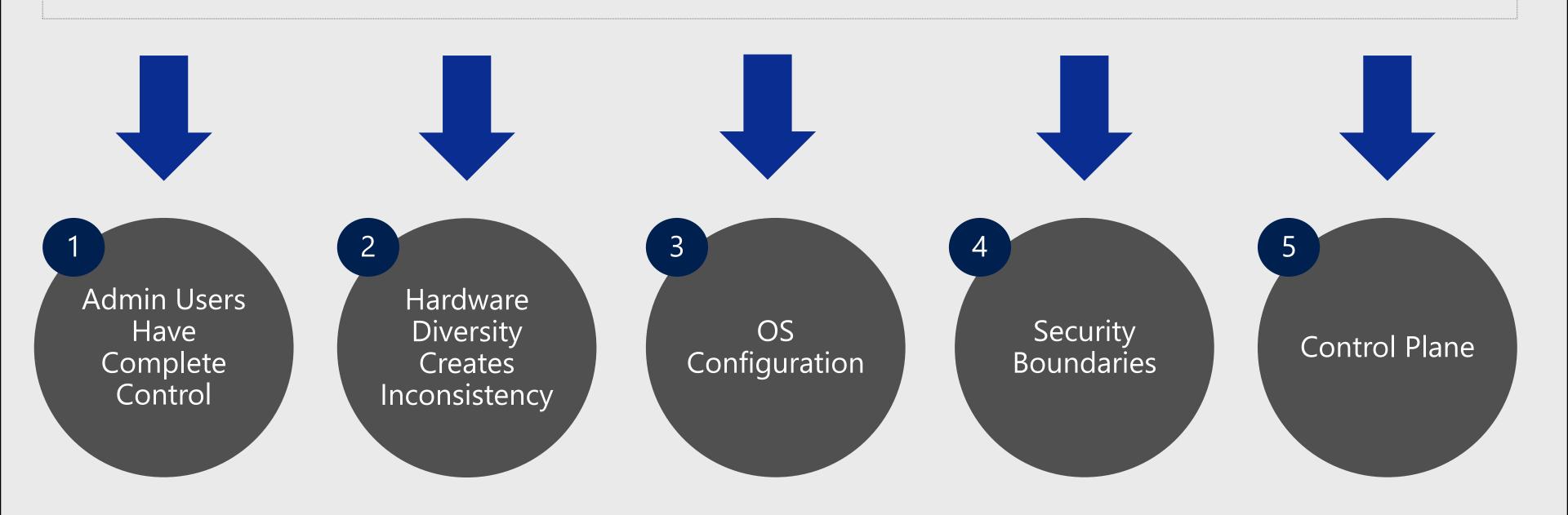
Malicious code cannot persist on a device.

Violations of promises are observable.

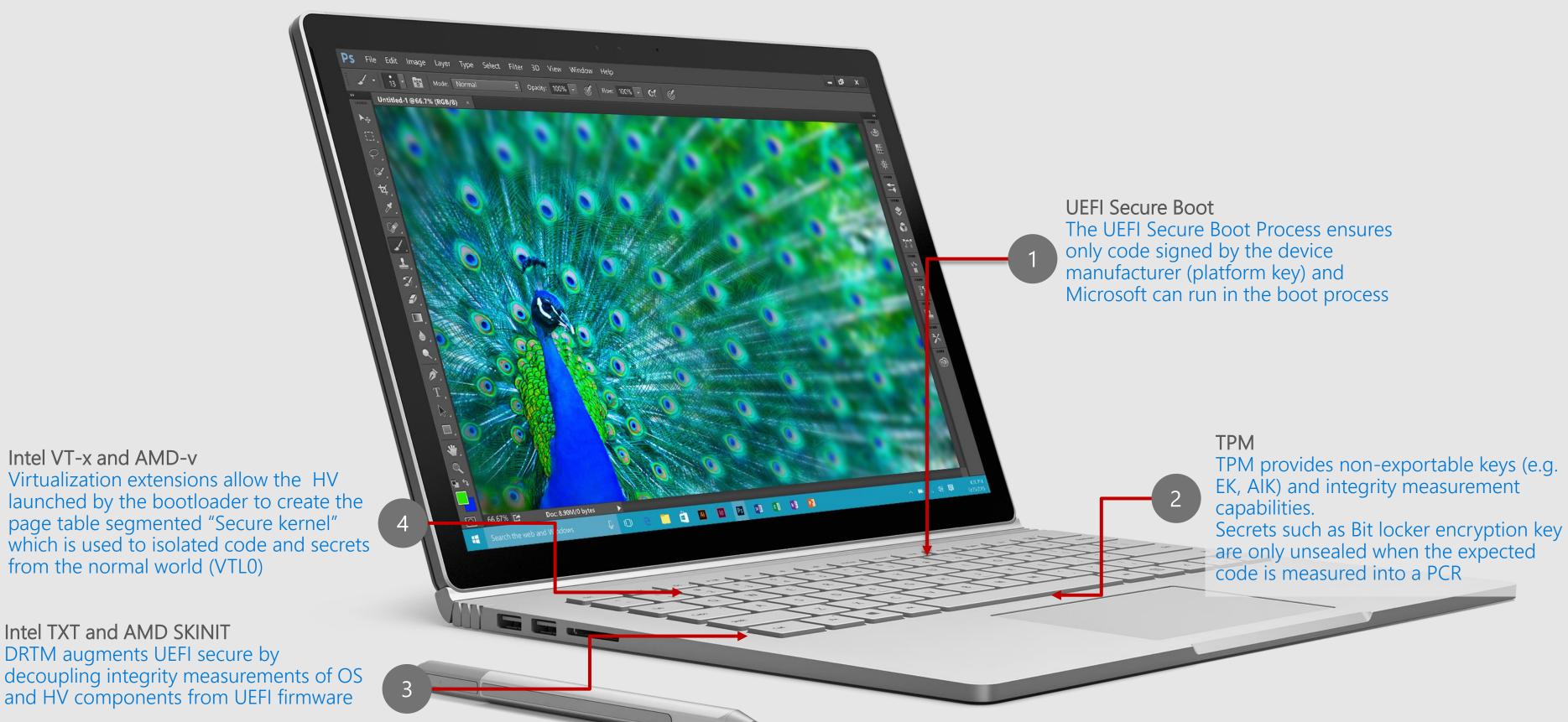
All apps and system components have only the privilege they need.



## Trust Challenges



## Hardware Anchored Trust

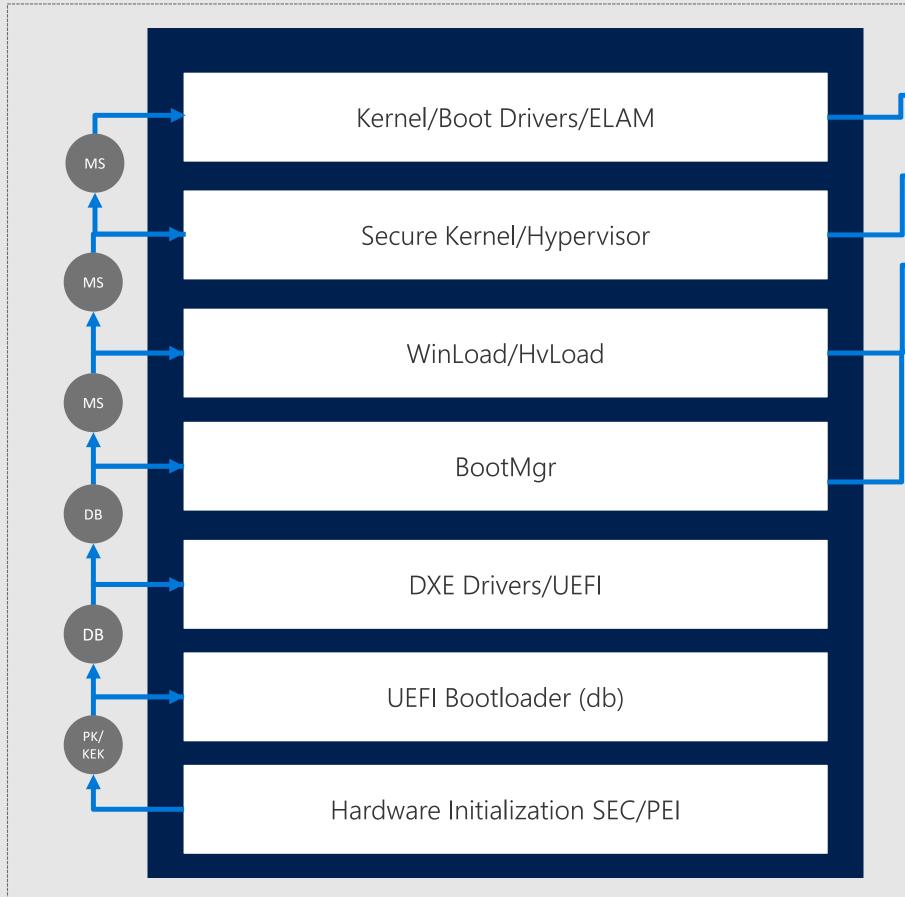


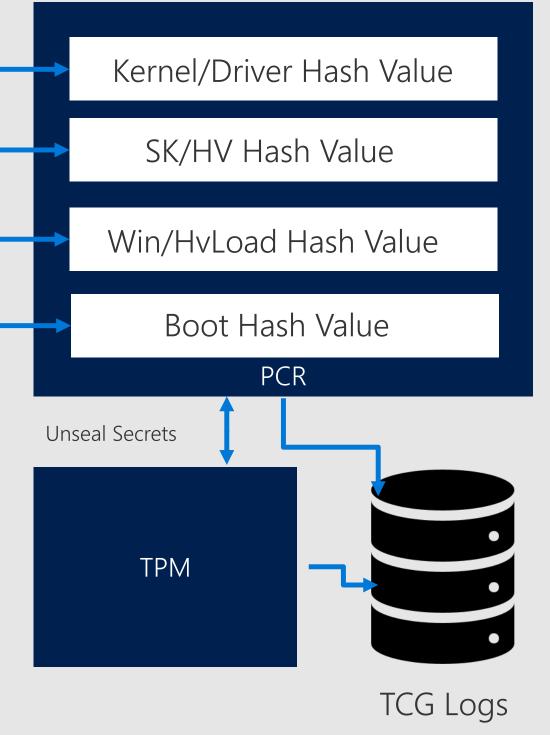
#### Intel VT-x and AMD-v

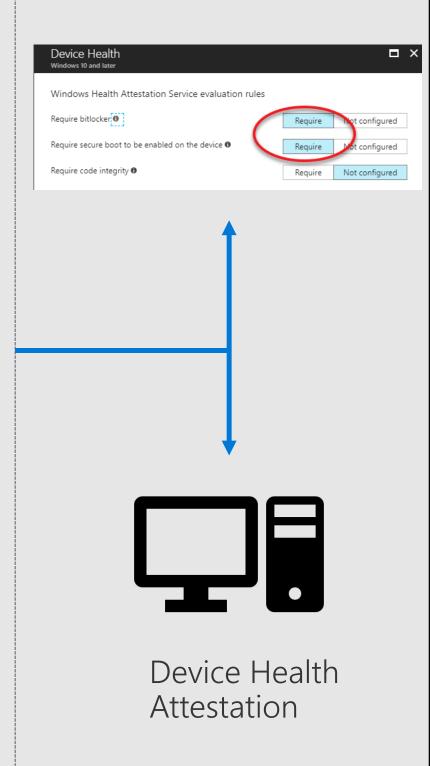
Virtualization extensions allow the HV launched by the bootloader to create the page table segmented "Secure kernel" which is used to isolated code and secrets from the normal world (VTL0)

Intel TXT and AMD SKINIT DRTM augments UEFI secure by decoupling integrity measurements of OS

## Attesting to Trust







See <a href="here">here</a> for additional details

## Limitations of current trust model

#### Secure boot and static root of trust (SRTM)

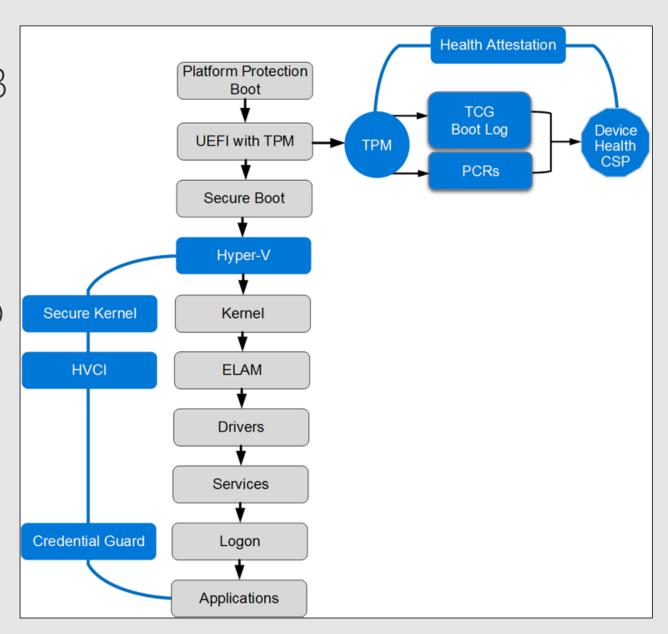
Strong dependency on the security UEFI firmware Widespread issues with OEM firmware are used to undermine TCB

#### High-integrity attestation is at boot time

DHA only measures and attests to the integrity of system bring up Simple runtime attacks do not impact health state

#### Limited runtime solutions

EDR/AV can identify runtime attacks limited to integrity (software) PatchGuard/HyperGuard great but not easily extensible



## Building a better model

## What is System Guard Runtime Attestation?

#### Platform Tamper Detection for Windows

Spanning device boot to ongoing runtime tampering

Designed for remote assessment of device health

Platform-driven approach to benefit a variety of 3<sup>rd</sup> party customers and scenarios

#### Why are we building it?

Challenging to build tamper detection schemes on top of Windows

Leverage the VBS security boundary to raise the bar on anti-tampering

Evolving security approach to keep Windows users and their data safe

#### How Does it Work?

Boot time attestation assures integrity of Windows & environment where tamper detection sensor runs

Sensor runs in isolated VBS environment constantly monitoring User/Kernel process space for signs of tampering

Relying parties calling into runtime attestation APIs in normal world get integrity protected attestation reports

Relying party backend verifies attestation report to ascertain device tamper status

## Improving Boot Trust

#### System Guard with DRTM

Utilize DRTM (Intel, AMD, QC) to measure TCB from a Microsoft MLE "Assume Breach" of UEFI and measure it from a hardware-rooted MLE

#### DRTM + Remote Attestation

Measurements of key properties available in PCRs Attest TCB components w/ System Guard runtime attestation + Azure Conditional Access + WDATP

#### SMM Attacks

Can be used to tamper HV and SK post-MLE

<u>SMM paging protections</u> + attestation on roadmap



Security features available on your device that use virtualization-based security.

This setting is managed by your administrator.

#### Memory integrity

Prevents attacks from inserting malicious code into high-security processes.



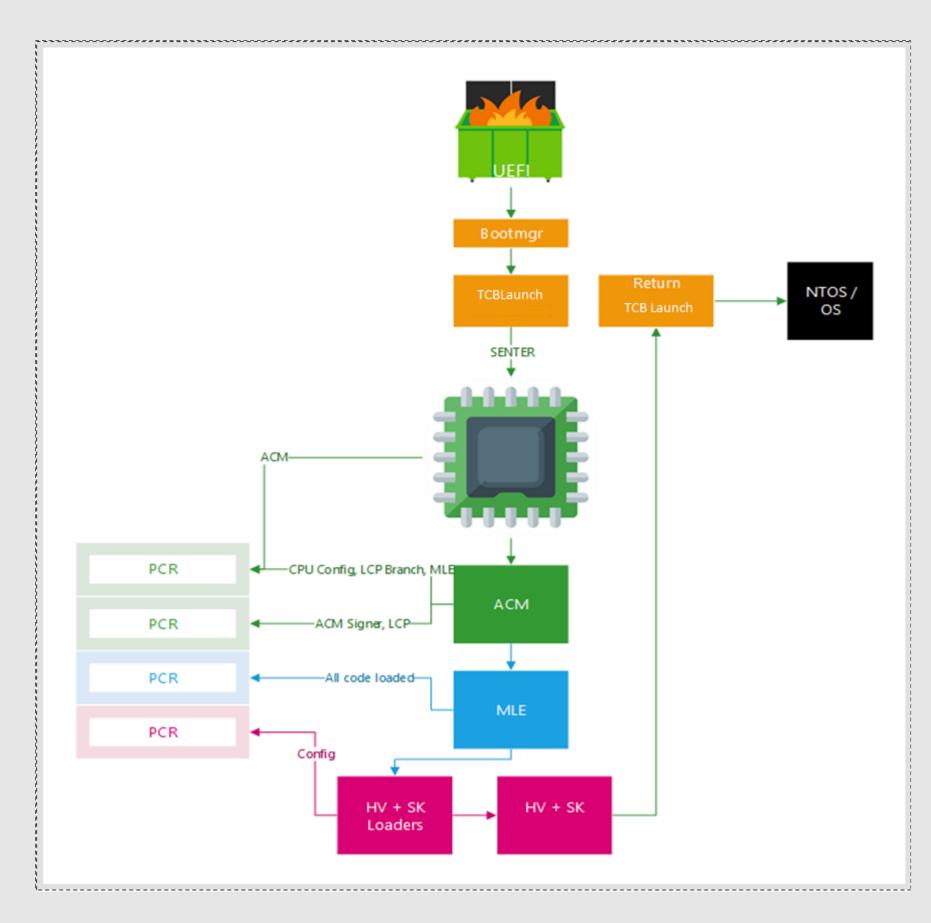
#### Learn more

#### Firmware protection

Windows Defender System Guard is protecting your device from compromised firmware.

#### Learn more

## System Guard - DRTM



## Regular Launch

Boot manager (**BootMgr**) is responsible for initial platform data collection Parses Secure Boot policy, ACPI/UEFI runtime tables and parts of registry Initializes BitLocker, displays boot/recovery menu Launches Windows loader, resume, or other boot app (such as **MemTest**)

# Windows loader (WinLoad) does the remainder of the data collection Constructs large structure (loader parameter block) to pass to the kernel In the case of Hyper-V load, passes data to HvLoader instead HvLoader initializes Hyper-V, Secure Kernel (for VSM/VBS), then returns

## Entire process executes at VTL 0

Secure Boot is the only root chain of trust – TPM used only for BitLocker

## Secure Launch - SINIT & ACM

Boot manager now checks for Intel TXT support & enablement
If present, TXT heap is allocated, GETSEC initialized, MTRRs configured

TxtLoadAcmModule will start the SINIT Authenticated Code Module

Authenticated Code Module provided by Intel for each CPU SKU Lives in \Windows\System32\acm.bin, or more likely, in TcbRes.wim WIM contains a VID/PID pair directory and matching ACM (KBL\_SINIT)

Parsed and loaded by TxtLoadAcmFromPackage

ACM validates chipset registers, PCH settings, PCI MMIO space If SINIT is successful, we are now in TXT mode and MLE can be loaded Otherwise, TXT launch fails, and the machine boots in *degraded mode* 

## Measured Launch Environment

MLE validates OS-sensitive data structures and execution environment Specific settings that Hyper-V may be susceptible to, MSR values, etc...

#### Transferring execution to the MLE

**BlarchLaunchMle** copies data into the **TRANSIT** section of **TcbLaunch TcbLaunch** contains the MLE embedded as the "DOS Stub" part of the PE header (the real PE binary has an **e\_lfanew** PE offset of 12KB+!)

#### After MLE execution

Calls the first ordinal of **TcbLaunch** (**BlSlEntryPoint**) which validates the input data and parameters of the **TcbLaunch** Boot Application itself Eventually calls the application entry point (**TcbMain**) of the real PE file

## TCB Launcher Flow

TcbLaunch serves as export module containing the Boot Library API

OS boot logic now lives in **TcbLoader.dll** which has two exports **TcbLoadEntry** used for boot flow, **TcbResumeEntry** for S5 resume flow

Loading Hyper-V is done by HvLoader.dll, compiled as dynamic library

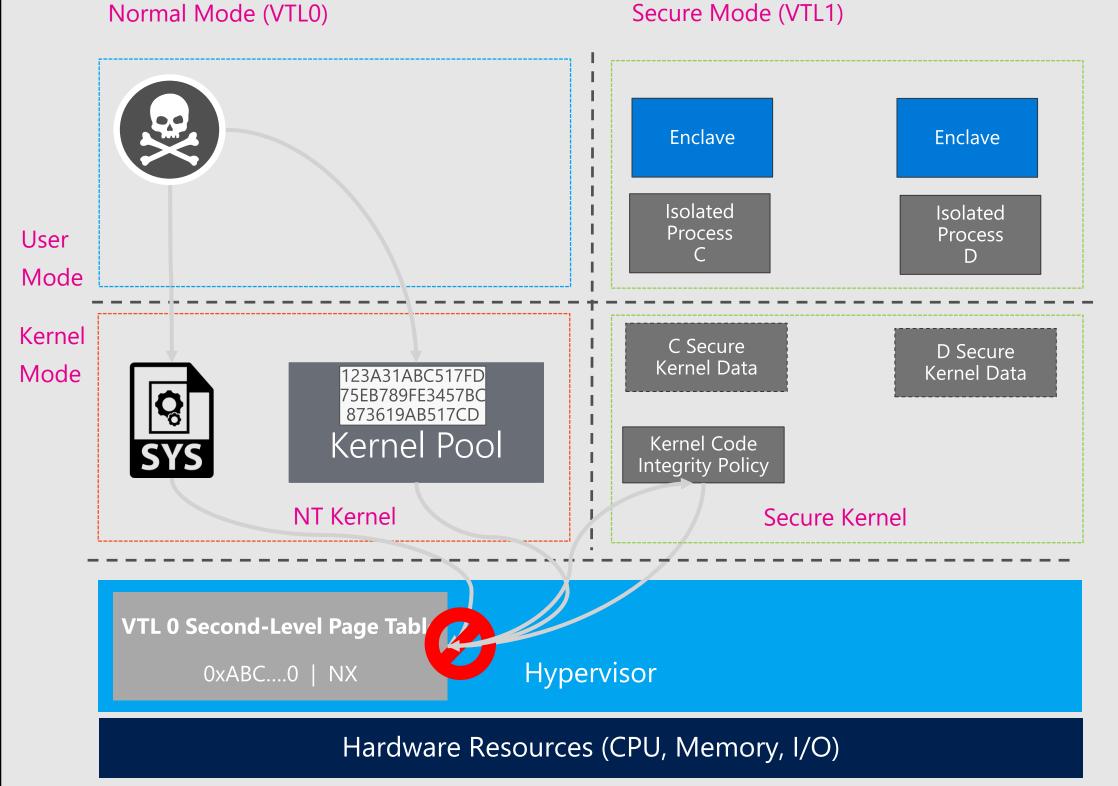
This factoring allows **TcbLoader** & **HvLoader** to share an address space State, memory and APIs can thus be global to both loaders – as if they were drivers running within the same kernel

#### HvLoader factored to handle both TXT vs. non-TXT boot

**BlGetExecutionEnvironment** used to determine memory ownership Multiple validation stages were added to all pointers and registers in ACPI tables and all other ephemeral transfer structures

### Virtual Secure Mode

Leverages nested page tables managed by VTL1 to eliminate W^X memory in VTL0 kernel-mode



#### **SLAT Used to Enforce RX-Only**

HVCl running in SK validates code pages
If valid, set GPA bits to R=1 W=0 KMX=UMX=1

#### **Mode-Based Execute (MBE) Control**

Extended-Extended Page Tables (EPT)

**XU** for user pages

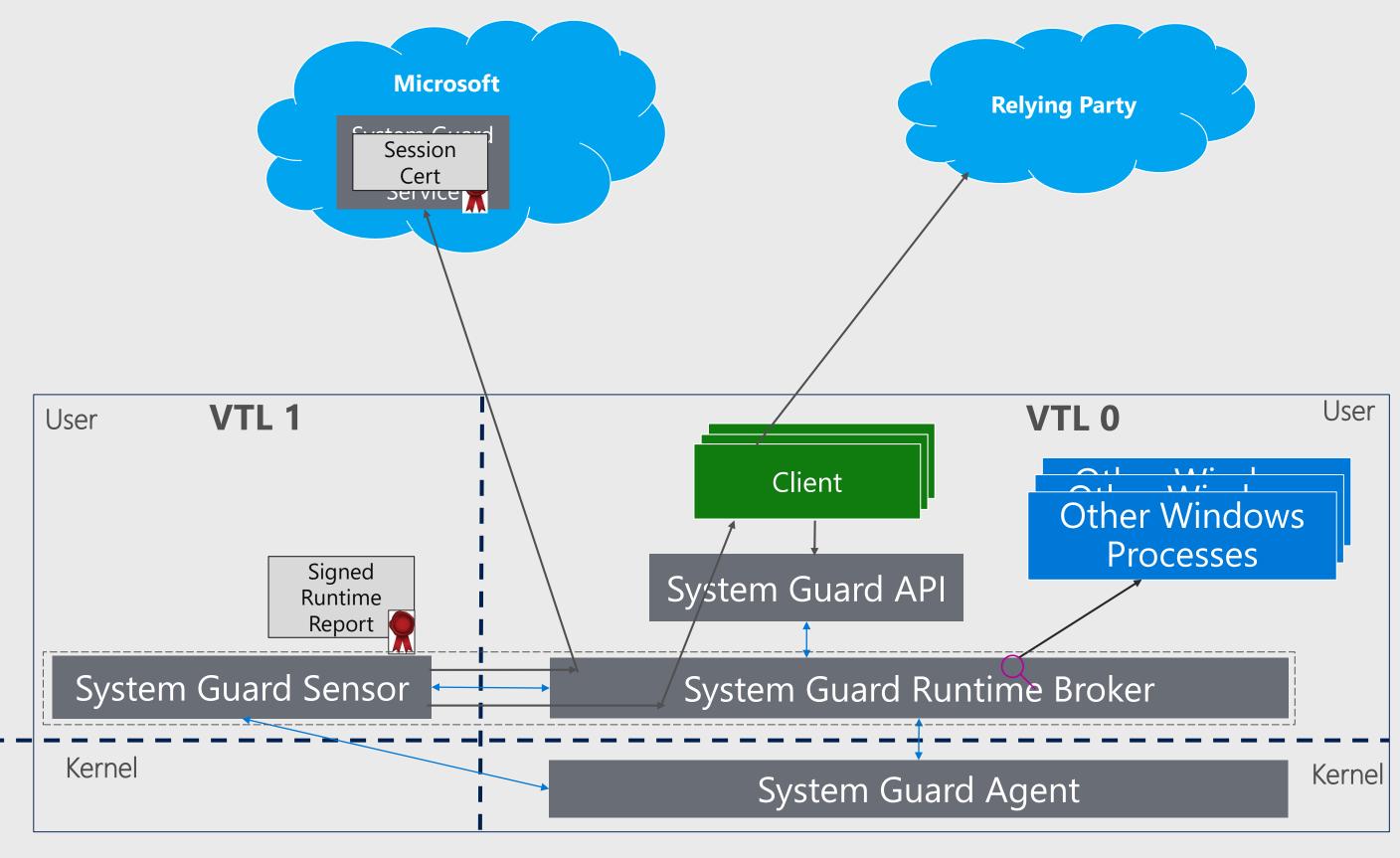
**XS** for supervisor pages

KMX and UMX become separate hardware bits

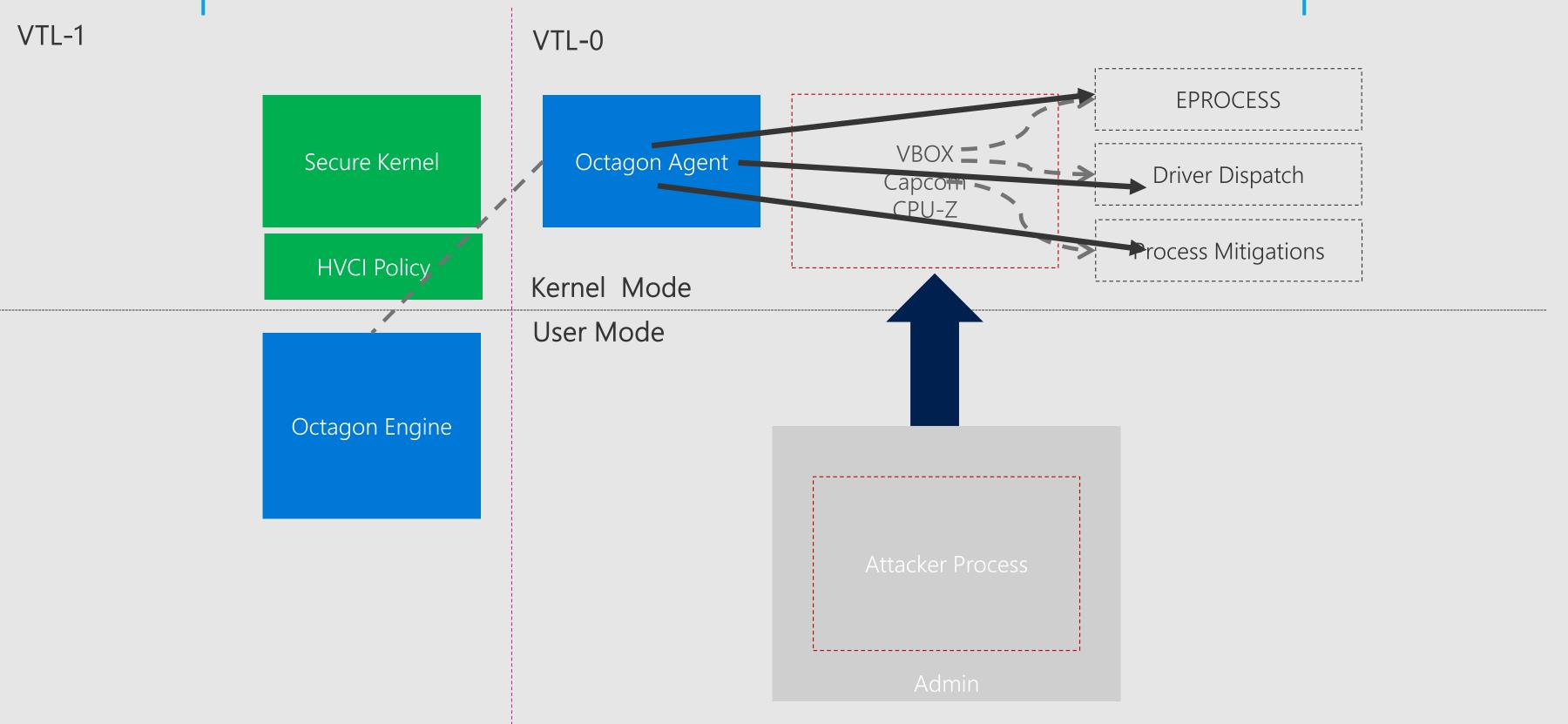
#### **VSM Enclaves**

Extensible development model for VSM
Supports VSM/SGX API documented <a href="here">here</a>
Requires code to be signed by Microsoft

## System Guard Runtime Attestation



## Example: Failed Attestation via Data Corruption



## Octagon Internals

## Activation

#### Agent Driver Component

SgrmAgent.sys loads early as a boot time driver – must win potential race against earlier name-squatting driver

Registers Kernel Extension Host to consume internal kernel API access for race-free process enumeration and thread CONTEXT acquisition

Creates \Device\MSSGRMAGENTSYS tied to "SgrmBroker" Service SID:

```
->Dacl : ->Ace[0]: ->AceType: ACCESS_ALLOWED_ACE_TYPE
->Dacl : ->Ace[0]: ->AceFlags: 0x0
->Dacl : ->Ace[0]: ->AceSize: 0x28
->Dacl : ->Ace[0]: ->Mask : 0x001201bf
->Dacl : ->Ace[0]: ->SID: S-1-5-80-3706850399-3459138796-2835936764-562029542-397710147 (Well Known Group: NT SERVICE\SgrmBroker)
```

But – IRP\_MJ\_CREATE dispatcher only allows handle to be opened by WinTCB Protected Process (Level 0x62) and restricts to a single owner per boot

## Activation

#### **Broker Activation**

SgrmBroker.exe ("System Guard Runtime Monitor Broker") loads as delayed started service, opens handle to Agent, and sends INIT IOCTL (0x9C402480)

This returns initial kernel image information (kernel base, size, checksum, timestamp, full path) which will be used to initialize the Assertion Engine

This Engine runs wrapped as a DLL – SgrmEnclave\_secure.dll for the VSM Enclave which can both execute the assertions and perform system attestation, and SgrmEnclave.dll which is a 'Shim' for assertion-only usage

Finally, registers an RPC Endpoint ("SgrmRpc") to provide Client API

## Agent Driver

#### Mailbox Interface

SgrmAgent provides a "Mailbox" IOCTL (0x9c402484) for requesting:

- User-mode read-only mapping of any piece of kernel-mode VA
- Race-free, referenced access to EPROCESS, ETHREAD, DRIVER\_OBJECT
- Suspend/resume of any thread
- User-mode copy of any physical RAM, kernel VA, or MMIO region
- Map/unmap of any file
- User-mode memory region bounds and mapped file name, if any
- User-mode CONTEXT of any thread
- Read-only MSR access
- Internal structure symbol offset lookup by name

## Agent Driver

#### Looked up Symbols

```
EPROCESS: (User)DirectoryTableBase, Peb, Token, Protection, SequenceNumber, CreateTime, SignatureLevel, MitigationFlags, SectionSignatureLevel, (InheritedFrom)UniqueProcessId, Flags2
```

TOKEN: TokenSource, Privileges, IntegrityLevelIndex, UserAndGroups, Flags

DRIVER\_OBJECT: DriverStart/Size/Name/Unload/StartIo,
MajorFunction[], and DEVICE\_OBJECT's DriverObject

These lookups can provide hints as to what type of data the runtime attestation engine is looking out for – and possibly ignoring

## System Guard Runtime Client API

#### Provide an Extensible Mechanism for WDATP & 3rd-party AM

Communicates with SYSTEM, Local Administrator, and WDATP Service SID

Internal RS4 API has a generic "Report" API and internal registration for WDATP

New RS5 API provides dedicated

GetSessionReport, GetSessionCertificate
and GetRuntimeReport functions that can be consumed by 3rd party AM
to both validate runtime assertions, as well as that they were generated by
a Secure Launch'ed VSM Enclave (or not!)

```
//Application requests an attestation session from System Guard
DeviceAttestationSession session = await
DeviceAttestationManager.CreateSessionAsync(RPID, Nonce);

//Application obtains session report signed by System Guard
Attestation Service
String jwtSessionReport = await session.GetSessionReportAsJwtAsync();

//Application can request the session X.509 certificate to verify the session report
//and obtain the key to validate the runtime report signature
Certificate sessionCertificate = await session.GetSessionCertificateAsync();

//Application obtains runtime report signed by System Guard (called multiple times per session)
String jwtRuntimeReport = await session.GetRuntimeReportAsJwtAsync();
```

## Octagon Script Model

#### LUA Engine

Assertions should be rapidly deployable (preferably through Cloud) to react to real-time threats in the wild, while at the same time provide safety and security while potentially dangerous/untrusted Ring 0 data is being parsed

Running Octagon in Ring 3 mitigates the largest privilege escalation worries, but compiled assertions in a language with pointers and a full memory access model could still result in theoretical VTL1 compromise through constructed VTL 0 kernel data

Thus, the Assertion Engine is a full-blown LUA Script Interpreter – providing acceptable performance, good security, and a rapidly iteratable language commonly used in other AM tools (Windows Defender)

## LUA<->Broker Assist Wrapper API

#### How Does LUA Talk to Win32/NT?

When the Assertion Engine is spun up, the EngHostInitialize export is called, passing in g\_BrokerAssistCallbackTbl as an array of "assist" APIs

These are made available through an AssistWrapper class which uses the Win32 CallEnclave API to exit the enclave and execute the required assist

Most of these assists map 1:1 to the IOCTLs which are provided by the Agent, but some don't: GetSystemInformation/Time, NotifyAssertion/LuaFailure, NotifyEngine, SetTimer, TraceLog

These can simply use the regular Win32/TPM API exposed to user mode

## Broker<->LUA Host Wrapper API

#### How Does the Client API -> Broker Execute LUA?

The Enclave DLL exports EngHostNotify/DispatchThread/Initialize/ Shutdown for internal use. During initialization, LUA globals can also be set, such as GLOB\_KERNEL\_IMAGE\_BASE

EngHostGetReport (RS4), EngHostGetRuntimeReport (RS5) call into the LUA functions (such as "Notify" or "ExecuteNext")

EngHostGetSessionCertificate/Report use VTL1 Enclave API (Enclave GetAttestationReport) to get VSM-signed proof that the assertions came from trusted system. EngHostCreate/DestroyAttestationClient and EngHostAttest talk to relying party to submit the attestation data

## Decompiling the LUA Script

#### Until Cloud Deployed Scripts Arrived...

The entire LUA script bytecode is in the .luaseg PE section of SgrmEnclave

No modern, well-supported LUA 5.3.4+ decompiler exists (and MS has used custom engine before)

But with some hacks... some of the logic can be exposed

```
ERROR unknown upvalue 0[1\ 0\ 14] = 1\ 0\ 15
1 0 14 = "EprocessObject HandleProcessCreation"
l_0_{15} = function(l_{123_0})
  -- function num : 0_122 , upvalues : ERROR_unknown_upvalue_0
  (ERROR_unknown_upvalue_0.EPROCESS_OBJECT_ENFORCED_ASSERTIONS):Begin("Creation")
  local l_123_1 = nil
  local 1 123 3 = nil
  local l_123_2 = l_123_0:ReadObjectCached()
  l_123_1 = (ERROR_unknown_upvalue_0.Symbols):GetMember("EPROCESS_Protection")
  local l_123_4 = (ERROR_unknown_upvalue_0.Util):ByteArrayToNumber(<math>l_123_2, l_123_1, l_123_3)
  l_123_0:SaveValue("Protection", l_123_4)
  -- DECOMPILER ERROR at PC27: Overwrote pending register: R2 in 'AssignReg'
  l_123_1 = (ERROR_unknown_upvalue_0.Symbols):GetMember("EPROCESS_SignatureLevel")
  local l_123_5 = (ERROR_unknown_upvalue_0.Util):ByteArrayToNumber(<math>l_123_2, l_123_1, l_123_3)
  l_123_0:SaveValue("SignatureLevel", l_123_5)
  -- DECOMPILER ERROR at PC43: Overwrote pending register: R2 in 'AssignReg'
  l_123_1 = (ERROR_unknown_upvalue_0.Symbols):GetMember("EPROCESS_SectionSignatureLevel")
  local l_123_6 = (ERROR_unknown_upvalue_0.Util):ByteArrayToNumber(<math>l_123_2, l_123_1, l_123_3)
  l_123_0:SaveValue("SectionSignatureLevel", l_123_6)
  -- DECOMPILER ERROR at PC59: Overwrote pending register: R2 in 'AssignReg'
  1 123 1 = (ERROR unknown upvalue 0.Symbols):GetMember("EPROCESS MitigationFlags")
  local l_123_7 = (ERROR_unknown_upvalue_0.Util):ByteArrayToNumber(<math>l_123_2, l_123_1, l_123_3)
  l_123_0:SaveValue("MitigationFlags", l_123_7)
  (ERROR_unknown_upvalue_0.EPROCESS_OBJECT_ENFORCED_ASSERTIONS):End("Creation")
  return ERROR_unknown_upvalue_0.EPROCESS_OBJECT_CREATION_COST
```

## Decompiling the Cloud-Deployed LUA Script

#### In RS5 Builds 17704 and later...

The LUA script is now deployed by cloud, dropped in \Windows\System32 \Sgrm\SgrmAssertions.bin with a catalog file (SgrmAssertions.cat) containing its hash and the Isolated User Mode EKU (1.3.6.1.4.1.311.10.3.37)

File is validated & checked against catalog hash by the enclave, not broker. It contains a LUA\_FILE\_HEADER followed by the LUA Bytecode

### Current LUA Assertions

#### Subject to Change & Rapid Iteration

- Process Image Base memory contents & integrity
- Code Integrity for MsSense.exe, MsMpEng.exe & SgrmBroker.exe
- Driver Dispatch Routines for MsSecFlt, SgrmAgent, WdFilter
- Device Objects for HackSysExtremeVulnerableDriver, mimidrv, Htsysm
- EPROCESS Data Modification (Mitigations, Token, Protection Level)
- Primary TOKEN Data Modification (Groups, IL) and SYSTEM Steal/Swap
- Firmware (specifically, SMRAM Unlock on AMD Systems w/ special MSR)

Consider these a useful "tech demo" for the initial implementation as obviously the checks are nowhere exhaustive yet – the magic behind the technology isn't what hardcoded LUA script runs – it's the engine around it

### Current LUA Assertion GUIDs

#### Encoded in Runtime Report

- FIRMWARE\_ENFORCEMENT/AUDITMODE\_GUID 5fd851c7-e688-4887-bf30-13d2c6b6d6c3 b9a781f3-d473-4e10-b871-d32465c4d572
- DRIVEROBJECT\_DISPATCHES\_GUID0efb8b25-8b47-4993-8a44-69e4b732c105
- BLACKLISTED\_DEVICES\_GUID9817a40a-69b7-4e95-af06-4eef53005660
- PRIMARY\_TOKEN\_ENFORCEMENT/AUDIT\_GUID 33b38db6-f83a-4709-bfc0-1b917d03b2bf 1200f800-eb82-4abf-868f-ef03acb58fd7
- EPROCESS\_OBJECT\_ENFORCEMENT/AUDIT\_GUID to version mode status e57680b5-3440-47cb-b9dc-49c0ae9db073 e8e0e9a8-6238-4331-a4a5-06779f2bd033
- CODE\_INTEGRITY\_GUID30dafe52-80ac-4530-a388-6507719e4e5e

```
S C:\Users\Admin\Desktop\Invoke-Sgrm> $runtimeReport.Payload.context.assertions
                : 30dafe52-80ac-4530-a388-6507719e4e5e
version
                : audit
                : good
                : b9a781f3-d473-4e10-b871-d32465c4d572
                : enforce
status
                : good
                : 5fd851c7-e688-4887-bf30-13d2c6b6d6c3
version
                : good
                : 1200f800-eb82-4abf-868f-ef03acb58fd7
                : good
lastExecutionTime : 1530546013
                : e8e0e9a8-6238-4331-a4a5-06779f2bd033
                : audit
                : 33b38db6-f83a-4709-bfc0-1b917d03b2bf
                : enforce
status
                : good
lastExecutionTime : 1530546013
                : e57680b5-3440-47cb-b9dc-49c0ae9db073
                : enforce
```

## Cloud Communications

#### Broker for HTTP Network Communications

Connecting to an HTTPS server is not only impossible from a Protected Process, it is also a bad idea

But Session Reports require connecting to sgrm.microsoft.com to send an HTTP POST request with the attestation (sent to /v1.0/Attestation)

Hence, a custom sandboxed process (SgrmLpac.exe) is used, which exposes an RPC interface that is used to connect to the service – it runs as a Low Privilege App Container (LPAC), having barely any attack surface

Nit: provides an interesting "living off the land" HTTP POST RPC API...

## Octagon Threat Model

## Enclave Technologies

#### Enclave models

Windows supports two enclave types: Intel SGX and Microsoft VSM

#### Intel SGX

Fully encrypted memory and bus traffic, resilient to SMM (but not ME) <u>but</u> requires special UEFI enablement, Intel Sky Lake+ SKU, AESM service

#### Microsoft VSM

Susceptible to runtime SMM attacks and bus snooping/RAM analysis but runs on any CPU platform with hardware virtualization support

#### Physical attacker with a quantum laser bit flipper is out of scope

SMM attacks are a realistic attack vector, however. What else is there?

## Attack Surface

#### TPM assumed trusted

Reliance on it for reports, attestation, secret sealing through PCRs

#### Hypervisor assumed trusted

VTL1 security and secrets vulnerable to Hyper-V bugs and VSM bugs

#### SMM assumed trusted

Unmeasured SMM runtime components, privileged SMI handlers, access outside SMRAM-backed communication buffers can mess with TXT

#### Management Engine assumed trusted

TXT guarantees and IOMMU boundaries vulnerable to ME vulnerabilities

## Limitations

#### Vulnerable to race conditions

Any runtime assertion technology will be limited by an attacker restoring normal/expected state within the timing window of the assertion (already an issue with PatchGuard today

## Privileged runtime VTL 0 attacker can theoretically supply 'fake' data

Measured data for runtime report comes through VTL 0 / Ring 3 broker process and provided by VTL 0 / Ring 0 driver

#### DRTM and runtime assertions can only measure/verify 'known' unknowns

A magic "God Mode" MSR won't be measured if nobody knows it exists. Similarly, data-only corruption techniques and effects must be known

## Defense-in-Depth

#### Fast I/O Dispatch with probe & locked MDLs reduces race conditions

Memory-based assertions to detect data attacks are backed by Memory Descriptor Lists mapped into user mode, simulating a shared data buffer

#### Ring 3 Broker runs as Windows TCB-level protected process

Minimizes risks of user-mode-only admin attacks, requiring either a protected process bypass or kernel-driven attack

#### Self-assertion and validation of Octagon's own components

These checks validate that Octagon components haven't been compromised, forcing a successful bypass to have to win a Time-of-Check-To-Time-of-Use (TOCTTOU) race

## Future Improvements

#### Cloud-provisioned scripts with encryption and/or obfuscation

Would create information asymmetry for attackers wishing to fake assertion (must know every possible assertion – a single one leads to detection) – this model has proven successful for PatchGuard

#### Hypervisor-backed intercepts and guarantees

Gathering data (or enhancing existing data) through hypervisor could reduce reliance on VTL 0 truth provenance

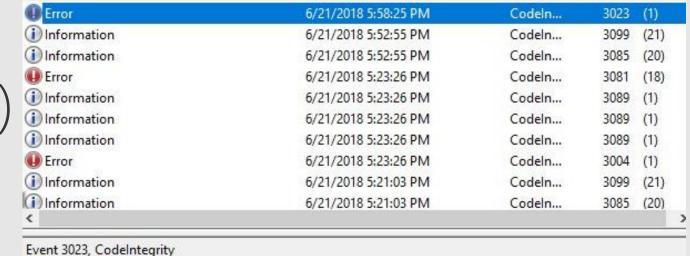
### Improve measurement, signing and attestation guarantees of platform

Vendors continuously innovating with stronger guarantees at firmware and chipset level

## Improving Kernel Security

#### VBS enabled by default in all Windows SKUs

Needs Kaby Lake+ due to perf reasons (MBEC) Activates Kernel CFG, DMA Protection, etc..



#### Blocking malicious, vulnerable, abused drivers

Windows is unable to verify the integrity of the file \Device\HarddiskVolume3\Windows\Temp\cpuz143 \cpuz143\_x64.sys because the signing certificate has been revoked. Check with the publisher to see if a new signed version of the kernel module is available.

The error code 0x11C7 (4551) occured during initialization

General Details

With HVCI, drivers that are known sources of abuse blocked via blacklist Future is to enable more secure signing model (WHQL, EV) by default

#### Read-only kernel memory

HyperGuard enforces memory/MSRs as RO

Looking at expanding to protect current attestation points

## Improving Boot Security

#### System Guard with DRTM

Hardware support for DRTM and measured boot available in WIP

OEM support for TXT is coming in future devices

#### SMM Isolation and Attestation

SMM is a key attack vector for TXT and VBS Working with IHVs to use page table protection against SMM attacks targeting OS & HV memory

Ability to attest memory protection is available

Memory integrity

Prevents attacks from inserting malicious code into high-security processes.



#### Learn more

#### Firmware protection

Windows Defender System Guard is protecting your device from compromised firmware.

#### Learn more

#### Windows Defender Credential Guard

Credential Guard is protecting your account login from attacks.

#### Learn more

#### Memory access protection

Protects your device's memory from attacks by malicious external devices.

Learn more

#### Early boot DMA protection

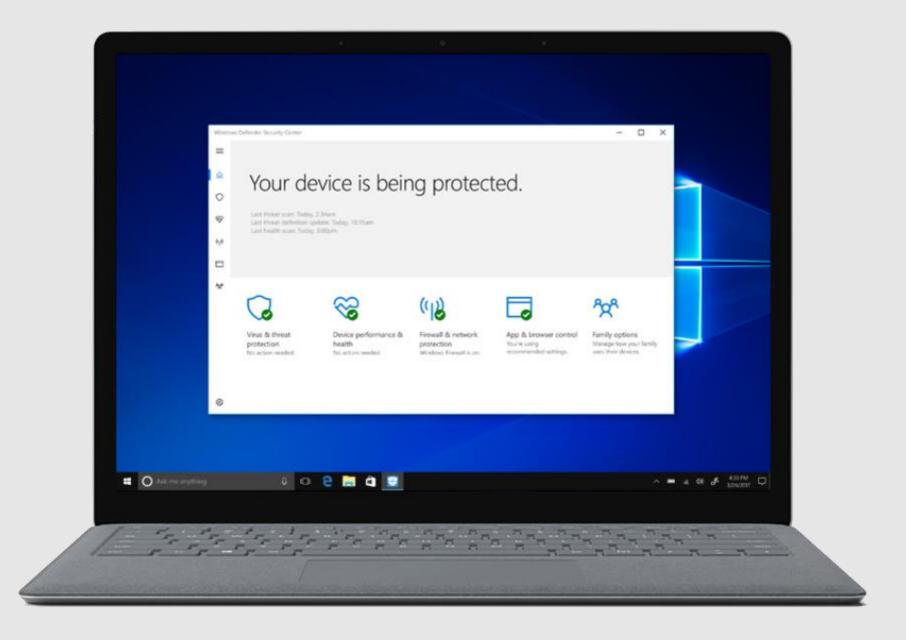
DMA-r, Early boot IOMMU, and BME Clear are all supported in WIP

## Wrap up

## Windows security promises are increasing

System Guard Runtime Attestation is critical to providing transparency to users & cloud services

Available now in 1803/1809 with a public attestation API coming soon



Aspirational security promises are the guiding principles for security investments